$\uparrow C R$ number as allocated by MCC support team
For submission to: TAG-RAN\#6
list expected approval meeting \# here $\uparrow$

| for approval | $\mathbf{X}$ |
| ---: | ---: |
| for information |  |



Form: CR cover sheet, version 2 for 3GPP and SMG The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc


## Work item:

| Category: | F | Correction | X | Release: | Phase 2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | A | Corresponds to a correction in an earlier release |  |  | Release 96 |
| (only one category | B | Addition of feature |  |  | Release 97 |
| shall be marked | C | Functional modification of feature |  |  | Release 98 |
| with an $X$ ) | D | Editorial modification |  |  | Release 99 |
|  |  |  |  |  | Release 00 |

Reason for

$\underline{\text { change: }} \quad$| Description of TFCl coding for normal and short length cases is quite different and |
| :--- |
| does not specify coding/puncturing process completely. Also dependancy on OVSF |
| sequences defined externally which should be removed. |

## Clauses affected: $\quad 4.2 .13,4.3 .1,4.3 .1 .1,4.3 .1 .2$



## Other comments:

### 4.2.13 Transport format detection

Transport format detection can be performed both with and without Transport Format Combination Indicator (TFCI). If a TFCI is transmitted, the receiver detects the transport format combination from the TFCI. When no TFCI is transmitted, so called blind transport format detection may beis used, i.e. the receiver side uses the possible transport format combinations as a priori information.

### 4.2.13.1 Blind transport format detection

Blind transport format detection may be performed in the receiver by trying all possible combinations of the transport format.

### 4.2.13.2 Explicit transport format detection based on TFCI

### 4.2.13.2.1 Transport Format Combination Indicator (TFCI)

The Transport Format Combination Indicator (TFCI) informs the receiver of the transport format combination of the CCTrCHs . As soon as the TFCI is detected, the transport format combination, and hence the individual transport channels' transport formats are known, and decoding of the transport channels can be performed.

### 4.3 Coding for layer 1 control

### 4.3.1 Coding of transport format combination indicator (TFCI)

The number of TFCI bits is variable and is set at the beginning of the call via higher layer signalling. Encoding of the TFCI bits depends on the number of them. If there are 6-10 bits of TFCI the channel encoding is done as described in section 4.3.1.1. Also specific coding of less than 6 bits is possible as explained in section 4.3.1.23. For improved TFCI detection reliability repetition is used to increase the number of TFCI bits. Additionally, with any TFCI coding scheme it is assumed that in the receiver combining of two successive TFCI words will be performed if the shortest transmission time interval of any TrCH is at least 20 ms .

### 4.3.1.1 Default TFCI wordCoding of long TFCI lengths

The TFCI bits are encoded using a $(30,10)$ punctured sub-code of the second order Reed-Muller code. The coding procedure is as shown in figure 4.3.3.1-1.


Figure 4.3.3.1-1: Channel coding of TFCI bits
TFCI is encoded by the $(32,10)$ sub-code of second order Reed-Muller code. The code words of the $(32,10)$ sub-code of second order Reed-Muller code are linear combination of some among 10 basis sequences: all 1's, 5 HadamardOVSF codes $\left(\underline{H}_{5,1}, \mathrm{H}_{5,2}, \mathrm{H}_{5,4}, \mathrm{H}_{5,8}, \mathrm{H}_{5,16} \mathrm{C}_{5}(1), \mathrm{C}_{5}(2), \mathrm{C}_{5}(4), \mathrm{C}_{5}(8), \mathrm{C}_{5}(16)\right.$, and 4 masks (Mask1, Mask2, Mask3, Mask4). The 4 mask sequences are as follows in table 4.3.1-1.

