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Source: Motorola**Title:** Layer III MAP Message Body in Protection Mode 2**Document for:****Agenda Item:**

1. Proposal

We propose that in protection mode 2 of layer III MAP security, we encrypt first and then add a MAC for integrity. The reason for doing it in this order is that integrity can be checked without the need to decrypt first, so a false MAP message can be discarded with much less computation. This more efficient integrity protection provides a degree of protection against denial of service attack by flooding a node with false MAP messages.

We propose to replace TSGS3-000312 with the following:

7.4.2.3 Protection Mode 2

The Layer III Message Body in protection mode 2 takes the following form:

$\text{TVP} \parallel E_{K_{\text{SXY}(\text{con})}}(\text{Cleartext}) \parallel H_{K_{\text{SXY}(\text{int})}}(\text{TVP} \parallel \text{MAP Header} \parallel \text{Security Header} \parallel \text{Ciphertext})$

where "Cleartext" is the original MAP message in cleartext. Message confidentiality is achieved by encrypting cleartext, TVP and integrity check value with the confidentiality session key K_{SXY(con)}. Authentication of origin and message integrity are achieved by applying the message authentication code (MAC) function H to the concatenation of Time Variant Parameter TVP, MAP Header, Security Header and Ciphertext. The integrity is performed on the encrypted message so that integrity can be checked before decryption.

[Note1: There is need for replay protection of Layer III messages; it is envisaged to use TVP for this purpose. The precise definition of the use of TVP is ffs.]