# 25.222 CR 008rev2 <br> Current Version: 3.0.0 

GSM (AA.BB) or $3 G(A A . B B B)$ specification number $\uparrow$
$\uparrow$ CR number as allocated by MCC support team
For submission to: TAG-RAN\#6
list expected approval meeting \# here $\uparrow$

| for approval | $\mathbf{X}$ |
| ---: | ---: |
| information |  |



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## Work item:

| Category: | F | Correction | X | Release: | Phase 2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | A | Corresponds to a correction in an earlier release |  |  | Release 96 |
| (only one category | B | Addition of feature |  |  | Release 97 |
| shall be marked | C | Functional modification of feature |  |  | Release 98 |
| with an $X$ ) | D | Editorial modification |  |  | Release 99 |
|  |  |  |  |  | Release 00 |


| Reason for | Description of TFCI coding for normal and short length cases is quite different and <br> does not specify coding process completely. Also dependency on OVSF sequences <br> change: <br> defined externally which should be removed. |
| :--- | :--- |
|  | Remove narrative text in 4.3.1 <br> Change 4.3.1.1 to more specification-like form |
|  | Add 4.3.1.2.1, 4.3.1.2.2 section headings to make clearer different rules <br> Change order of basis vectors and input bit positions in tables 4.3.1-1 and 4.3.1-2 to <br> correspond with changes in FDD TFCI coding. |

## Clauses affected: $\quad 4.2 .13,4.3 .1,4.3 .1 .1,4.3 .1 .2$

## Other specs affected:

Other 3G core specifications Other GSM core specifications MS test specifications BSS test specifications O\&M specifications


Other $\quad$ Due to a limitation of MS-Word the original 4.3.1-1 cannot be deleted in markup mode.

### 4.2.13 Transport format detection

Transport format detection can be performed both with and without Transport Format Combination Indicator (TFCI). If a TFCI is transmitted, the receiver detects the transport format combination from the TFCI. When no TFCI is transmitted, so called blind transport format detection may beis used, i.e. the receiver side uses the possible transport format combinations as a priori information.

### 4.2.13.1 Blind transport format detection

Blind transport format detection may be performed in the receiver by trying all possible combinations of the transport format.

### 4.2.13.2 Explicit transport format detection based on TFCI

### 4.2.13.2.1 Transport Format Combination Indicator (TFCI)

The Transport Format Combination Indicator (TFCI) informs the receiver of the transport format combination of the CCTrCHs . As soon as the TFCI is detected, the transport format combination, and hence the individual transport channels' transport formats are known, and decoding of the transport channels can be performed.

### 4.3 Coding for layer 1 control

### 4.3.1 Coding of transport format combination indicator (TFCI)

The number of TFCI bits is variable and is set at the beginning of the call via higher layer signalling. Encoding of the TFCI bits depends on the number of them. If there are 6-10 bits of TFCI the channel encoding is done as described in section 4.3.1.1. Also specific coding of less than 6 bits is possible as explained in section 4.3.1.23.- For improved TFCI detection reliability repetition is used to increase the number of TFCI bits. Additionally, with any TFCI coding seheme it is assumed that in the receiver $\underline{C}$ eombining of two successive TFCI words shallwill be performed if the shortest transmission time interval of any constituent TrCH is at least 20 ms .

### 4.3.1.1 Default TFCI wordCoding of long TFCI lengths

The TFCI bits are encoded using a $(32,10)$ sub-code of the second order Reed-Muller code. The coding procedure is as shown in figure 4.3.3.1-1.


Figure 4.3.3.1-1: Channel coding of TFCI bits
TFCI is encoded by the $(32,10)$ sub-code of second order Reed-Muller code. The code words of the $(32,10)$ sub-code of second order Reed-Muller code are linear combination of some among 10 basis sequences: all 1's, 5 OVSF codes ( $\mathrm{C}_{5}(1), \mathrm{C}_{5}(2), \mathrm{C}_{5}(4), \mathrm{C}_{5}(8), \mathrm{C}_{5}(16)$ ), and 4 masks (Mask1, Mask2, Mask3, Mask4). The basis 4 mask sequences are as follows in table 4.3.1-1.

