3GPP TSG R Dresden, Ge		leeting #9 Nov - 3 Dec 19	999		Docum	e.g. for 30	R1-99J1 GPP use the format T SMG, use the format F	P-99xxx
		CHANGE	REQI	JEST	Please see embedo page for instruction			
		25.222	CR	008	Current	t Versio	n: <u>3.0.0</u>	
GSM (AA.BB) or 3	G (AA.BBB) specific	cation number ↑		↑ CR /	number as allocated	by MCC si	upport team	
For submission	meeting # here \uparrow	for info	pproval rmation	X		strateg -strateg	gic use o	nly)
Proposed char (at least one should be	nge affects:	ersion 2 for 3GPP and SMG	ME		m is available from: ftp:// FRAN / Radio		Core Network	
Source:	Siemens					Date:	1999-11-22	
Subject:	TFCI codin	g						
Work item:								
(only one category shall be marked	B Addition of	ds to a correction f feature modification of fe		rlier releas		ease:	Phase 2 Release 96 Release 97 Release 98 Release 99 Release 00	X
<u>Reason for</u> change:	does not s	of TFCI coding for pecify coding/punct defined externally	cturing pi	ocess con	npletely. Also c			
Clauses affecte	ed: 4.2.13	8 <mark>, 4.3.1, 4.3.1.1, 4</mark> .	.3.1.2					
Other specs affected:	Other 3G co Other GSM of specifica MS test speci BSS test specific O&M specific	tions cifications ecifications	-	→ List of C → List of C	CRs: CRs: CRs:			
<u>Other</u> comments:								

4.2.13 Transport format detection

Transport format detection can be performed both with and without Transport Format Combination Indicator (TFCI). If a TFCI is transmitted, the receiver detects the transport format combination from the TFCI. When no TFCI is transmitted, so called blind transport format detection <u>may beis</u> used, i.e. the receiver side uses the possible transport format combinations as a priori information.

4.2.13.1 Blind transport format detection

Blind transport format detection may be performed in the receiver by trying all possible combinations of the transport format.

4.2.13.2 Explicit transport format detection based on TFCI

4.2.13.2.1 Transport Format Combination Indicator (TFCI)

The Transport Format Combination Indicator (TFCI) informs the receiver of the transport format combination of the CCTrCHs. As soon as the TFCI is detected, the transport format combination, and hence the individual transport channels' transport formats are known, and decoding of the transport channels can be performed.

4.3 Coding for layer 1 control

4.3.1 Coding of transport format combination indicator (TFCI)

The number of TFCI bits is variable and is set at the beginning of the call via higher layer signalling. Encoding of the TFCI bits depends on the number of them. If there are 6-10 bits of TFCI the channel encoding is done as described in section 4.3.1.1. Also specific coding of less than 6 bits is possible as explained in section 4.3.1.23. For improved TFCI detection reliability repetition is used to increase the number of TFCI bits. Additionally, with any TFCI coding scheme it is assumed that in the receiver combining of two successive TFCI words will be performed if the shortest transmission time interval of any TrCH is at least 20 ms.

4.3.1.1 Default TFCI wordCoding of long TFCI lengths

The TFCI bits are encoded using a (30, 10) punctured sub-code of the second order Reed-Muller code. The coding procedure is as shown in figure 4.3.3.1-1.

