

CHANGE REQUEST

⌘ **26.346 CR 013** ⌘ rev **1** ⌘ Current version: **6.0.0** ⌘

For **HELP** on using this form, see bottom of this page or look at the pop-up text over the ⌘ symbols.

Proposed change affects: UICC apps⌘ ME Radio Access Network Core Network

| | | | |
|------------------------|---|--|--|
| Title: | ⌘ | Specification of Raptor Forward Error Correction and Streaming User Service bundling (Combination of CR 6 rev1, CR 7 rev 1 and CR 12) | |
| Source: | ⌘ | TSG SA WG4 Codec | |
| Work item code: | ⌘ | MBMS-TSMBMS | Date: ⌘ 06/06/2005 |
| Category: | ⌘ | C | Release: ⌘ Rel-6 |
| | | <i>Use <u>one</u> of the following categories:</i> F (correction) A (corresponds to a correction in an earlier release) B (addition of feature), C (functional modification of feature) D (editorial modification) Detailed explanations of the above categories can be found in 3GPP TR 21.900 . | <i>Use <u>one</u> of the following releases:</i> Ph2 (GSM Phase 2) R96 (Release 1996) R97 (Release 1997) R98 (Release 1998) R99 (Release 1999) Rel-4 (Release 4) Rel-5 (Release 5) Rel-6 (Release 6) Rel-7 (Release 7) |

| | | |
|--------------------------------------|---|--|
| Reason for change: | ⌘ | Introduction of Raptor Forward Error Correction for MBMS file download and streaming services and Specification of Streaming User Service Bundling |
| Summary of change: | ⌘ | MBMS FEC scheme definition added to main body of TS Raptor specification included as Annex B. The User Service description XML is updated to allow for bundling of streaming user services. The service protection description is updated to indicate if the key management stream is FEC protected. A new FEC protection stream description is included. The Session Description is updated to describe the presence of the FEC in combination with the source stream and its parameter. The streaming FEC framework is changed to support the bundling of streams. This includes changes to source, repair packets, and the source block format. Needed SDP modifications and additions are introduced. The IANA registration information is also changed to specify SDP protocol identifiers instead of media types. |
| Consequences if not approved: | ⌘ | No Forward Error Correction code included for MBMS. No streaming user service bundling would be supported and the substantial gain in media bit-rate that this produces would not be available. |

| Clauses affected: | ⌘ | 2, 4.4, 5.2.1, 5.2.2, 5.5, 7.2.2, 7.2.3, 7.2.7, 7.2.9, 7.2.10, 7.2.11, 7.3.2, 8.2.2, 8.3.1, 8.3.2, Annex A, Annex B, Annex C. | | | | | | | | | | |
|------------------------------|---------------------|---|---|---|--|---|--|---|--|---|---------------------------|---|
| Other specs affected: | ⌘ | <table border="1"><tr><th>Y</th><th>N</th></tr><tr><td></td><td>X</td></tr><tr><td></td><td>X</td></tr><tr><td></td><td>X</td></tr></table> | Y | N | | X | | X | | X | Other core specifications | ⌘ |
| | | Y | N | | | | | | | | | |
| | | | X | | | | | | | | | |
| | X | | | | | | | | | | | |
| | X | | | | | | | | | | | |
| | Test specifications | | | | | | | | | | | |
| | O&M Specifications | | | | | | | | | | | |
| Other comments: | ⌘ | This CR is a purely editorial combination of CR 6 rev1, CR 7 rev 1 and CR 12, since these three CRs address substantially overlapping sections of the TS. Additionally some erroneous section numbering and text styles in those CRs are corrected. | | | | | | | | | | |

How to create CRs using this form:

Comprehensive information and tips about how to create CRs can be found at <http://www.3gpp.org/specs/CR.htm>. Below is a brief summary:

- 1) Fill out the above form. The symbols above marked ⌘ contain pop-up help information about the field that they are closest to.
- 2) Obtain the latest version for the release of the specification to which the change is proposed. Use the MS Word "revision marks" feature (also known as "track changes") when making the changes. All 3GPP specifications can be downloaded from the 3GPP server under <ftp://ftp.3gpp.org/specs/> For the latest version, look for the directory name with the latest date e.g. 2001-03 contains the specifications resulting from the March 2001 TSG meetings.
- 3) With "track changes" disabled, paste the entire CR form (use CTRL-A to select it) into the specification just in front of the clause containing the first piece of changed text. Delete those parts of the specification which are not relevant to the change request.

***** Start of Changes *****

2 References

The following documents contain provisions which, through reference in this text, constitute provisions of the present document.

- References are either specific (identified by date of publication, edition number, version number, etc.) or non-specific.
- For a specific reference, subsequent revisions do not apply.
- For a non-specific reference, the latest version applies. In the case of a reference to a 3GPP document (including a GSM document), a non-specific reference implicitly refers to the latest version of that document *in the same Release as the present document*.

- [1] 3GPP TR 21.905: "Vocabulary for 3GPP Specifications".
- [2] 3GPP TS 22.146: "Multimedia Broadcast/Multicast Service; Stage 1".
- [3] 3GPP TS 22.246: "Multimedia Broadcast/Multicast Service (MBMS) user services; Stage 1".
- [4] 3GPP TS 23.246: "Multimedia Broadcast/Multicast Service (MBMS); Architecture and functional description".
- [5] 3GPP TS 25.346: "Introduction of Multimedia Broadcast/Multicast Service (MBMS) in the Radio Access Network (RAN); Stage 2".
- [6] IETF RFC 3550 (July 2003): "RTP: A Transport Protocol for Real-Time Applications", H. Schulzrinne, S. Casner, R. Frederick, V. Jacobson.
- [7] IETF STD 0006/RFC 0768 (August 1980): "User Datagram Protocol", J. Postel.
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- [9] IETF RFC 3926 (October 2004): "FLUTE - File Delivery over Unidirectional Transport", T. Paila, M. Luby, R. Lehtonen, V. Roca, R. Walsh.
- [10] IETF RFC 3450 (December 2002): "Asynchronous Layered Coding (ALC) Protocol Instantiation", M. Luby, J. Gemmell, L. Vicisano, L. Rizzo, J. Crowcroft.
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- [14] IETF RFC 2327 (April 1998): "SDP: Session Description Protocol", M. Handley and V. Jacobson.
- [15] IETF draft-ietf-mmusic-sdp-srcfilter-06: "Session Description Protocol (SDP) Source Filters".
- [16] IETF RFC 3266 (June 2002): "Support for IPv6 in Session Description Protocol (SDP)", S. Olson, G. Camarillo, A. B. Roach.
- [17] IETF RFC 3048 (January 2001): "Reliable Multicast Transport Building Blocks for One-to-Many Bulk-Data Transfer", B. Whetten, L. Vicisano, R. Kermode, M. Handley, S. Floyd, M. Luby.
- [18] IETF RFC 2616 (June 1999): "Hypertext Transfer Protocol -- HTTP/1.1".
- [19] IETF RFC 1738 (December 1994): "Uniform Resource Locators (URL)".