Table 4.3.1-1: BasisMask sequences for $(32,10)$ TFCI code

| 1 | $\mathrm{M}_{1,0}$ | $\mathrm{M}_{\mathrm{i}, 1}$ | $\underline{M}_{i, 2}$ | $\mathrm{M}_{\underline{\text { i, }} \text {, }}$ | $\underline{\mathrm{M}}_{1,4}$ | $\underline{M}_{1,5}$ | $\mathrm{M}_{\mathrm{i}, 6}$ | $\mathrm{M}_{\mathrm{i}, 7}$ | $\mathrm{M}_{\mathrm{i}, 8}$ | $\mathrm{M}_{\mathrm{i}, 9}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |
| 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |
| 2 | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 1 |
| 3 | 1 | $\underline{0}$ | 0 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 0 | 1 | 1 |
| 4 | 1 | 1 | - | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 1 |
| 5 | 1 | 0 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 0 | 1 | 0 |
| 6 | 1 | 1 | 1 | 1 | $\underline{0}$ | 0 | $\underline{0}$ | 1 | 0 | 0 |
| 7 | 1 | 0 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ |
| 8 | 1 | 1 | 0 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | 1 | 1 | 0 |
| 9 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | 1 |
| 10 | 1 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 1 | 1 |
| 11 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ |
| 12 | 1 | 1 | $\underline{0}$ | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ | 1 |
| 13 | 1 | $\underline{0}$ | 1 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | 1 |
| 14 | 1 | 1 | 1 | 1 | 1 | $\underline{0}$ | 1 | 1 | 1 | 1 |
| 15 | 1 | 1 | $\underline{0}$ | 0 | $\underline{0}$ | 1 | 1 | 1 | 0 | 0 |
| 16 | 1 | 0 | 1 | 0 | $\underline{0}$ | 1 | 1 | 1 | 0 | 1 |
| 17 | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ |
| 18 | 1 | $\underline{0}$ | 0 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | 1 | 1 |
| 19 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 |
| 20 | 1 | 0 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 |
| $\underline{21}$ | 1 | 1 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | 1 | 1 |
| $\underline{22}$ | 1 | $\underline{0}$ | 0 | $\underline{0}$ | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ |
| $\underline{23}$ | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | 1 | 1 | $\underline{0}$ | 1 |
| 24 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | 1 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ |
| $\underline{25}$ | 1 | 1 | 1 | $\underline{0}$ | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 |
| 26 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ |
| 27 | 1 | 1 | $\underline{0}$ | 1 | 1 | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ |
| $\underline{28}$ | 1 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |
| $\underline{29}$ | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 30 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |
| 31 | 1 | $\underline{0}$ | 0 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |


| Mask 1 | 00101000011000111111000001110114 |
| :--- | :--- |
| Mask 2 | 00000001110011010110110111000114 |
| Mask 3 | 00001010111110010001101100101014 |
| Mask 4 | 00011100001101110010111101010001 |

For TFCIinformation bits $a_{0}, a_{1}, a_{2}, a_{3}, a_{4}, a_{5}, a_{6}, a_{7}, a_{8}, a_{9}\left(a_{0}\right.$ is LSB and $a_{9}$ is MSB $)$, the encoder structure is as follows in figure 4 . 7 .output code word bits $b_{i}$ are given by:
$b_{i}=\sum_{n=0}^{9}\left(a_{n} \times M_{i, n}\right) \bmod 2$
where $\mathrm{i}=0 \ldots 31$.


Figure 4-7: Encoder structure for $(32,10)$ sub-code of second order Reed-Muller code

### 4.3.1.2 Coding of short TFCI lengths

### 4.3.1.2.1 Coding very short TFCIs by repetition

If the number of TFCI bits is 1 or 2 , then repetition will be used for coding. In this case each bit is repeated to a total of 43 times giving 4-bit transmission for a single TFCI bit and 8-bit transmission for 2 TFCI bits.

