Source: SA WG3 (Security)

Title: CRs to 33.203: Correcting the text on sending an authentication

response (Rel-5 / Rel-6)

Document for: Approval

Agenda Item: 7.3.3

| SA Doc number | Spec | CR | Rev | Phase | Subject | Cat | Version-Current | SA WG3 Doc number | Workitem |
|------------------|--------|-----|-----|-------|---|-----|-----------------|----------------------|----------|
| SP-030596 | 33.203 | 047 | 1 | | Correcting the text on sending an authentication response | F | 5.7.0 | S3-030799 | IMS-ASEC |
| SP-030596 | 33.203 | 048 | 1 | | Correcting the text on sending an authentication response | A | 6.0.0 | S3-030800 | IMS-ASEC |

| 18''' – 21 st N | | | , , , , , , , , , , , , , , , , , , , | | | | | | | CR-Form-v7 |
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| CHANGE REQUEST | | | | | | | | | | GIV-I GIIII-VI |
| ж | TS | 33.203 | CR <mark>047</mark> | ≋rev | 1 | ж | Current versi | on: | 5.7.0 | * |
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| Proposed change affects: UICC apps# ME X Radio Access Network Core Network X | | | | | | | | | | |
| Title: | * | Correcting | g the text on se | nding an authe | enticat | ion r | esponse | | | |
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| Reason for change: # Parts of the text states that RES is sent from the UE to the S-CSCF rather than a response calculated from RES and other parameters. There are CN1 CRs pending the approval of the text in this CR. | | | | | | | | | | |
| Summary of o | change | param | arified that an a aters is sent fo ding RES in the | r the UE to the | | | | | | |
| Consequence not approved | | | patiable implen RFC 3310 as c | | possil | ble w | vithout change | e. Spe | cification | n will not |
| Clauses affec | rted: | ¥ 511 | , 6.1.1, 6.1.3 | | | | | | | |
| Other specs affected: | neu. | ¥ X X | Other core specificates O&M Specificates | tions | ж | TS 2 | 4.229, TS 24 | .228 | | |
| Other comme | ents: | * | | | | | | | | |

5.1.1 Authentication of the subscriber and the network

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An IM-subscriber will have its subscriber profile located in the HSS in the Home Network. The subscriber profile will contain information on the subscriber that may not be revealed to an external partner, cf. [3]. At registration an S-CSCF is assigned to the subscriber by the I-CSCF. The subscriber profile will be downloaded to the S-CSCF over the Cx-reference point from the HSS (Cx-Pull). When a subscriber requests access to the IP Multimedia Core Network Subsystem this S-CSCF will check, by matching the request with the subscriber profile, if the subscriber is allowed to continue with the request or not i.e. Home Control (Authorization of IM-services).

All SIP-signaling will take place over the PS-domain in the user plane i.e. IP Multimedia Core Network Subsystem is essentially an overlay to the PS-domain. Hence the Visited Network will have control of all the subscribers in the PS-domain i.e. Visited Control (Authorization of bearer resources) since the Visited Network provides the subscriber with a transport service and its associated QoS.

For IM-services a new security association is required between the mobile and the IMS before access is granted to IM-services.

The mechanism for mutual authentication in UMTS is called UMTS AKA. It is a challenge response protocol and the AuC in the Home Stratum derives the challenge. A Quintet containing the challenge is sent from the Home Stratum to the Serving Network. The Quintet contains the expected response XRES and also a message authentication code MAC. The Serving Network compares the response from the UE with the XRES and if they match the UE has been authenticated. The UE calculates an expected MAC, XMAC, and compares this with the received MAC and if they match the UE has authenticated the Serving Network.

The AKA-protocol is a secure protocol developed for UMTS and the same concept/principles will be reused for the IP Multimedia Core Network Subsystem, where it is called IMS AKA.

Note: Although the method of calculating the parameters in UTMS AKA and IMS AKA are identical, the parameters are transported in slightly different ways. In UMTS, the UE's response RES is sent in the clear, while in IMS RES is not sent in the clear but combined with other parameters to form an authentication response and the authentication response is sent to the network (as described in [17]).

The Home Network authenticates the subscriber at anytime via the registration or re-registration procedures.