- [20] 3GPP TS 33.246: "3G Security; Security of Multimedia Broadcast/Multicast Service (MBMS)".
- [21] OMG: "Unified Modeling Language (UML), version 1.5" (formal/03-03-01).
- [22] W3C Recommendation 28 October 2004: "XML Schema Part 2: Datatypes Second Edition".
- [23] IETF RFC 2234 (November 1997): "Augmented BNF for Syntax Specifications: ABNF", D. Crocker and P. Overell.
- [24] 3GPP TS 26.290: "Audio codec processing functions; Extended Adaptive Multi-Rate - Wideband (AMR-WB+) codec; Transcoding functions".
- [25] 3GPP TS 26.304: "Floating-point ANSI-C code for the Extended Adaptive Multi-Rate - Wideband (AMR-WB+) codec".
- [26] 3GPP TS 26.273: "Speech codec speech processing functions; Extended Adaptive Multi-Rate - Wideband (AMR-WB+) speech codec; Fixed-point ANSI-C code".
- [27] Void.
- [28] 3GPP TS 26.401: "General audio codec audio processing functions; Enhanced aacPlus general audio codec; General description".
- [29] 3GPP TS 26.410: "General audio codec audio processing functions; Enhanced aacPlus general audio codec; Floating-point ANSI-C code".
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- [32] 3GPP TS 26.244: "Transparent end-to-end streaming service; 3GPP file format (3GP)".
- [33] IETF RFC 3267 (June 2002): "Real-Time Transport Protocol (RTP) Payload Format and File Storage Format for the Adaptive Multi-Rate (AMR) Adaptive Multi-Rate Wideband (AMR-WB) Audio Codecs", J. Sjoberg, M. Westerlund, A. Lakaniemi, Q. Xie.
- [34] IETF draft-ietf-avt-rtp-amrwbplus-06 (September 2004): "RTP Payload Format for Extended AMR Wideband (AMR-WB+) Audio Codec", J. Sjoberg, M. Westerlund and A. Lakaniemi.
- [35] IETF RFC 3984 (February 2005): "RTP payload Format for H.264 Video", S. Wenger, M.M. Hannuksela, T. Stockhammer, M. Westerlund, D. Singer.
- [36] Void.
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- [39] IETF RFC 3556 (July 2003): "Session Description Protocol (SDP) Bandwidth Modifiers for RTP Control Protocol (RTCP) Bandwidth", S. Casner.
- [40] 3GPP TS 24.008: "Mobile radio interface Layer 3 specification; Core network protocols; Stage 3".
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- [44] ISO/IEC 14496-10/FDAM1: "AVC Fidelity Range Extensions".

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- [46] ITU-T Recommendation H.263 - Annex X (04/01): "Annex X: Profiles and levels definition".
- [47] 3GPP TS 26.234: "Transparent end-to-end streaming service; Protocols and codecs".
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- [50] 3GPP TS 26.073: "AMR speech Codec; C-source code".
- [51] 3GPP TS 26.104: "ANSI-C code for the floating-point Adaptive Multi-Rate (AMR) speech codec".
- [52] 3GPP TS 26.171: "AMR speech codec, wideband; General description".
- [53] 3GPP TS 26.190: "Mandatory Speech Codec speech processing functions AMR Wideband speech codec; Transcoding functions".
- [54] 3GPP TS 26.173: "ANSI-C code for the Adaptive Multi Rate - Wideband (AMR-WB) speech codec".
- [55] 3GPP TS 26.204: "ANSI-C code for the floating-point Adaptive Multi-Rate Wideband (AMR-WB) speech codec".
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***** Next Change *****

4.4 Functional Entities to support MBMS User Services

Figure 3 depicts the MBMS network architecture showing MBMS related entities involved in providing MBMS user services.

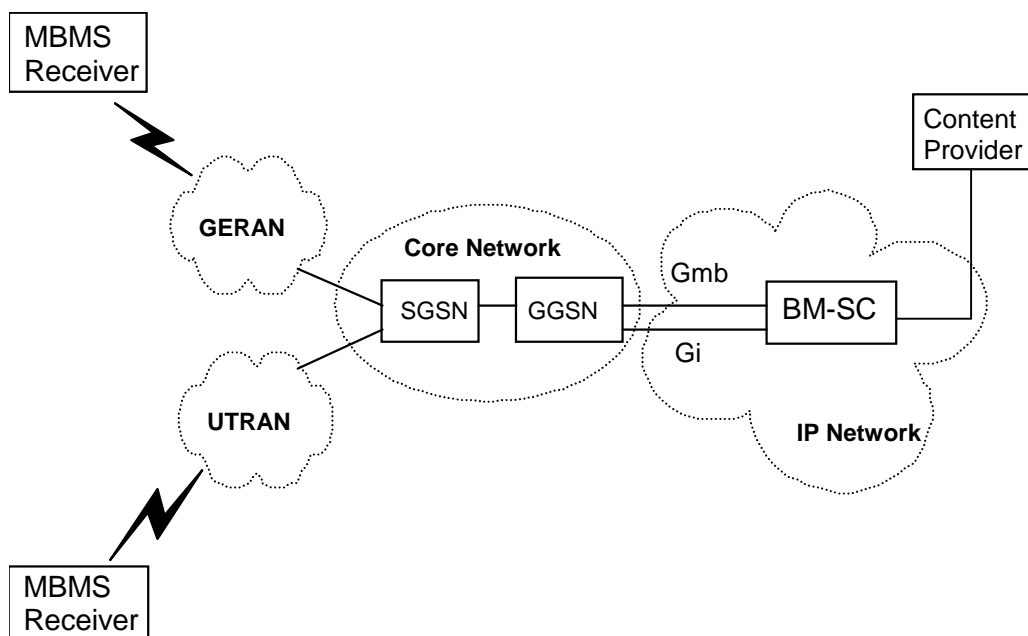


Figure 3: MBMS network architecture model

MBMS User Service architecture is based on an MBMS receiver on the UE side and a BM-SC on the network side.

The use of the Gmb and Gi interface in providing IP multicast traffic and managing MBMS bearer sessions is described in detailed in TS 23.246 [4].

Details about the BM-SC functional entities are given in figure 4.

