Ad-hoc on aSIP and NDS WIS Munich, 8th – 9th October 2000

Agenda Item: -

Source: Ericsson

Title: Authentication and protection mechanisms for IM CN SS

Document for: Discussion and decision

1 Scope and objectives

The scope for this document is to define mechanisms and trust relations for secure access to IM CN SS, IP Multimedia Core Network subsystem, in UMTS. The IM CN SS in UMTS will support IP Multimedia applications such as video, audio and multimedia conferences. 3GPP has chosen SIP, Session Initiation Protocol as the signalling protocol for creating and terminating Multimedia sessions.

The following are proposed in the document:

- 1. An authentication mechanism for IM CN SS according to the following principles
 - Re-use of UMTS AKA concept
 - Authentication performed in the HSS
- 2. A scheme, including Key Management, for protection of man-in-the-middle attack in IM CN SS
 - Confidentiality and integrity protection between the UE and the P-CSCF (Optional)
 - Integrity protection between the UE and the S-CSCF (Mandatory)
 - Inter operator protection by using the SEG and IPSec
- 3. That the IMSI is not used for User Identity instead it is proposed that the NAI as in RFC 2486 shall be used.

2 Introduction

3GPP has defined three CSCFs, Call State Control Functions, the P-CSCF, the I-CSCF and the S-CSCF where P stands for Proxy and I for Interrogating and S for Serving, cf. figure 1. All of these CSCFs will act like SIP servers. The role of the S-CSCF is to provide the (roamed or non-roamed) subscriber with service control. The S-CSCF is assigned to the subscriber at registration. Depending on the policy in the HN, Home Network, the S-CSCF is in the HN or in the VN, Visited Network. It is a requirement in 3GPP that it shall be possible to hide the network topology from other operators e.g. the number of S-CSCFs. The entry point for the VN to the HN is the I-CSCF. The I-CSCF is responsible for choosing the S-CSCF based on information the I-CSCF gets from the P-CSCF in the VN and the HSS, Home Subscriber Server, in the HN. The information the I-CSCF will need for the selection process is amongst other things the subscriber identity, VN capabilities, required capabilities of the S-CSCF, if the VN or the HN will be the SN, Serving Network etc. The P-CSCF shall enable the call control to be sent to the HN through the I-CSCF and also enable the SIP-messages to be sent to the UE in the VN.

The registration procedure of the UE is divided into three flows:

- 1. Common initiation
- 2. HN-control
- 3. VN-control

At start of the registration the UE sends a SIP-register to the P-CSCF in the VN which then sends the SIP-registration to the HN. The I-CSCF will the make the selection of the SN and S-CSCF. The subscriber data etc is sent to the S-CSCF by the I-CSCF. It is necessary that the HN can authenticate the subscriber in a secure way using e.g. a shared secret. The authentication protocol which is used in UMTS provides mutual authentication i.e. the subscriber can authenticate the serving network and the SN can authenticate the subscriber. This mechanism is reused in the proposal presented in this document. However this requires a new mode in SIP as

described in S3-000456. An analysis of different protocols for authentication is made in S3-000447 and in S3-000588.

When the UE makes a call a SIP INVITE is sent to the P-CSCF in the VN which will be forwarded to the S-CSCF which in turn forwards the INVITE based on the destination.

It is important to ensure that the subscriber gets the services he is entitled to and that he gets the services he asks for. A threat towards the UE and the HN is a man-in-the-middle attack where the attacker tries to modify the SIP-messages in whatever way his purposes might be. Therefore it is necessary to check that the SIP-messages sent from the UE have not been tampered with. This can be fulfilled by using integrity protection on the application layer using a shared secret, the Integrity Key IK, between the UE and the S-CSCF which has the subscriber data. Note that the SN i.e. the S-CSCF can be located in either VN or in HN.

The integrity protection does not protect the messages from being read by an attacker who can if no confidentiality protection is provided get information like the IP-address of the UE and what Codecs will be used in the Multimedia session etc.

The trust model proposed in this document is that 1) the subscriber trusts the VN and the HN, 2) the VN trusts the HN and 3) the HN trusts the VN and other networks it has a relation or a roaming agreement with. Since the SIP-proxies have to be able to read certain parts of the SIP-messages and assuming the trust model above hop-by-hop encryption is proposed. Either TLS or IPSec can provide hop-by-hop encryption.

Protection between the nodes within a network e.g. between the I-CSCF and the S-CSCF in the HN is open and a subject for the network owner to solve.