6.1.1 Authentication of an IM-subscriber

Before a user can get access to the IM services at least one IMPU needs to be registered and the IMPI authenticated in the IMS at application level. In order to get registered the UE sends a SIP REGISTER message towards the SIP registrar server i.e. the S-CSCF, cf. Figure 1, which will perform the authentication of the user. The message flows are the same regardless of whether the user has an IMPU already registered or not.

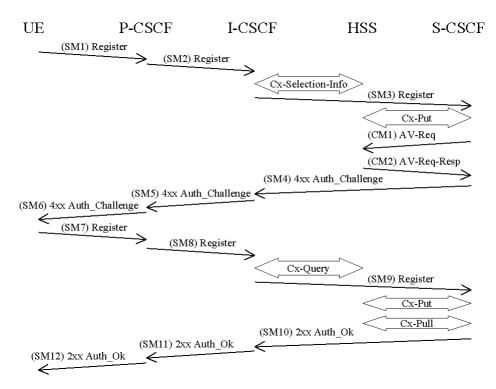


Figure 4: The IMS Authentication and Key Agreement for an unregistered IM subscriber and successful mutual authentication with no synchronization error

The detailed requirements and complete registration flows are defined in [8] and [11].

SMn stands for SIP Message n and CMm stands for Cx message m which has a relation to the authentication process:

SM1: REGISTER(IMPI, IMPU)

In SM2 and SM3 the P-CSCF and the I-CSCF respectively forwards the SIP REGISTER towards the S-CSCF.

After receiving SM3, if the IMPU is not currently registered at the S-CSCF, the S-CSCF needs to set the registration flag at the HSS to initial registration pending. This is done in order to handle mobile terminated calls while the initial registration is in progress and not successfully completed. The registration flag is stored in the HSS together with the S-CSCF name and user identity, and is used to indicate whether a particular IMPU of the user is unregistered or registered at a particular S-CSCF or if the initial registration at a particular S-CSCF is pending. The registration flag is set by the S-CSCF sending a Cx-Put to the HSS. If the IMPU is currently registered, the S-CSCF shall leave the registration flag set to *registered*. At this stage the HSS has performed a check that the IMPI and the IMPU belong to the same user.

Upon receiving the SIP REGISTER the S-CSCF CSCF shall use an Authentication Vector (AV) for authenticating and agreeing a key with the user. If the S-CSCF has no valid AV then the S-CSCF shall send a request for AV(s) to the HSS in CM1 together with the number m of AVs wanted where m is at least one.

CM1: Cx-AV-Reg(IMPI, m) Upon receipt of a request from the S-CSCF, the HSS sends an ordered array of *n* authentication vectors to the S-CSCF using CM2. The authentication vectors are ordered based on sequence number. Each authentication vector consists of the following components: a random number RAND, an expected response XRES, a cipher key CK, an integrity key IK and an authentication token AUTN. Each authentication vector is good for one authentication and key agreement between the S-CSCF and the IMS user.

CM2

Cx-AV-Req-Resp(IMPI, RAND1||AUTN1||XRES1||CK1||IK1,....,RANDn||AUTNn||XRESn||CKn||IKn)

When the S-CSCF needs to send an authentication challenge to the user, it selects the next authentication vector from the ordered array, i.e. authentication vectors in a particular S-CSCF are used on a first-in / first-out basis.

The S-CSCF sends a SIP 4xx Auth_Challenge i.e. an authentication challenge towards the UE including the challenge RAND, the authentication token AUTN in SM4. It also includes the integrity key IK and the cipher key CK for the P-CSCF. Draft ietf sip digest aka 01RFC 3310 [17] specifies how the fields to populate the corresponding parameters of an authenticatione challenge.

The verification of the SQN by the USIM and ISIM will cause the UE to reject an attempt by the S-CSCF to re-use a AV. Therefore no AV shall be sent more than once.

NOTE: This does not preclude the use of the normal SIP transaction layer re-transmission procedures.

SM4:

4xx Auth_Challenge(IMPI, RAND, AUTN, IK, CK)

When the P-CSCF receives SM5 it shall store the key(s) and remove that information and forward the rest of the message to the UE i.e.