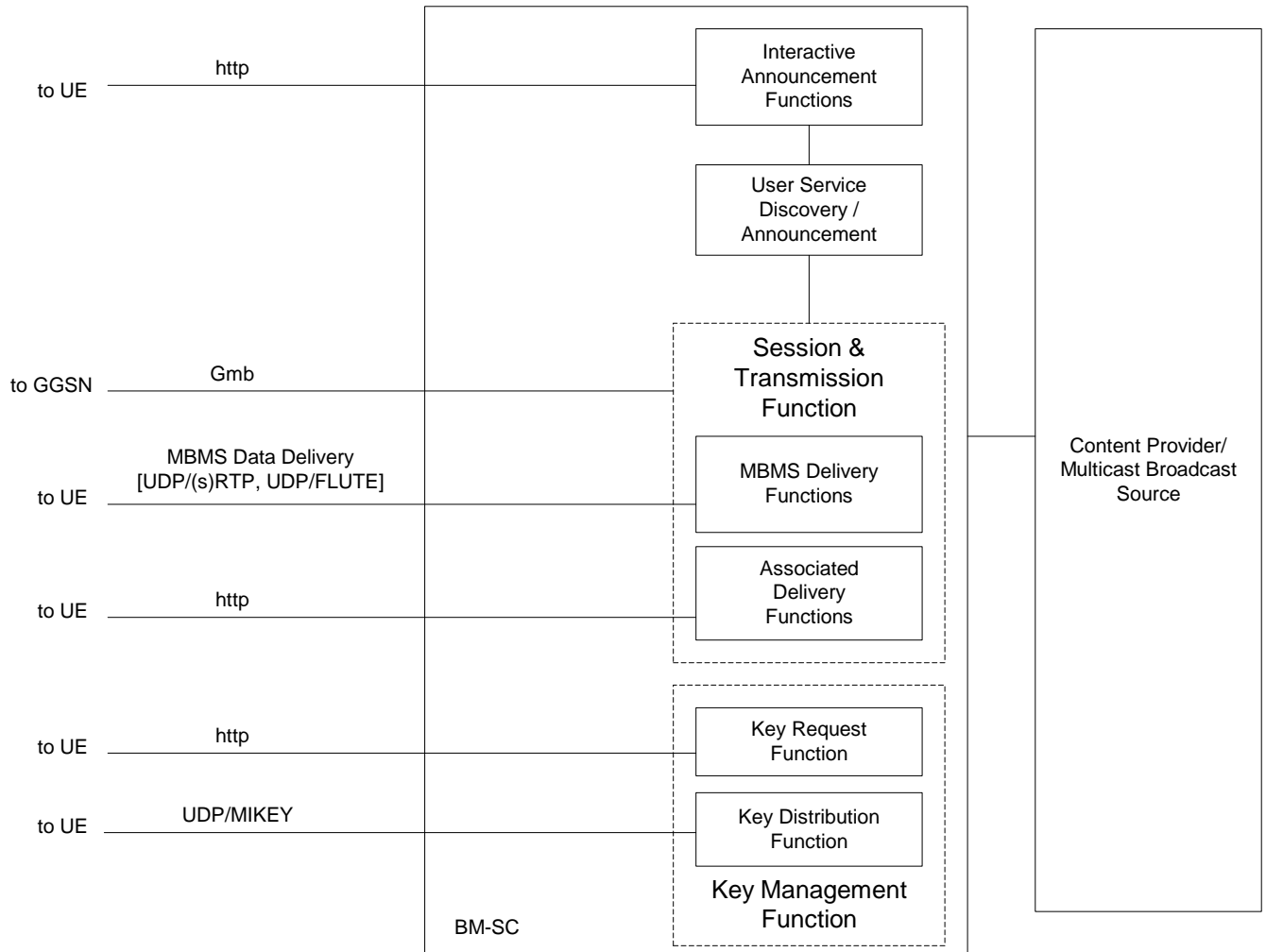


Figure 4: BM-SC sub-functional structure

The Session and Transmission function is further subdivided into the MBMS Delivery functions and the Associated Delivery functions.

The BM-SC and UE may exchange service and content related information either over point-to-point bearers or MBMS bearers whichever is suitable. To that end the following MBMS procedures are provided:

- User Service Discovery / Announcement providing service description material to be presented to the end-user as well as application parameters used in providing service content to the end-user
- MBMS-based delivery of data/content (~~optionally confidentiality and/or integrity protected~~) from the BM-SC to the UE over IP multicast.
 - o [The data/content is optionally confidentiality and/or integrity protected](#)
 - o [The data/content is optionally protected by an forward error correction code](#)
- Key Request and Registration procedure for receiving keys and key updates.

- Key distribution procedures whereby the BM-SC distributes key material required to access service data and delivered content.
- Associated Delivery functions are invoked by the UE in relation to the MBMS data transmission. The following associated delivery functions are available:
 - o File repair for download delivery method used to complement missing data.
 - o Delivery verification and reception statistics collection procedures

The interfaces between internal BM-SC functions are outside the scope of this specification.

A “Proxy and Transport function” may be located between the “Session and Transmission Function” and the GGSN. The “Proxy and Transport function” is transparent to the “Session and Transmission function”.

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4.4.3 MBMS Session and Transmission Function

The MBMS Session and Transmission function transfers the actual MBMS session data to the group of MBMS UEs. The MBMS Session and Transmission function interacts with the GGSN through the Gmb Proxy function to activate and release the MBMS transmission resources.

The function contains the MBMS delivery methods, which use the MBMS bearer service for distribution of content. Further this function contains a set of Associated-Delivery Functions, which may be invoked by the UE in relation to the MBMS data transmission (e.g. after the MBMS data transmission).

The BM-SC Session and Transmission function is further described in later clauses of this specification as well as in TS 23.246 [4].

MBMS user services data may be integrity and/or confidentiality protected as specified within TS 33.246 [20], and protection is applied between the BM-SC and the UE. This data protection is based on symmetric keys, which are shared between the BM-SC and the UEs accessing the service.

[MBMS user services may also be protected against packet loss between BM-SC and UE using a forward error correction code.](#)

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5.2.1 Introduction

User service discovery refers to methods for the UE to obtain a list of available MBMS user services [or user service bundles](#) along with information on the user services. Part of the information may be presented to the user to enable service selection.

User service announcement refers to methods for the MBMS service provider to announce the list of available MBMS user services [and user service bundles](#), along with information on the user services, to the UE.

In order for the user to be able to initiate a particular service, the UE needs certain metadata information. The required metadata information is described in clause 5.2.2.

According to 3GPP TS 23.246 [4], in order for this information to be available to the UE operators/service providers may consider several service discovery mechanisms. User service announcement may be performed over a MBMS bearer or via other means. The download delivery method is used for the user service announcement over a MBMS bearer. The user service announcement mechanism based on the download delivery method is described in clause 5.2.3. Other user service announcement and discovery mechanisms by other means than the download delivery method are out of scope of the present document.

5.2.2 MBMS User Service Description metadata fragments

MBMS User Service Discovery/Announcement is needed in order to advertise MBMS Streaming and MBMS Download User Services [and User Service Bundles](#) in advance of, and potentially during, the User Service sessions described. The User Services are described by metadata (objects/files) delivered using the download delivery method as defined in clause 7 or using interactive announcement functions.

MBMS User Service Discovery/Announcement involves the delivery of fragments of metadata to many receivers in a suitable manner. The metadata itself describes details of services. A *metadata fragment* is a single uniquely identifiable block of metadata. An obvious example of a metadata fragment would be a single SDP file (RFC 2327 [14]).

The metadata consists of:

- a metadata fragment object describing details of [a single or a bundle of](#) MBMS user services;
- a metadata fragment object(s) describing details of MBMS user service sessions;
- a metadata fragment object(s) describing details of Associated delivery methods;
- a metadata fragment object(s) describing details of service protection;
- [a metadata fragment object describing details of the FEC repair data stream.](#)

Metadata management information consists of:

- ~~—~~a metadata envelope object(s) allowing the identification, versioning, update and temporal validity of a metadata fragment.