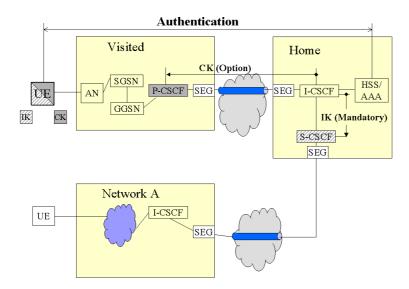


Figure 1 An overview of the architecture.

In figure 1 it is shown that the protection mechanism between the different networks is IPSec using the SEG as specified in the WI NDS. Furthermore the UE and the P-CSCF can optionally get confidentiality protection by using the shared key CK by using either IPSec or TLS. In this document WTLS is proposed. Note also that the UE and the S-CSCF share the Key IK which will be used for integrity protection between the UE and the S-CSCF. The S-CSCF is in this picture located in the HN however it is the policy of the HN that decides if the S-CSCF shall be located in either HN or VN.

3 Background

One assumption that has been made in TR 33.8xx "Access security for IP-based services" is that it is possible to reuse the AKA concept as defined in R'99 for UMTS. However this does not mean necessarily that the same algorithms are used in the multimedia domain as in UMTS.

In 23.228 v 1.1.0 it is assumed that the user is authenticated before the I-CSCF sends the Cx-Select-pull to the HSS. Therefore it is proposed in this document that the HSS authenticate the user. Another good reason for authenticating the user in the HSS is that the HN has the knowledge and control on the outcome of the authentication. Besides this allows for possible enhancements/modifications in the procedure without affecting nodes in the rest of the NW, especially those at the visited NW (P-CSCF).

The I-CSCF is also responsible for assigning an S-CSCF to the user and deciding if the home network or the visited network is chosen for service control. It is assumed in this document that the HSS also calculates a quintet, which has to be distributed to different nodes as described in this document.

It is a requirement in TR 33.8xx that the IM CN SS shall not rely on the security provided by the PS-domain. Furthermore the integrity protection shall take place between the UE and the S-CSCF according to TR33.8xx. As a working assumption it is assumed that SIP AKA is provided.

4 Description of the mechanisms

4.1 Access independence

The home environment shall make the authentication of the user. In this document only the aspects important for the UMTS standardisation using SIP AKA for IM CN SS have been evaluated. The core network also has to be IETF capable.

4.2 Inter operator protection

Since sensitive SIP signalling messages are sent between e.g. the visited network and the home network it has to be protected. In the WI SAWG3 NDS, Network Domain Security, it is specified that a Security Gateway, SEG, shall be at the border of a network, providing IP security for IP communication between different networks. It is therefore proposed in this document that it shall be mandatory to use the SEG for SIP signalling, cf. Figure 1.

4.3 User identity

It is proposed not to use the IMSI as the identifier of the user in the IM CN SS. There are several reasons for that. One being the requirement to be access independent and an other being the requirement to give user identity confidentiality in UMTS and hence the IMSI should not be used for IM CN SS. Instead it is proposed that the NAI, Network Access Identifier, cf. RFC 2486, shall be used which has the format *user@realm*. This is still FFS in S2.

4.4 Configurability and visibility

In UMTS R'99 the link between the UE and the RNC may be confidentiality protected and integrity protected. If the user receives a non-ciphered call then the user may reject the non-ciphered call this is worked out in R'00. This might also be an issue for aSIP since in this discussion paper it is proposed that the confidentiality protection between the UE and the P-CSCF is optional. The user should be notified whenever the confidentiality protection is turned off and also be given the opportunity to reject a non-ciphered SIP-session. This shall be FFS.

4.5 Registration

In this proposal the authentication takes place in the HSS/AAA according to the scheme below. The proposal seems to meet the current requirements in TR 33.8xx. Note that the CK, the Cipher Key, is sent from the HSS to the P-CSCF via the I-CSCF together with a challenge using a mode called SIP AKA described in S3-000456. The subscriber/user can check that the challenge originates from the HN and respond with a RES, which is checked by the HSS. The subscriber can derive the CK from the challenge and the shared secret.

4.5.1 Common initiation of the registration

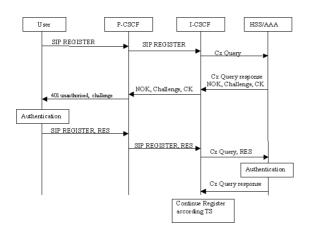


Figure 2 The user is authenticated in the HSS/AAA and the CK is sent to the P-CSCF. The HN is authenticated by the UE.