Table 4.3.1-1: Mask and Hadamard sequences

|  | $\underline{\text { Mask }}$ | $\underline{\text { Mask } 2}$ | $\underline{\text { Mask3 }}$ | $\underline{\text { Mask }}$ |  | $\underline{H}_{5,1}$ | $\underline{H_{5,2}}$ | $\underline{\mathrm{H}_{5,4}}$ | $\underline{H_{5,8}}$ | $\underline{\mathrm{H}_{5,16}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\underline{\text { bit0 }}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |  | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |
| $\underline{\text { bit1 }}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |  | $\underline{1}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |


| bit2 | 1 | $\underline{0}$ | 0 | 0 | 0 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| bit3 | $\underline{0}$ | 0 | 0 | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ |
| bit4 | 1 | 0 | 1 | 1 | $\underline{0}$ | 0 | 1 | 0 | $\bigcirc$ |
| bit5 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| bit6 | $\underline{0}$ | $\underline{0}$ | 1 | 0 | 0 | 1 | 1 | $\underline{0}$ | $\underline{0}$ |
| bit7 | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 0 | 0 |
| bit8 | $\underline{0}$ | 1 | 1 | $\underline{0}$ | 0 | 0 | $\underline{0}$ | 1 | $\bigcirc$ |
| bit9 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| bit10 | 1 | $\underline{0}$ | 1 | 1 | 0 | 1 | $\underline{0}$ | 1 | $\underline{0}$ |
| bit11 | $\underline{0}$ | 0 | 1 | 1 | 1 | 1 | $\underline{0}$ | 1 | $\underline{0}$ |
| bit12 | $\underline{0}$ | 1 | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ |
| bit13 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 0 |
| bit14 | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 1 | $\bigcirc$ |
| bit15 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |
| bit16 | 1 | 0 | 0 | $\bigcirc$ | 0 | $\bigcirc$ | $\bigcirc$ | 0 | 1 |
| bit17 | 1 | 1 | 0 | $\bigcirc$ | 1 | 0 | 0 | 0 | 1 |
| bit18 | 1 | 1 | 0 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | 1 |
| bit19 | 1 | - | 1 | $\underline{0}$ | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 |
| bit20 | $\underline{0}$ | 1 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ | 1 |
| bit21 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 |
| bit22 | $\underline{0}$ | $\underline{0}$ | 1 | 1 | $\underline{0}$ | 1 | 1 | $\underline{0}$ | 1 |
| bit23 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | $\bigcirc$ | 1 |
| bit24 | $\underline{0}$ | 1 | 0 | $\underline{0}$ | 0 | 0 | $\underline{0}$ | 1 | 1 |
| bit25 | 1 | 1 | 0 | 1 | 1 | $\underline{0}$ | $\underline{0}$ | 1 | 1 |
| bit26 | 1 | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ | 1 | 1 |
| bit27 | 1 | 0 | $\bigcirc$ | 1 | 1 | 1 | $\underline{0}$ | 1 | 1 |
| bit28 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ | $\underline{0}$ | $\underline{0}$ | 1 | 1 | 1 |
| bit29 | 1 | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 1 |
| bit30 | 1 | 1 | 1 | 0 | $\bigcirc$ | 1 | 1 | 1 | 1 |
| bit31 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
|  | Mask 1 | 0010100001100011111100000111014 |  |  |  |  |  |  |  |
|  | Mask 2 | 00000001110011010110110111000111 |  |  |  |  |  |  |  |
|  | Mask 3 | 00001010111110010001101100101014 |  |  |  |  |  |  |  |
|  | Mask 4 | 00011100001101110010111101010001 |  |  |  |  |  |  |  |

For information bits $a_{0}, a_{1}, a_{2}, a_{3}, a_{4}, a_{5}, a_{6}, a_{7}, a_{8}, a_{9}\left(a_{0}\right.$ is LSB and $a_{9}$ is MSB $)$, the encoder structure is as follows in figure 4-7.


Figure 4-7: Encoder structure for $(32,10)$ sub-code of second order Reed-Muller code
Then, the output words of the $(32,10)$ sub-code of second order Reed-Muller coder are punctured into length 30 by puncturing output bits $\mathrm{c}_{0}$ and $\mathrm{c}_{16}$. The remaining punctured bits are denoted by $b_{\underline{k}}, k=0,1,2, \ldots, 29(k=29$ corresponds to the MSB bit).

### 4.3.1.2 Coding of short TFCI lengths

### 4.3.1.2.1 Coding very short TFCls by repetition

If the number of TFCI bits is 1 or 2 , then repetition will be used for coding. In this case each bit is repeated to a total of 43 times giving 4-bit transmission for a single TFCI bit and 8-bit transmission for 2 TFCI bits.