The metadata envelope and metadata fragment objects are transported as file objects in the same download session either as separate referencing files or as a single embedding file - see clause 5.2.3.6). A single metadata envelope shall describe a single metadata fragment, and thus instances of the two are paired. An service announcement sender shall make a metadata envelope instance available for each metadata fragment instance. The creation and use of both an embedded envelope instance and a referenced envelope instance for a particular fragment instance is not recommended.

The metadata envelope and metadata fragment objects may be compressed using the generic GZip algorithm RFC 1952 [42] as content/transport encoding for transmission. Where used over an MBMS bearer, this shall be according to Download delivery content encoding using FLUTE - see clause 7.2.5.

NOTE 1: It was agreed in principle at SA4#34 that MBMS user service description allows the association of delivery methods to one or more access systems. The specification text is FFS

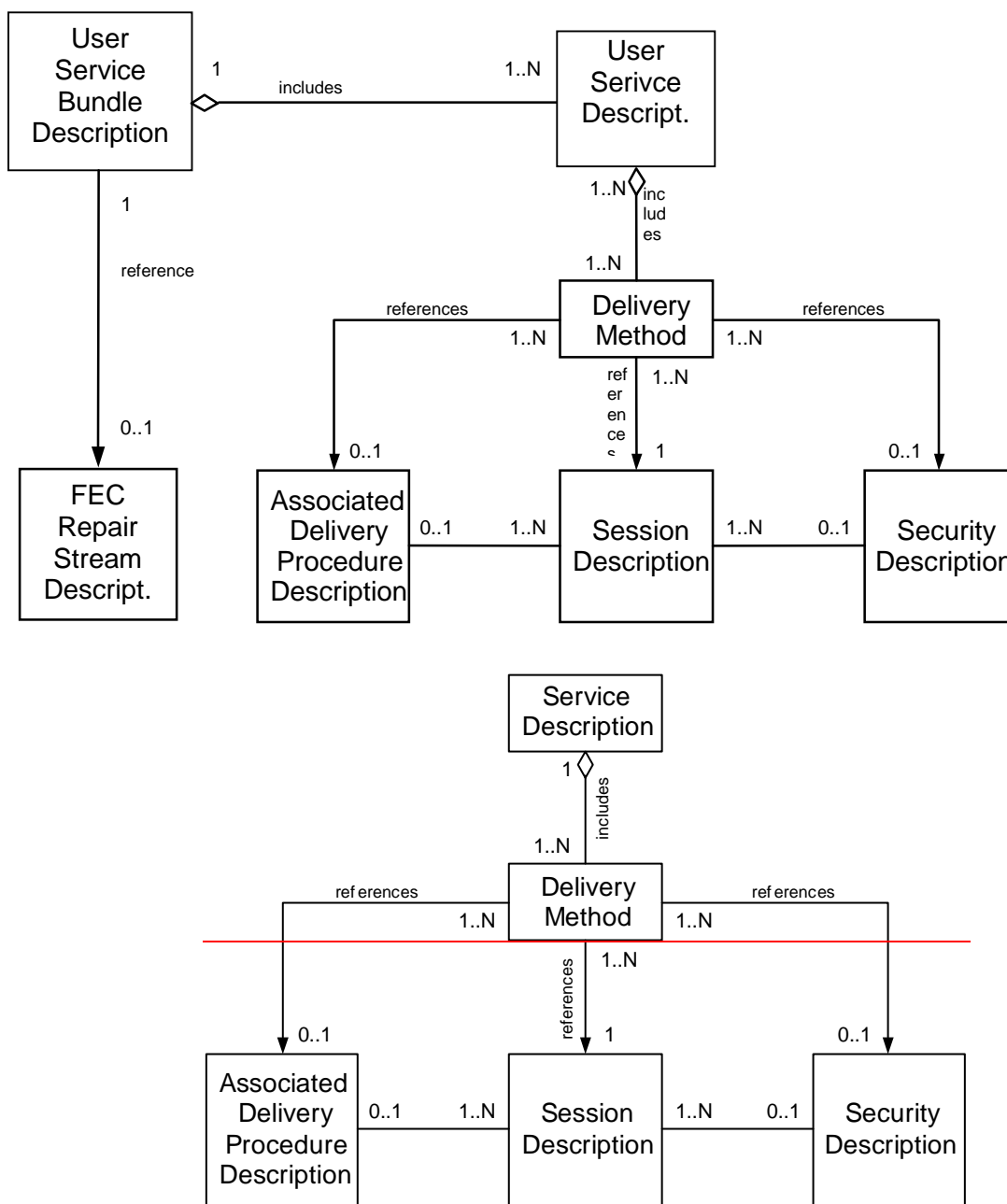


Figure 5: Simple Description Data Model

Figure 5 illustrates the simple data model relation between these description instances using UML [21] for a single User Service [Bundle](#) Description.

NOTE 2: "N" means any number in each instance.

[One MBMS User Service Bundle Description shall contain at least one User Service Description instances and may contain several. The User Service Bundle Description may refer to a single FEC Repair Stream Description.](#)

One MBMS User Service Description instance shall include at least one delivery method description instance. The delivery method description shall refer to one session description instance.

The delivery method description may contain references to a service protection description and an associated delivery procedure description. Several delivery methods may reference the same service protection description, in case the same encryption keys are used across delivery methods.

If the associated delivery procedure description is present in the user service description instance, it may be referenced by one or more delivery methods.

If the service protection description is present in the user service description instance, it may be referenced by one or more delivery methods.

Multipart MIME may be used to concatenate the descriptions one file for transport.

5.2.2.1 Session Description

One or more session descriptions are contained in one session description object. The session description instance shall be formatted according to the Session Description Protocol (SDP) [14]. Each session description instance must describe either one Streaming session or one FLUTE Download session. A session description for a Streaming session may include multiple media descriptions for RTP sessions. The *sessionDescriptionURI* references the session description object. The session description is specified in clause 7.3 for the MBMS download delivery method and in clause 8.3 for the MBMS streaming delivery method.

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5.2.2.3 Service Protection Description

The security description fragment contains the key identifiers and procedure descriptions for one delivery method. When different delivery methods use the same security description, the same security description document is referenced from the different delivery method elements.

The security description is reference by the protectionDescriptionURI of the deliveryMethod element. The security description fragment shall use the MIME type application/mbms-protection-description.

The security description contains key identifiers and the server address to request the actual key material. The key management servers are protected against overload situations like the associated delivery procedures. Associated delivery procedures are defined in clause 9.4.

The root element of the security description is the securityDescription element. It contains the key identities, which are required for one delivery method. Further the security description contains one or more key management server addresses (i.e. BM-SC).

The keyManagement element defines the list of key management servers (i.e. BM-SC). The MBMS UE must register with the key management server to receive key material.

The key management server is protected like the associated delivery procedures against overload conditions. The key management server shall be selected as defined in clause 9.3.3. The back-off time shall be determined as defined in clause 9.3.2.

The attribute confidentialityProtection defines whether a confidentiality protection scheme is used.

The attribute integrityProtection defines whether an integrity protection scheme is used.

The attribute uiccKeyManagement defines the UICC key management in the MBMS.