### 4.3.1.2.2 Coding short TFCIs using bi-orthogonal codes

If the number of TFCI bits is in the range of 3 to 5 , then one word of the biorthogonal $(16,5)$ block code will be used.
The code words of the biorthogenal $(16,5)$ code are from two mutually biorthogenal sets,
$S_{C_{4}}=\left\{C_{4}(0), C_{4}(1), \ldots, C_{4}(15)\right\}$ and its binary complement, $\bar{S}_{C_{4}}=\left\{\bar{C}_{4}(0), \bar{C}_{4}(1), \ldots, \bar{C}_{4}(15)\right\}$. Words of set
$S_{C_{4}}$ are from the level 4 of the code three, which is generated, using the short code generation method defined in TS 25.223. The mapping of information bits to code words is shown in the table 4.3.12.

Table-4.3.1-2: Mapping of information bits to code words for biorthogonal $(16,5)$ code

| Information bits | Code word |
| :---: | :---: |
| 00000 | $\overline{C_{4}(\theta)}$ |
| 00001 | $\overline{C_{4}(0)}$ |
| 00010 | $\overline{C_{4}(1)}$ |
| $\bar{\ldots}$ | $\overline{C_{4}(14)}$ |
| 11101 | $\overline{C_{4}(15)}$ |
| 11110 | $\overline{C_{4}(15)}$ |
| 11111 |  |

If the number of TFCI bits is in the range 3 to 5 the TFCI bits are encoded using a $(16,5)$ bi-orthogonal (or first order Reed-Muller) code. The coding procedure is as shown in figure 4-8.


## Figure 4-8: Channel coding of short length TFCI bits

The code words of the $(16,5)$ bi-orthogonal code are linear combinations of 5 basis sequences as defined in table 4.3.12 below.

## Table 4.3.1-2: Basis sequences for $(16,5)$ TFCI code

| I | $\underline{M} \mathrm{M}_{\mathrm{i}, \mathrm{O}}$ | $\mathrm{M}_{\mathrm{i}, 1}$ | $\underline{M_{i 2}}$ | $\underline{M_{i, 3}}$ | $\underline{\mathrm{M}} \mathrm{M}, 4^{\text {a }}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\underline{0}$ | 1 | 1 | $\underline{0}$ | 0 | $\bigcirc$ |
| 1 | 1 | 0 | 1 | $\underline{0}$ | 0 |
| $\underline{2}$ | 1 | 1 | 1 | $\bigcirc$ | $\underline{0}$ |
| $\underline{3}$ | 1 | 0 | $\underline{0}$ | 1 | 0 |
| 4 | 1 | 1 | $\underline{0}$ | 1 | 0 |
| $\underline{5}$ | 1 | $\underline{0}$ | 1 | 1 | 0 |
| $\underline{6}$ | 1 | 1 | 1 | 1 | $\underline{0}$ |
| 7 | 1 | 0 | $\underline{0}$ | $\underline{0}$ | 1 |
| $\underline{8}$ | 1 | 1 | $\underline{0}$ | 0 | 1 |
| $\underline{9}$ | 1 | 0 | 1 | $\underline{0}$ | 1 |
| 10 | 1 | 1 | 1 | 0 | 1 |
| 11 | 1 | 0 | $\underline{0}$ | 1 | 1 |
| $\underline{12}$ | 1 | 1 | $\underline{0}$ | 1 | 1 |
| $\underline{13}$ | 1 | 0 | 1 | 1 | 1 |
| 14 | 1 | 1 | 1 | 1 | 1 |
| 15 | 1 | 0 | $\underline{0}$ | 0 | 0 |

For TFCI information bits $\mathrm{a}_{\underline{0}}, \mathrm{a}_{\underline{1}}, \mathrm{a}_{\underline{2}}, \mathrm{a}_{\underline{3}}, \mathrm{a}_{4}\left(\mathrm{a}_{\underline{0}}\right.$ is LSB and $\left.\mathrm{a}_{\underline{4}} \underline{\text { is MSB }}\right)$, the ), the output code word bits $\mathrm{b}_{\underline{1}}$ are given by: $b_{i}=\sum_{n=0}^{4}\left(a_{n} \times M_{i, n}\right) \bmod 2$
where $\mathrm{i}=0 \ldots 15(\mathrm{i}=15$ corresponds to the MSB).