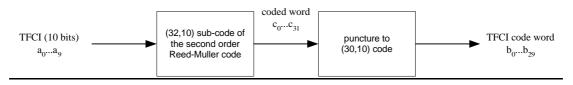


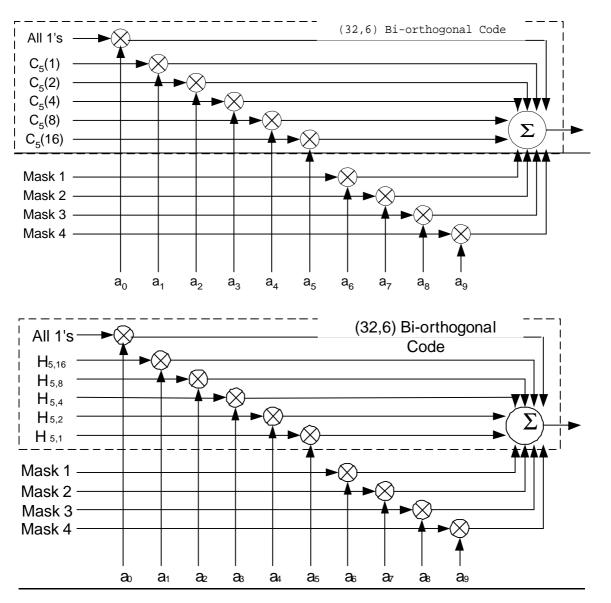
Figure 4.3.3.1-1: Channel coding of TFCI bits

TFCI is encoded by the (32,10) sub-code of second order Reed-Muller code. The code words of the (32,10) sub-code of second order Reed-Muller code are linear combination of some among 10 basis sequences: all 1's, 5 <u>HadamardOVSF</u> codes ($\underline{H}_{5,1}, \underline{H}_{5,2}, \underline{H}_{5,4}, \underline{H}_{5,8}, \underline{H}_{5,16}, \underline{C}_{s}(1), \underline{C}_{s}(2), \underline{C}_{s}(4), \underline{C}_{s}(8), \underline{C}_{s}(16)$), and 4 masks (Mask1, Mask2, Mask3, Mask4). The 4 mask-sequences are as follows in table 4.3.1-1.

	Mask1	Mask2	Mask3	Mask4	<u>H_{5,1}</u>	<u>H_{5,2}</u>	<u>H_{5,4}</u>	<u>H_{5,8}</u>	<u>H_{5.16}</u>
<u>bit0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>
<u>bit1</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>

<u>bit2</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>		<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>
<u>bit3</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>		<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>
bit4	<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>		<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>
bit5	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>		<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>
<u>bit6</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>		<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>
<u>bit7</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>		<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>
<u>bit8</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>		<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>
bit9	<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>		<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>
<u>bit10</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>		<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>
<u>bit11</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>		<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>
<u>bit12</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>		<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>
<u>bit13</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>		<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>
bit14	<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>		<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>
<u>bit15</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>		<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>
<u>bit16</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>		<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>
<u>bit17</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>		<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>
<u>bit18</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>		<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>
<u>bit19</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>		<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>
<u>bit20</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>		<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>
<u>bit21</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>		<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>
<u>bit22</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>		<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>
<u>bit23</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>		<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>
<u>bit24</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>		<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>
<u>bit25</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>		<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>
<u>bit26</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>		<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>
<u>bit27</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>		<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>
<u>bit28</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>		<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>
<u>bit29</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>		<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>
<u>bit30</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>		<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>
<u>bit31</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>		<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>
1	Mask 1	001	0100001	1000111	1110000	011101	11	1		
ŀ	Mask 2									
F					01101101					
ŀ	Mask 3				0011011					
L	Mask 4	000	1110000	1101110	0101111	0101000)1			

For information bits a_0 , a_1 , a_2 , a_3 , a_4 , a_5 , a_6 , a_7 , a_8 , a_9 (a_0 is LSB and a_9 is MSB), the encoder structure is as follows in figure 4-7.



4

Figure 4-7: Encoder structure for (32,10) sub-code of second order Reed-Muller code

Then, the output words of the (32,10) sub-code of second order Reed-Muller coder are punctured into length 30 by puncturing output bits c_0 and c_{16} . The remaining punctured bits are denoted by b_k , k = 0, 1, 2, ..., 29 (k = 29 corresponds to the MSB bit).

4.3.1.2 Coding of short TFCI lengths

4.3.1.2.1 Coding very short TFCIs by repetition

If the number of TFCI bits is 1 or 2, then repetition will be used for coding. In this case each bit is repeated to a total of 43 times giving 4-bit transmission for a single TFCI bit and 8-bit transmission for 2 TFCI bits.

4.3.1.2.2 Coding short TFCIs using bi-orthogonal codes

If the number of TFCI bits is in the range of 3 to 5, then one word of the biorthogonal (16,5) block code will be used.