The element keyId contain the key identifications and the mapping to RTP media flows or FLUTE channels sessions. The identity element identifies the key as defined in clause 6.3.2.1 of 3GPP TS 33.246 [20]. The mediaFlow attribute specifies the RTP media flow or FLUTE channel. The value shall be of form <IP-destination-address>:<destination-port>. The mediaFlow element shall be present when more than one RTP media flow or FLUTE channel is defined in one session description element as defined in clause 5.2.2.1. When only one RTP media flow or one FLUTE channel is defined in the session description, then the mediaFlow attribute may not be present. The delivery method element defines the mapping between the RTP media flow or the FLUTE channel and the key identification.

The presence of the *fecProtection* element indicates that any MIKEY packet with an multicast destination IP address equal to any of the used destination address in the User Service Bundle Description instance's delivery methods, are FEC protected and encapsulated in FEC source packets, see section 8.2.2.3. The attributes *fecEncodingId*, *fecInstanceID*, and *fecOtiExtension* specifies the FEC payload ID used in the source packet. All service protection descriptions referenced by a User Service Bundle Description instance shall use the same FEC parameters.

XML schema for Security Description:

```

<?xml version="1.0" encoding="UTF-8"?>
<xs:schema elementFormDefault="qualified"
  targetNamespace="urn:3gpp:metadata:2004:securitydescription"
  xmlns="urn:3gpp:metadata:2004:securitydescription"
  xmlns:xs="http://www.w3.org/2001/XMLSchema">
  <xs:element name="securityDescription">
    <xs:element name="keyManagement" type="keyManagementType" minOccurs="0" maxOccurs="1"/>
    <xs:sequence>
      <xs:element name="keyId" type="keyIdType" minOccurs="1" maxOccurs="unbounded"/>
    </xs:sequence>
    <xs:attribute name="confidentialityProtection"
      type="xs:boolean" use="optional" default="true"/>
    <xs:attribute name="integrityProtection"
      type="xs:boolean" use="optional" default="true"/>
    <xs:attribute name="uiccKeyManagement"
      type="xs:boolean" use="optional" default="true"/>
    <xs:element name="fecProtection" type="fecProtectionType" minOccurs="0" maxOccurs="1"/>
  </xs:element>

  <xs:complexType name="keyManagementType">
    <xs:sequence>
      <xs:element name="serverURI" type="xs:anyURI" minOccurs="1" maxOccurs="unbounded"/>
    </xs:sequence>
    <xs:attribute name="waitTime" type="xs:unsignedLong" use="optional" default="0"/>
    <xs:attribute name="maxBackOff" type="xs:unsignedLong" use="optional" default="0"/>
  </xs:complexType>
  <xs:complexType name="keyIdType">
    <xs:attribute name="identity" type="xs:string" use="required"/>
    <xs:attribute name="mediaFlow" type="xs:string" use="optional"/>
  </xs:complexType>
  <xs:complexType name="fecProtectionType">
    <xs:attribute name="fecEncodingId" type="xs:unsignedLong" use="required" default="0"/>
    <xs:attribute name="fecInstanceId" type="xs:unsignedLong" use="optional"/>
    <xs:attribute name="fecOtiExtension" type="xs:string" use="optional"/>
  </xs:complexType>
</xs:schema>

```

Example of a security description:

```

<?xml version="1.0" encoding="UTF-8"?>
<securityDescription
  xmlns="www.example.com/3gppSecurityDescription"
  xmlns:xs="http://www.w3.org/2001/XMLSchema-instance"
  confidentialityProtection="true"
  integrityProtection="true"
  uiccKeyManagement="true">
  <keyManagement
    waitTime="5"
    maxBackOff="10">
    <serverURI = "http://register.operator.umts/" -/>
    <serverURI = "http:// register2.operator.umts/" -/>
  </keyManagement>
  <keyId identity="<someMSKIdA>" mediaFlow=224.1.2.3:4002 />
  <keyId identity="<someMSKIdB>" mediaFlow=224.1.2.3:4004 />
  <fecProtection
    fecEncodingId="130"
    fecInstanceId="0"
    fecOtiExtension="1SCxWEMNe397m24SwgyRhg==" />
</securityDescription>

```

5.2.2.4 XML-Schema for MBMS User Service [Bundle](#) Description

The root element of the MBMS [User Service Bundle](#) description is the [bundleDescription](#) element. The element is of the [bundleDescriptionType](#). The [bundleDescription](#) contains one or several [userServiceDescription](#) elements and optionally a reference to the FEC repair stream description. The element is of type [userServiceDescriptionType](#).

Each [userServiceDescription](#) element shall have a unique identifier. The unique identifier shall be offered as [serviceId](#) attribute within the [userServiceDescription](#) element and shall be of URN format.

The [userServiceDescription](#) element may contain one or more [name](#) elements. The intention of a [Name](#) element is to offer a title of the user service. For each name elements, the language shall be specified according to XML datatypes (XML Schema Part 2 [22]).

The *userServiceDescription* element may contain one or more *ServiceLanguage* elements. Each *serviceLanguage* element represents the available languages of the user services. The language shall be specified according to XML datatypes (XML Schema Part 2 [22]).

Each *userServiceDescription* element shall contain at least one *deliveryMethod* element. The *deliveryMethod* element contains the description of one delivery method. The element shall contain one reference to a session description and may contain references to one associated delivery procedure and/or one service protection descriptions. The session description is further specified in clause 5.2.2.1.

The *deliveryMethod* element may contain a reference to an associated delivery procedure description. The description and configuration of associated delivery procedures is specified in clause 5.2.2.5.

The *deliveryMethod* element may contain a reference to a service protection description. The service protection description is specified in clause 5.2.2.3.

```
<?xml version="1.0" encoding="UTF-8"?>
<xs:schema elementFormDefault="qualified"
  targetNamespace="urn:3gpp:metadata:2004:userservicedescription"
  xmlns="urn:3gpp:metadata:2004:userservicedescription"
  xmlns:xs="http://www.w3.org/2001/XMLSchema">
  <xs:element name="bundleDescription" type="bundleDescriptionType"/>
  <xs:complexType name="bundleDescriptionType">
    <xs:sequence>
      <xs:element name="userServiceDescription" type="userServiceDescriptionType"
        minOccurs="1" maxOccurs="unbound"/>
    </xs:sequence>
    <xs:attribute name="fecDescriptionURI" type="xs:anyURI" use="optional"/>
  </xs:complexType>

  <xs:complexType name="userServiceDescriptionType">
    <xs:sequence>
      <xs:element name="name" type="nameType" minOccurs="0"
        maxOccurs="unbounded"/>
      <xs:element name="serviceLanguage" type="xs:language" minOccurs="0"
        maxOccurs="unbounded"/>
      <xs:element name="deliveryMethod" type="deliveryMethodType"
        maxOccurs="unbounded"/>
    </xs:sequence>
    <xs:attribute name="serviceId" type="xs:anyURI" use="required"/>
  </xs:complexType>

  <xs:complexType name="deliveryMethodType">
    <xs:attribute name="associatedProcedureDescriptionURI"
      type="xs:anyURI" use="optional"/>
    <xs:attribute name="protectionDescriptionURI" type="xs:anyURI"
      use="optional"/>
    <xs:attribute name="sessionDescriptionURI" type="xs:anyURI"
      use="required"/>
  </xs:complexType>

  <xs:complexType name="nameType">
    <xs:simpleContent>
      <xs:extension base="xs:string">
        <xs:attribute name="lang" type="xs:language" use="optional"/>
      </xs:extension>
    </xs:simpleContent>
  </xs:complexType>
</xs:schema>
```