4.5.2 The S-CSCF is located in the HN

When the authentication has succeeded the I-CSCF selects an S-CSCF and in this example it is located in the HN. The HSS sends the Integrity Key, IK, to the S-CSCF. The subscriber has derived the IK from the challenge. All subsequent SIP messages shall be integrity protected by using appropriate algorithms and the IK.

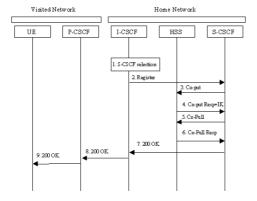


Figure 3 The registration process proceeds and the HSS sends the IK to the S-CSCF

4.5.3 The S-CSCF is located in the HN

When the authentication has succeeded the I-CSCF in the VN selects an S-CSCF. The HSS sends the Integrity Key, IK, to the S-CSCF. The subscriber has derived the IK from the challenge. All subsequent SIP messages shall be integrity protected by using appropriate algorithms and the IK.

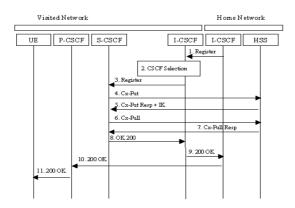


Figure 4 The registration process when the S-CSCF is located in the VN.

4.6 Integrity protection

It is proposed in this document that the SIP-messages like e.g. INVITE is integrity protected on the application layer. In draft-ietf-sip-rfc2543bis-01 it is a requirement that all SIP implementations should support PGP-based encryption and may implement other schemes. PGP may also be used for authentication. In TR 33.8xx an analysis over different protocols like PGP is outlined. However it is concluded that PGP being a PKI based system does not fulfil the 3GPP requirements. Hence another scheme, which does not necessarily have to be a 3GPP scheme, has to be defined.

The integrity key IK is derived by the 3GPP AKA scheme and sent from the HSS to the S-SCSF. The UE in turn can calculate the IK (and CK) from the challenge sent from the HSS.

If the S-CSCF is in visited network then the IK has to be sent through the I-CSCF in home to the I-CSCF in the visited domain and finally to the S-CSCF in the visited network.

For example, if the SIP request is to be:

INVITE sip:watson@boston.bell-telephone.com SIP/2.0

Via: SIP/2.0/UDP 169.130.12.5

Authorization: PGP version=5.0, signature=...

From: A. Bell <sip:a.g.bell@bell-telephone.com>
To: T. A. Watson <sip:watson@bell-telephone.com>

Call-ID: 187602141351@worcester.bell-telephone.com

Subject: Mr. Watson, come here.

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Content-Type: application/sdp

Content-Length: ...

v=0

o=bell 53655765 2353687637 IN IP4 128.3.4.5

s=Mr. Watson, come here.

t=0 0

c=IN IP4 135.180.144.94

m=audio 3456 RTP/AVP 0 3 4 5
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Then the text in italic and in bold is integrity protected (using a canonical form) and the underlined text is not protected since some parts of it will be changed by proxies (of course the Authorization part has to be unchanged).

4.7 Confidentiality protection

As an option it is proposed that confidentiality and integrity protection take place between the UE and the P-CSCF. Using appropriate mechanisms defined in WTLS may protect the SIP messages between the UE and the P-CSCF. WTLS may then use the CK as a pre-master secret, which in turn is used for calculation of a key block containing material for encryption key, MAC-secret and IV. Using this option then the UE and the P-CSCF will have a security association. Furthermore it fulfils the requirement in 33.8xx that the protection in IM CN SS shall not rely on the protection mechanisms in the PS-domain. There is a requirement in the TR 33.8xx stating that it shall be an option that confidentiality takes place between the UE and the S-CSCF but since several fields in the SIP messages can not be encrypted we have to rely on hop-by-hop protection using either TLS or IPSec. It is proposed that WTLS be used since it is optimised for low-bandwidth bearer networks with relatively long latency, which is not the case for IPSec. The exact mechanisms and the use of WTLS is FFS.

5 Discussion

There are some issues that have not been discussed in this document but they are worthwhile to mention here.

Since in this proposal the SEG is used for protection between 1) P-CSCF and I-CSCF and 2) I-CSCF and I-CSCF there is an issue regarding interoperability between a 3GPP-network and a 3GPP2-network since 3GPP2 does not support a SEG.

Furthermore no end-to-end solution is provided in this proposal i.e. UE-UE. However using security mechanisms like e.g. an RTP-cipher for the multimedia application itself can possibly solve this. Then a key management scheme has to be defined for end-to-end security.

It is assumed in this document that the principles according to S3-000456 are available i.e. that an AKA mode exists for SIP. This mode does not exist today.