The code words of the biorthogonal (16, 5) code are from two mutually biorthogonal sets,

 $S_{C_4} = \{C_4(0), C_4(1), \dots, C_4(15)\} \text{ and its binary complement, } \overline{S}_{C_4} = \{\overline{C}_4(0), \overline{C}_4(1), \dots, \overline{C}_4(15)\}. \text{ Words of set}$

3GPP

 S_{C_4} are from the level 4 of the code three, which is generated, using the short code generation method defined in TS 25.223. The mapping of information bits to code words is shown in the table 4.3.1 -2.

Table 4.3.1-2: Mapping of information bits to code words for biorthogonal (16, 5) code

	r
Information bits	Code word
00000	$-C_4(0)$
00001	$\overline{C_4(0)}$
00010	$-C_{4}(1)$
	
11101	$\overline{C_4(14)}$
11110	$-C_4(15)$
11111	$\overline{C_4(15)}$

The TFCI bits are encoded using a (15, 5) punctured bi-orthogonal code. The coding procedure is as shown in figure <u>4-8</u>.

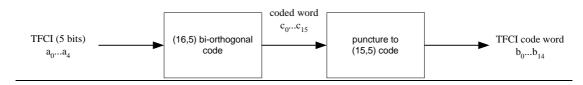


Figure 4-8: Channel coding of short length TFCI bits

<u>Firstly, TFCI is encoded by the (16,5) bi-orthogonal (or first order Reed-Muller) code. The code words of the (16,5) bi-orthogonal code are linear combinations of 5 basis sequences: the all 1's sequence and 4 Hadamard codes ($H_{4,1}$, $H_{4,2}$, $H_{4,4}$, $H_{4,4}$, $H_{4,8}$) as defined in table 4.3.1-2 below.</u>

	<u>H_{4,1}</u>	<u>H_{4,2}</u>	<u>H_{4,4}</u>	<u>H_{4,8}</u>
<u>bit0</u>	<u>0</u>	<u>0</u>	<u>0</u>	<u>0</u>
<u>bit1</u>	<u>1</u>	<u>0</u>	<u>0</u>	<u>0</u>
bit2	<u>0</u>	<u>1</u>	<u>0</u>	<u>0</u>
bit3	<u>1</u>	<u>1</u>	<u>0</u>	<u>0</u>
bit4	<u>0</u>	<u>0</u>	<u>1</u>	<u>0</u>
bit5	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>
bit6	<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>
bit7	<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>
bit8	<u>0</u>	<u>0</u>	<u>0</u>	<u>1</u>
bit9	<u>1</u>	<u>0</u>	<u>0</u>	<u>1</u>
<u>bit10</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>
<u>bit11</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>
<u>bit12</u>	<u>0</u>	<u>0</u>	<u>1</u>	<u>1</u>
<u>bit13</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>

Table 4.3.1-2: Hadamard sequences

<u>bit14</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>
<u>bit15</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>

6

For information bits a_0 , a_1 , a_2 , a_3 , a_4 (a_0 is LSB and a_4 is MSB), the encoder structure is as following figure 13 (summation is modulo 2).

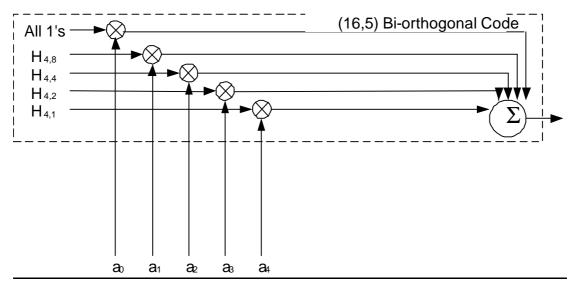


Figure 13: Encoder structure for (16,5) bi-orthogonal code

Then the output words of the (16,5) bi-orthogonal coder are punctured to length 15 by puncturing bits c_0 (the lsb).

The bits in the punctured code word are denoted by b_k , where subscript k indicates bit position in the code word (k =14 corresponds to the MSB bit).