5.2.2.5 Example MBMS User Service [Bundle](#) Description Instances

The following User Service [Bundle](#) Description instance is an example of a simple fragment. This fragment includes only the mandatory elements.

```
<?xml version="1.0" encoding="UTF-8"?>
<bundleDescription
  <userServiceDescription
    xmlns="http://www.example.com/3gppUserServiceDescription"
    xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance">
    <userServiceDescription
      userServiceId="urn:3gpp:0010120123hotdog">
      <deliveryMethod
        sessionDescriptionURI="http://www.example.com/3gpp/mbms/session1.sdp"/>
```

```

</userServiceDescription>
</bundleDescription>

```

The following User Service Description instance is an example of a fuller fragment.

```

<?xml version="1.0" encoding="UTF-8"?>
<bundleDescription<userServiceDescription
  xmlns="www.example.com/3gppUserServiceDescription"
  xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"
  fecDescriptionURI="http://www.example.com/3gpp/mbms/session1-fec.sdp">
  <userServiceDescription
    serviceId="urn:3gpp:1234567890coolcat">
    <name lang="EN">something in english</name>
    <name lang="DE">something in german</name>
    <name lang="FR">something in french</name>
    <name lang="FI">something in finnish</Name>
    <serviceLanguage>EN</serviceLanguage>
    <serviceLanguage>DE</serviceLanguage>
    <deliveryMethod
      sessionDescriptionURI="http://www.example.com/3gpp/mbms/session1.sdp"/>
    <deliveryMethod
      sessionDescriptionURI="http://www.example.com/3gpp/mbms/session2.sdp"
      associatedProcedureDescriptionURI=
        "http://www.example.com/3gpp/mbms/procedureX.xml"/>
    <deliveryMethod
      sessionDescriptionURI="http://www.example.com/3gpp/mbms/session3.sdp"
      associatedProcedureDescriptionURI=
        "http://www.example.com/3gpp/mbms/procedureY.xml"/>
    <deliveryMethod
      sessionDescriptionURI="http://www.example.com/3gpp/mbms/session4.sdp"
    </userServiceDescription>
  </bundleDescription>

```

5.2.2.6 FEC Repair Stream Description

The streaming delivery method's FEC has separate stream for repair data, which is described by the FEC Repair Stream Description. The FEC Repair Stream Description shall be done using SDP [14]. This SDP file is referenced by the *bundleDescription* element in the service description. The FEC Repair Stream described is common for all FEC protected packet flows within the MBMS User Service Bundle Description instance.

***** Next Change *****

5.5 MBMS Protocols

Figure 9 illustrates the protocol stack used by MBMS User services. The grey-shaded protocols and functions are outside of the scope of this specification. MBMS security functions and the usage of HTTP-digest and SRTP are defined in TS 33.246 [20].

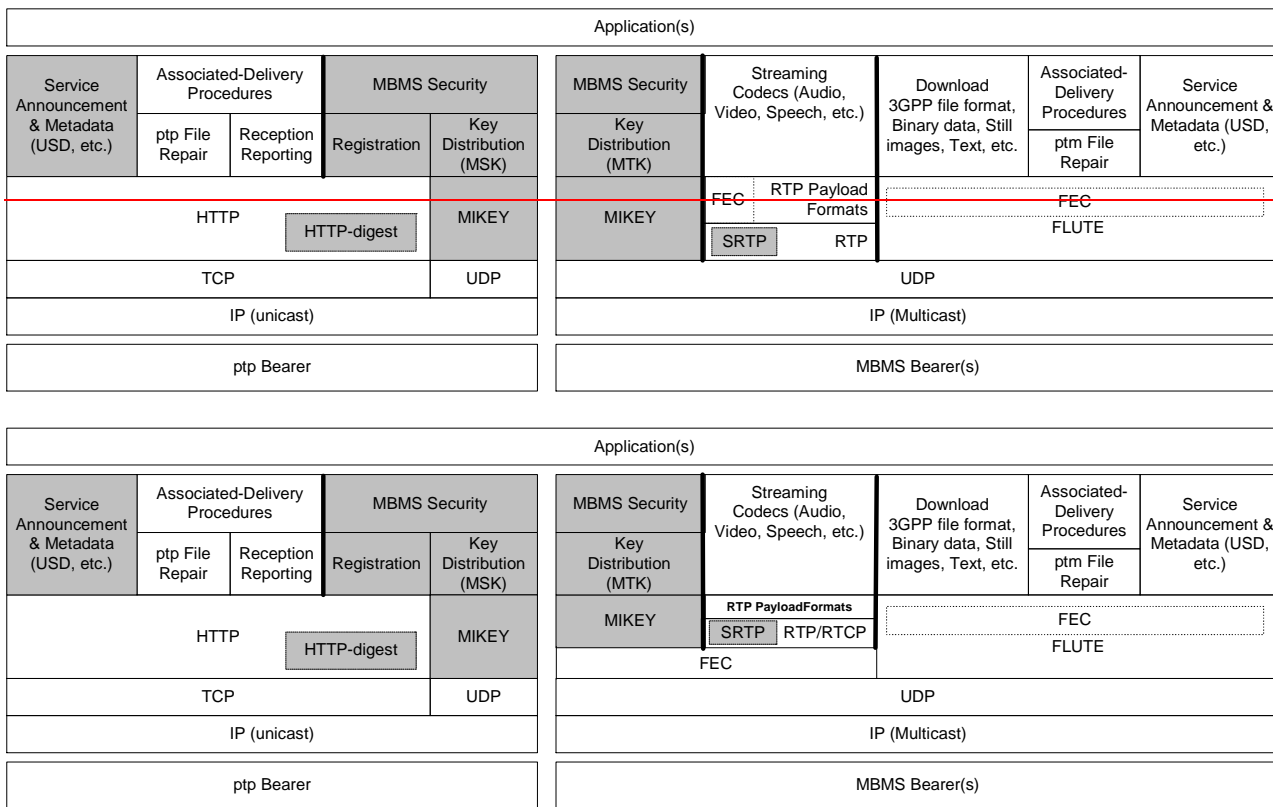


Figure 9: Protocol stack view of the MBMS User Services

***** Next Change *****

7.2.1 Fragmentation of Files

Fragmentation of files shall be provided by a blocking algorithm (which calculates source blocks from source files) and a symbol encoding algorithm (which calculates encoding symbols from source blocks).

~~Exactly one encoding symbol shall be carried in the payload of one FLUTE packet.~~

7.2.2 Symbol Encoding Algorithm

The "Compact No-Code FEC scheme" [12] (FEC Encoding ID 0, also known as "Null-FEC") shall be supported.