SM6:

4xx Auth_Challenge(IMPI, RAND, AUTN)

Upon receiving the challenge, SM6, the UE takes the AUTN, which includes a MAC and the SQN. The UE calculates the XMAC and checks that XMAC=MAC and that the SQN is in the correct range as in [1]. If both these checks are successful the UE <u>uses</u>calculates the response, RES; and some other parameters to calculate an authentication response. This response is puts it into the Authorization header and sentes it back to the registrar in SM7. RFC 3310 Draft ietf-sip-digest aka 01 [17] specifies howthe fields to populate the the corresponding parameters of the response. It should be noted that the UE at this stage also computes the session keys CK and IK.

SM7:

REGISTER(IMPI, <u>Authentication response</u>RES)

The P-CSCF forwards the <u>authentication response</u>RES in SM8 to the I-CSCF, which queries the HSS to find the address of the S-CSCF. In SM9 the I-CSCF forwards the <u>authentication response</u>RES to the S-CSCF.

Upon receiving SM9 containing the response, the S-CSCF retrieves the active XRES for that user and uses this to check the <u>authentication</u> response sent by the UE as described in <u>RFC 3310 Draft ietf sip digest aka 01 [17]</u>. If the check is successful then the user has been authenticated and the IMPU is registered in the S-CSCF. If the IMPU was not currently registered, the S-CSCF shall send a Cx-Put to update the registration-flag to *registered*. If the IMPU was currently registered the registration-flag is not altered.

It shall be possible to implicitly register IMPU(s). The implicitly registered IMPU(s) all belong to the same Service Profile. All the IMPU(s) being implicitly registered shall be delivered by the HSS to the S-CSCF and subsequently to the P-CSCF. The S-CSCF shall regard all implicitly registered IMPU(s) as registered IMPU(s).

When an IMPU has been registered this registration will be valid for some period of time. Both the UE and the S-CSCF will keep track on a timer for this purpose but the expiration time in the UE is smaller than the one in the S-CSCF in order to make it possible for the UE to be registered and reachable without interruptions. A successful registration of a

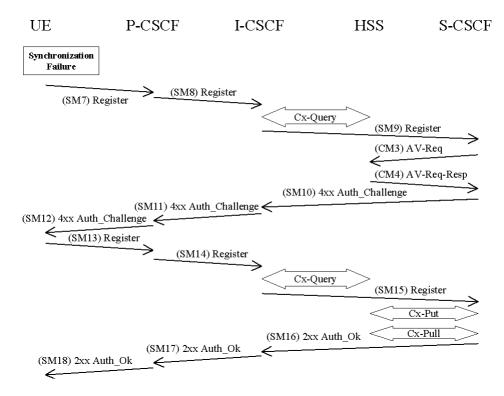
previously registered IMPU (including implicitly registered IMPUs) means the expiry time of the registration is refreshed.

It should be noted that the UE initiated re-registration opens up a potential denial-of-service attack. That is, an attacker could try to register an already registered IMPU and respond with an incorrect wrong RES authentication response and in order to make the HN de-register the IMPU. For this reason a subscriber should not be de-registered if it fails an authentication. It shall be defined by the policy of the operator when successfully registered IMPU(s) are to be de-registered.

The lengths of the IMS AKA parameters are specified in chapter 6.3.7 in [1].

6.1.3 Synchronization failure

In this section the case of an authenticated registration with synchronization failure is described. After resynchronization, authentication may be successfully completed, but it may also happen that in subsequent attempts other failure conditions (i.e. user authentication failure, network authentication failure) occur. In below only the case of synchronization failure with subsequent successful authentication is shown. The other cases can be derived by combination with the flows for the other failure conditions.



The flow equals the flow in 6.1.1 up to SM6. When the UE receives SM6 it detects that the SQN is out of range and sends a synchronization failure back to the S-CSCF in SM7. RFC 3310 Draft ietf-sip-digest-aka-01 [17] describes the fields to populate corresponding parameters of synchronization failure.

SM7: REGISTER(Failure = *Synchronization Failure*, AUTS, IMPI)

Upon receiving the *Synchronization Failure* and the AUTS the S-CSCF sends an Av-Req to the HSS in CM3 including the required number of Avs, m.

CM3: Cx-AV-Req(IMPI, RAND,AUTS, m)

The HSS checks the AUTS as in section 6.3.5 in [1]. After potentially updating the SQN, the HSS sends new AVs to the S-CSCF in CM4.