### 4.3.1.2.2 Coding short TFCIs using bi-orthogonal codes

If the number of TFCI bits is in the range of 3 to 5 , then one word of the biorthogenal $(16,5)$ block code will be used.
The code words of the biorthogonal $(16,5)$ code are from two mutually biorthogonal sets,
$S_{C_{4}}=\left\{C_{4}(0), C_{4}(1), \ldots, C_{4}(15)\right\}$ and its binary complement, $\bar{S}_{C_{4}}=\left\{\bar{C}_{4}(\theta), \bar{C}_{4}(1), \ldots, \bar{C}_{4}(15)\right\}$. Werds of set
$S_{C_{4}}$ are from the level 4 of the code three, which is generated, using the short code generation method defined in TS 25.223. The mapping of information bits to code words is shown in the table 4.3.1 2.

Table-4.3.1-2: Mapping of information bits to code words for biorthogonal $(16,5)$ code

| Information bits | Code word |
| :---: | :---: |
| 00000 | $\overline{C_{4}(\theta)}$ |
| 00001 | $\overline{C_{4}(0)}$ |
| 00010 | $\overline{C_{4}(1)}$ |
| $\ldots$ | $\overline{-}$ |
| 11101 | $\overline{C_{4}(14)}$ |
| 11110 | $\overline{C_{4}(15)}$ |
| 11111 | $\overline{C_{4}(15)}$ |

The TFCI bits are encoded using a $(15,5)$ punctured bi-orthogonal code. The coding procedure is as shown in figure 4-8.


Figure 4-8: Channel coding of short length TFCI bits
Firstly, TFCI is encoded by the $(16,5)$ bi-orthogonal (or first order Reed-Muller) code. The code words of the $(16,5)$ biorthogonal code are linear combinations of 5 basis sequences: the all 1's sequence and 4 Hadamard codes $\left(\mathrm{H}_{4,1}, \mathrm{H}_{4,2}\right.$, $\underline{H}_{4,4}, \mathrm{H}_{4,8}$ ) as defined in table 4.3.1-2 below.

Table 4.3.1-2: Hadamard sequences

|  | $\mathrm{H}_{4,1}$ | $\mathrm{H}_{4,2}$ | $\mathrm{H}_{4,4}$ | $\mathrm{H}_{4,8}$ |
| :---: | :---: | :---: | :---: | :---: |
| bit0 | 0 | 0 | $\underline{0}$ | $\underline{0}$ |
| bit1 | 1 | 0 | $\bigcirc$ | $\underline{0}$ |
| bit2 | $\underline{0}$ | 1 | 0 | $\underline{0}$ |
| bit3 | 1 | 1 | $\bigcirc$ | $\underline{0}$ |
| bit4 | $\underline{0}$ | $\underline{0}$ | 1 | $\underline{0}$ |
| bit5 | 1 | 0 | 1 | $\underline{\square}$ |
| bit6 | 0 | 1 | 1 | $\underline{0}$ |
| bit7 | 1 | 1 | 1 | $\underline{0}$ |
| bit6 | $\underline{0}$ | 0 | 0 | 1 |
| $\underline{\text { bit9 }}$ | 1 | $\underline{0}$ | $\underline{0}$ | 1 |
| bit10 | $\underline{0}$ | 1 | $\bigcirc$ | 1 |
| bit11 | 1 | 1 | $\underline{0}$ | 1 |
| bit12 | $\underline{0}$ | 0 | 1 | 1 |
| bit13 | 1 | 0 | 1 | 1 |


| $\underline{\text { bit14 }}$ | $\underline{0}$ | 1 | 1 | 1 |
| :---: | :---: | :---: | :---: | :---: |
| $\underline{\text { bit15 }}$ | 1 | 1 | 1 | 1 |

For information bits $a_{0}, a_{1}, a_{2}, a_{3}, a_{4}\left(a_{0}\right.$ is LSB and $a_{4}$ is MSB), the encoder structure is as following figure 13 (summation is modulo 2).


Figure 13: Encoder structure for $(16,5)$ bi-orthogonal code
Then the output words of the $(16,5)$ bi-orthogonal coder are punctured to length 15 by puncturing bits $\mathrm{c}_{0}$ (the lsb).
The bits in the punctured code word are denoted by $b_{k}$, where subscript $k$ indicates bit position in the code word ( $k$ $\equiv 14$ corresponds to the MSB bit).