~~NOTE: the support of any other symbol encoding scheme is still under discussion~~

~~The "MBMS FEC scheme" is described in clause 7.2.11.~~

~~The UE supporting MBMS user service shall support a decoder for the "MBMS FEC scheme".~~

~~If the MBMS UE receives a mathematically sufficient set of encoding symbols generated according to the encoder specification in Annex B for reconstruction of a source block then the decoder shall recover the entire source block. Note that the example decoder described in annex B fulfils this requirement.~~

~~The "MBMS FEC scheme" shall be supported. MBMS forward error correction is indicated by the FEC Encoding ID defined in Section 7.2.11 below. The MBMS forward error correction code is defined in Annex B.~~

7.2.3 Blocking Algorithm

~~The~~ [In the case of the Compact No-Code FEC scheme \[12\] \(FEC Encoding ID 0\), then the "Algorithm for Computing Source Block Structure" described within the FLUTE specification \[9\] shall be used.](#)

[In the case of MBMS forward error correction, then the algorithm defined in Annex B shall be used.](#)

[The values of \$N\$, \$Z\$, \$T\$ and \$A\$ shall be set such that the sub-block size is less than 256KB.](#)

***** Next Change *****

7.2.7 Signalling of Parameters with Basic ALC/FLUTE Headers

FLUTE and ALC mandatory header fields shall be as specified in [9, 10] with the following additional specializations:

- The length of the CCI (Congestion Control Identifier) field shall be 32 bits and it is assigned a value of zero (C=0).
- The Transmission Session Identifier (TSI) field shall be of length 16 bits (S=0, H=1, 16 bits).
- The Transport Object Identifier (TOI) field should be of length 16 bits (O=0, H=1).
- Only Transport Object Identifier (TOI) 0 (zero) shall be used for FDT Instances.
- The following features may be used for signalling the end of session and end of object transmission to the receiver:
 - The Close Session flag (A) for indicating the end of a session.
 - The Close Object flag (B) for indicating the end of an object.

In FLUTE the following applies:

- The T flag shall indicate the use of the optional "Sender Current Time (SCT)" field (when T=1).
- The R flag shall indicate the use of the optional "Expected Residual Time (ERT)" field (when R=1).
- The LCT header length (HDR_LEN) shall be set to the total length of the LCT header in units of 32-bit words.
- For "Compact No-Code FEC scheme" [12], the ~~payload-FEC Payload~~ ID shall be set according to [13] such that a 16 bit SBN (Source Block Number) and then the 16 bit ESI (Encoding Symbol ID) are given.
- [For "MBMS FEC scheme", the FEC Payload ID shall be set according to Section 7.2.10 below.](#)

***** Next Change *****


7.2.9 Signalling of Parameters with FDT Instances


The FLUTE FDT Instance schema (RFC 3926 [9]) shall be used. In addition, the following applies to both the session level information and all files of a FLUTE session.

The inclusion of these FDT Instance data elements is mandatory according to the FLUTE specification:

- [_Content-Location](#) (URI of a file).
- [_TOI](#) (Transport Object Identifier of a file instance).
- [_Expires](#) (expiry data for the FDT Instance).

Additionally, the inclusion of these FDT Instance data elements is mandatory:


 Content-Length (source file length in bytes).


 Content-Type (content MIME type).

- [FEC Encoding ID](#).

[Other FEC Object Transmission Information specified by the FEC scheme in use :](#)

- [NOTE: The FEC Object Transmission Information elements used are dependent on the FEC scheme, as indicated by the FEC Encoding ID.](#)

 FEC-OTI-Maximum-Source-Block-Length.


 FEC-OTI-Encoding-Symbol-Length.

 FEC-OTI-Max-Number-of-Encoding-Symbols.


- [FEC-OTI-Scheme-Specific-Info](#)

NOTE 1: RFC 3926 [9] describes which part or parts of an FDT Instance may be used to provide these data elements.

These optional FDT Instance data elements may or may not be included for FLUTE in MBMS:

 Complete (the signalling that an FDT Instance provides a complete, and subsequently unmodifiable, set of file parameters for a FLUTE session may or may not be performed according to this method).

 ~~FEC-OTI-FEC-Instance-ID~~.

 Content-Encoding.

NOTE 2: The values for each of the above data elements are calculated or discovered by the FLUTE sender.

[The FEC-OTI-Scheme-Specific-Info FDT Instance data element contains information specific to the FEC scheme indicated by the FEC Encoding ID encoded using base64..](#)

7.2.10 FDT Schema

The following XML Schema shall be use for the FDT Instance:

```
<?xml version="1.0" encoding="UTF-8"?>
<xs:schema xmlns:xs="http://www.w3.org/2001/XMLSchema"
  xmlns:fl="http://www.example.com/flute"
  elementFormDefault:xs="qualified"
  targetNamespace:xs="http://www.example.com/flute">
  <xs:element name="FDT-Instance">
    <xs:complexType>
      <xs:sequence>
        <xs:element name="File" maxOccurs="unbounded">
          <xs:complexType>

            <xs:sequence>
              <xs:element name="Group" type="xs:string" minOccurs="0" maxOccurs="unbounded"/>
            </xs:sequence>
          <xs:any processContents="skip" minOccurs="0" maxOccurs="unbounded"/>

          <xs:attribute name="Content-Location" type="xs:anyURI" use="required"/>
          <xs:attribute name="TOI" type="xs:positiveInteger" use="required"/>
          <xs:attribute name="Content-Length" type="xs:unsignedLong" use="optional"/>
          <xs:attribute name="Transfer-Length" type="xs:unsignedLong" use="optional"/>
          <xs:attribute name="Content-Type" type="xs:string" use="optional"/>
          <xs:attribute name="Content-Encoding" type="xs:string" use="optional"/>
          <xs:attribute name="Content-MD5" type="xs:base64Binary" use="optional"/>
          <xs:attribute name="FEC-OTI-FEC-Encoding-ID" type="xs:unsignedLong" use="optional"/>
          <xs:attribute name="FEC-OTI-FEC-Instance-ID" type="xs:unsignedLong" use="optional"/>
          <xs:attribute name="FEC-OTI-Maximum-Source-Block-Length"
            type="xs:unsignedLong" use="optional"/>
          <xs:attribute name="FEC-OTI-Encoding-Symbol-Length"
            type="xs:unsignedLong" use="optional"/>
          <xs:attribute name="FEC-OTI-Max-Number-of-Encoding-Symbols"
            type="xs:unsignedLong" use="optional"/>
        
```

```

<xs:attribute name="FEC-OTI-Scheme-Specific-Info"
type="xs:base64Binary" use="optional"/>

<xs:anyAttribute processContents="skip"/>
</xs:complexType>
</xs:element>
</xs:sequence>

<xs:sequence>
<xs:element name="Group" type="xs:string" minOccurs="0" maxOccurs="unbounded"/>
</xs:sequence>
<xs:any processContents="skip" minOccurs="0" maxOccurs="unbounded"/>

<xs:attribute name="Expires" type="xs:string" use="required"/>
<xs:attribute name="Complete" type="xs:boolean" use="optional"/>
<xs:attribute name="Content-Type" type="xs:string" use="optional"/>
<xs:attribute name="Content-Encoding" type="xs:string" use="optional"/>
<xs:attribute name="FEC-OTI-FEC-Encoding-ID" type="xs:unsignedLong" use="optional"/>
<xs:attribute name="FEC-OTI-FEC-Instance-ID" type="xs:unsignedLong" use="optional"/>
<xs:attribute name="FEC-OTI-Maximum-Source-Block-Length"
type="xs:unsignedLong" use="optional"/>
<xs:attribute name="FEC-OTI-Encoding-Symbol-Length"
type="xs:unsignedLong" use="optional"/>
<xs:attribute name="FEC-OTI-Max-Number-of-Encoding-Symbols"
type="xs:unsignedLong" use="optional"/>
<xs:attribute name="FEC-OTI-Scheme-Specific-Info"
type="xs:base64Binary" use="optional"/>
<xs:anyAttribute processContents="skip"/>
</xs:complexType>
</xs:element>
</xs:schema>

```

***** Next Change *****

7.2.11 FEC Scheme definition

7.2.11.1 General

This clause defines an FEC encoding scheme for the MBMS forward error correction code defined in Annex B for the download delivery method. This scheme is identified by FEC Encoding ID [TBA]. The FEC Payload ID format and FEC Object Transmission Information format are as defined in the following clauses.