 $CM4: \\ Cx-AV-Req-Resp(IMPI, n,RAND_1||AUTN_1||XRES_1||CK_1||IK_1,....,RAND_n||AUTN_n||XRES_n||CK_n||IK_n) \\ \\ RAND_n||AUTN_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XR$

The rest of the messages i.e. SM10-SM18 including the Cx messages are exactly the same as SM4-SM12 and the corresponding Cx messages in 6.1.1.

3GPP TSG-SA-WG3 Meeting #31

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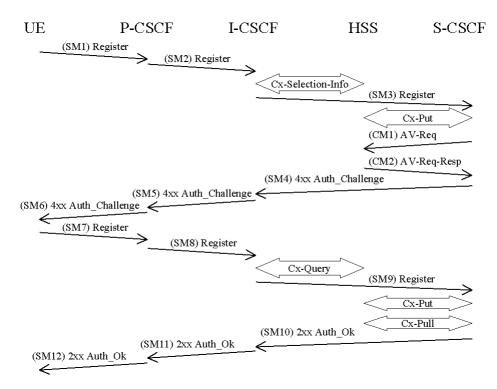


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CM2

Cx-AV-Req-Resp(IMPI, RAND1||AUTN1||XRES1||CK1||IK1,....,RANDn||AUTNn||XRESn||CKn||IKn)

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The S-CSCF sends a SIP 4xx Auth_Challenge i.e. an authentication challenge towards the UE including the challenge RAND, the authentication token AUTN in SM4. It also includes the integrity key IK and the cipher key CK for the P-CSCF. Draft ietf sip digest aka 01RFC 3310 [17] specifies how the fields to populate the corresponding parameters of an authenticatione challenge.

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NOTE: This does not preclude the use of the normal SIP transaction layer re-transmission procedures.

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4xx Auth_Challenge(IMPI, RAND, AUTN, IK, CK)

When the P-CSCF receives SM5 it shall store the key(s) and remove that information and forward the rest of the message to the UE i.e.

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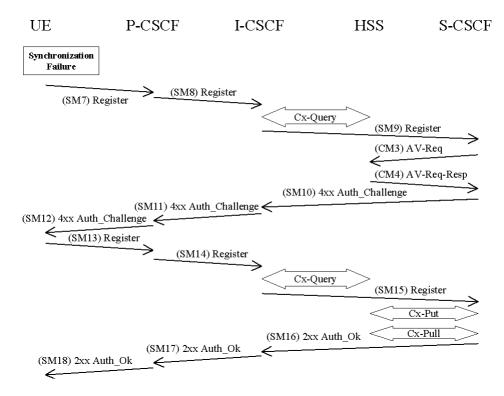
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The lengths of the IMS AKA parameters are specified in chapter 6.3.7 in [1].

6.1.3 Synchronization failure

In this section the case of an authenticated registration with synchronization failure is described. After resynchronization, authentication may be successfully completed, but it may also happen that in subsequent attempts other failure conditions (i.e. user authentication failure, network authentication failure) occur. In below only the case of synchronization failure with subsequent successful authentication is shown. The other cases can be derived by combination with the flows for the other failure conditions.



The flow equals the flow in 6.1.1 up to SM6. When the UE receives SM6 it detects that the SQN is out of range and sends a synchronization failure back to the S-CSCF in SM7. RFC 3310 Draft ietf-sip-digest-aka-01 [17] describes the fields to populate corresponding parameters of synchronization failure.

SM7: REGISTER(Failure = Synchronization Failure, AUTS, IMPI)

Upon receiving the *Synchronization Failure* and the AUTS the S-CSCF sends an Av-Req to the HSS in CM3 including the required number of Avs, m.

CM3: Cx-AV-Req(IMPI, RAND,AUTS, m)

The HSS checks the AUTS as in section 6.3.5 in [1]. After potentially updating the SQN, the HSS sends new AVs to the S-CSCF in CM4.

 $CM4: \\ Cx-AV-Req-Resp(IMPI, n,RAND_1||AUTN_1||XRES_1||CK_1||IK_1,....,RAND_n||AUTN_n||XRES_n||CK_n||IK_n) \\ \\ RAND_n||AUTN_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||IK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||CK_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XRES_n||XR$

The rest of the messages i.e. SM10-SM18 including the Cx messages are exactly the same as SM4-SM12 and the corresponding Cx messages in 6.1.1.