# 3GPP TSG RAN WG1#11 San Diego USA, Feb. 29<sup>th</sup> – March 3<sup>rd</sup>, 2000

Document R1-00-0237 e.g. for 3GPP use the format TP-99xxx or for SMG, use the format P-99-xxx

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<u>Reason for</u> <u>change:</u>	For the mos	t commonality be	tween F	DD and TI	DD TFCI basis	s, this C	R is proposed	
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# 4.3 Coding for layer 1 control

## 4.3.1 Coding of transport format combination indicator (TFCI)

Encoding of the TFCI bits depends on the number of them. If there are 6-10 bits of TFCI the channel encoding is done as described in section 4.3.1.1. Also specific coding of less than 6 bits is possible as explained in section 4.3.1.2.

### 4.3.1.1 Coding of long TFCI lengths

The TFCI bits are encoded using a (32, 10) sub-code of the second order Reed-Muller code. The coding procedure is as shown in figure 4.3.3.1-1.

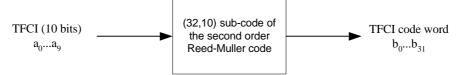


Figure 4.3.3.1-1: Channel coding of TFCI bits

TFCI is encoded by the (32,10) sub-code of second order Reed-Muller code. The code words of the (32,10) sub-code of second order Reed-Muller code are linear combination of some among 10 basis sequences. The basis sequences are as follows in table 4.3.1-1.

I	M <sub>i,0</sub>	M <sub>i,1</sub>	M <sub>i,2</sub>	M <sub>i,3</sub>	M <sub>I,4</sub>	M <sub>i,5</sub>	M <sub>i,6</sub>	M <sub>i,7</sub>	M <sub>i,8</sub>	<b>M</b> i,9
0	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	0	0	0	0
1	<u>0</u> 4	<u>1</u> 0	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	1	0	0	0
2	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	0	0	0	1
3	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	1	0	1	1
4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	0	0	0	1
5	<u>0</u> 4	<u>1</u> 0	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	0	0	1	0
6	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	0	1	0	0
7	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	0	1	1	0
8	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	1	1	1	0
9	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	1	0	1	1
10	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	0	0	1	1
11	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	0	1	1	0
12	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	0	1	0	1
13	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	1	0	0	1
14	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	1	1	1	1
15	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	1	1	0	0
16	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	1	1	0	1
17	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	1	0	1	0
18	<u>0</u> 4	<u>0</u> 0	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	0	1	1	1
19	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	0	1	0	1
20	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	0	0	1	1
21	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	0	1	1	1
22	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	0	1	0	0
23	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	1	1	0	1
24	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	1	0	1	0
25	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	1	0	0	1
26	<u>0</u> 4	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	0	0	1	0
27	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	1	1	0	0
28	<u>0</u> 4	<u>1</u> 0	<u>1</u> 1	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	1	1	1	0
29	<u>1</u> 4	1	1	1	1					
30	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	0	0	0	0
31	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	1	0	0	0

Table 4.3.1-1: Basis sequences for (32,10) TFCI code

For TFCI bits  $a_0, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9$  ( $a_0$  is LSB and  $a_9$  is MSB), the output code word bits  $b_i$  are given by:  $b_i = \sum_{n=0}^{9} (a_n \times M_{i,n}) \mod 2$ 

where i=0...31.  $N_{TFCI}$ =32.

# 4.3.1.2 Coding of short TFCI lengths

#### 4.3.1.2.1 Coding very short TFCIs by repetition

If the number of TFCI bits is 1 or 2, then repetition will be used for coding. In this case each bit is repeated to a total of 4 times giving 4-bit transmission ( $N_{TFCI}=4$ ) for a single TFCI bit and 8-bit transmission ( $N_{TFCI}=8$ ) for 2 TFCI bits. In the case of two TFCI bits denoted  $b_0$  and  $b_1$  the TFCI word shall be {  $b_0, b_1, b_0, b_1, b_0, b_1, b_0, b_1$  }.

#### 4.3.1.2.2 Coding short TFCIs using bi-orthogonal codes

If the number of TFCI bits is in the range 3 to 5 the TFCI bits are encoded using a (16, 5) bi-orthogonal (or first order Reed-Muller) code. The coding procedure is as shown in figure 4-8.

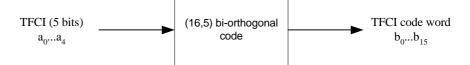


Figure 4-8: Channel coding of short length TFCI bits

The code words of the (16,5) bi-orthogonal code are linear combinations of 5 basis sequences as defined in table 4.3.1-2 below.

i	M <sub>i,0</sub>	M <sub>i,1</sub>	M <sub>i,2</sub>	M <sub>i,3</sub>	M <sub>i,4</sub>
0	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0
1	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0
2	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0
3	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0
4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0
5	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0
6	<u>1</u> 4	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0
7	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4
8	<u>1</u> 4	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4
9	<u>0</u> 1	<u>1</u> 0	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4
10	<u>1</u> 4	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4
11	<u>0</u> 1	<u>0</u> 0	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4
12	<u>1</u> 4	<u>0</u> 1	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4
13	<u>0</u> 4	<u>1</u> 0	<u>1</u> 4	<u>1</u> 4	<u>1</u> 1
14	<u>1</u> 4				
15	<u>0</u> 1	<u>0</u> 0	<u>0</u> 0	<u>0</u> 0	<u>1</u> 0

Table 4.3.1-2: Basis sequences for (16,5) TFCI code

For TFCI information bits  $a_0$ ,  $a_1$ ,  $a_2$ ,  $a_3$ ,  $a_4$  ( $a_0$  is LSB and  $a_4$  is MSB), the ), the output code word bits  $b_j$  are given by:  $b_i = \sum_{n=0}^{4} (a_n \times M_{i,n}) \mod 2$ 

where i=0...15.  $N_{TFCI}$ =16.