7.2.11.2 FEC payload ID

The FEC Payload ID shall be a 4 octet field defined as follows:

| | |
|----------------------------------|---------------------------------|
| <u>Source Block Number (SBN)</u> | <u>Encoding Symbol ID (ESI)</u> |
|----------------------------------|---------------------------------|

Source Block Number (SBN), (16 bits): An integer identifier for the source block that the encoding symbols within the packet relate to.

Encoding Symbol ID (ESI), (16 bits): An integer identifier for the encoding symbols within the packet.

The interpretation of the Source Block Number and Encoding Symbol Identifier is defined in Annex B.

7.2.11.3 FEC Object Transmission Information

The FEC Object Transmission information shall consist of:

- [The FEC Encoding ID](#)
- [The Transfer Length \(F\)](#)
- [The parameters T, Z, N and A defined in Annex B.](#)

[When EXT-FTI is used to communicate the Object Transmission Information, the FEC Encoding ID and Transfer Length shall be coded according to FLUTE \[9\]. The other parameters shall be encoded in the FEC Encoding ID specific portion of the EXT_FT I field as shown in Figure 10a below.](#)

| | | |
|---|--|--|
| General EXT_FT I format | Encoding Symbol Length (T) | |
| Number of Source Blocks (Z) | Number of Sub-Blocks (N) | Symbol Alignment Parameter (A) |

Figure 10a: FEC Encoding ID-specific EXT_FT I format

[The parameters T and Z are 16 bit unsigned integers, N and A are 8 bit unsigned integers.](#)

[When the FDT is used to deliver the FEC Object Transmission Information, then the FEC Encoding ID, Transfer Length \(F\) and Encoding Symbol Length \(T\) shall be encoded using the Transfer-Length, FEC-OTI-Encoding-ID and FEC-OTI-Encoding-Symbol-Length elements defined in Section 7.2.10. The remaining parameters Z, N, A, shall be encoded as a 4 byte field within the FEC-OTI-Scheme-Specific-Info field, according to the format specified in Figure 10a above, excepting the Encoding Symbol Length field.](#)

***** Next Change *****

7.3.2.8 FEC capabilities and related parameters

A new FEC-declaration attribute is defined which results in, e.g.:

- a=FEC-declaration:0 encoding-id=128; instance-id=0

This ~~attribute may~~ [can be used on both session-level \(and so the first instance \(fec-ref=0\) becomes the default for all media\)-and media-level. Multiple instances are allowed to specify several different FEC declarations. The attribute is used on session level to define FEC declarations used by multiple media blockscomponents, that e](#)Each media ~~block~~ [blockcomponent references an FEC declaration using the “a=FEC” attribute. On media level it is used to define FEC declarations which are only valid for a single media blockcomponent. If FEC declarations on both session and media level uses the same reference number \(fec-ref\) then the media level the valid one for that media blockdeclaration takes precedence for that media component.](#)

~~to specify differences between media-~~This [attribute](#) is optional [to use for the download delivery method](#) as the information will be available elsewhere (e.g. FLUTE FDT Instances). If this attribute is not used, and no other FEC-OTI information is ~~signaled~~ [signalled](#) to the UE by other means, the UE may assume that support for FEC id 0 is sufficient capability to enter the session.

A new FEC-declaration [reference](#) attribute ~~shall be is~~ defined which results in, e.g.:

- ☐ a=FEC:0

This is ~~only~~ a media-level [only](#) attribute, used as a short hand to ~~inherit~~ [reference](#) one of one or more ~~session-level~~ FEC-declarations ~~to a specific media.~~

The syntax for the attributes in ABNF - RFC 2234 [23] is:

- ☐ sdp-fec-declaration-line = "a=FEC-declaration:" fec-ref SP fec-enc-id ";" [SP fec-inst-id]_ CRLF
- ☐ fec-ref = 1*3DIGIT (value is the SDP-internal identifier for FEC-declaration).
- ☐ fec-enc-id = "encoding-id=" enc-id
- ☐ end-id = 1*DIGIT (value is the FEC Encoding ID used).

~~fec-inst-id = "instance-id=" inst-id~~

~~inst-id = 1*DIGIT (value is the FEC Instance ID used).~~

~~sdp-fec-line = "a=FEC:" fec-ref_ CRLF~~

~~fec-ref = 1*DIGIT (value is the FEC declaration identifier).~~

***** Next Change *****

8.2.2 FEC mechanism for RTP

~~NOTE1: This scheme is intended to be generic. If the selected FEC scheme(s) does not fit, it can be modified.~~

~~NOTE2: The specification text is FFS.~~

The “MBMS FEC scheme” is described in clause 8.2.2.8.

The MBMS UE shall support a decoder for the “MBMS FEC scheme”.

This section defines a generic mechanism for applying Forward Error Correction to streaming media. The mechanism consists of three components:

- (i) construction of an FEC source block from the source media packets belonging to one or several UDP packet flows related to a particular segment of the stream(s) (in time). The UDP flows ~~is~~ include RTP, RTCP, SRTP and MIKEY packets.
- (ii) modification of source packets to indicate the position of the source data from the source packet within the source block
- (iii) definition of repair packets, sent over UDP, which can be used by the FEC decoder to reconstruct missing portions of the source block.

The mechanism does not place any restrictions on the source data which can be protected together, except that the source data is carried over UDP. The data may be from several different UDP flows that are protected jointly.

~~The generic mechanism for systematic FEC of RTP streams consists of two RTP payload formats, one for FEC source packets, one for FEC repair packets, and their related signaling. In addition, the construction of a FEC source block is specified. A receiver supporting the streaming delivery method shall support the payload-packet format for FEC source packets and ~~shall~~ may also support the payload-packet format for FEC repair packets.~~

At the sender, the mechanism begins by processing original RTP-UDP packets to create:

- (i) a stored copy of the original packets in the form of a source block and
- (ii) FEC source packets for transmission to the receiver

After constructing the source block from the original RTP-UDP payloads~~packets~~ to be protected and their flow identity (based on destination IP address and UDP port), the FEC encoder generates the desired amount of FEC protection data, i.e. encoding symbols. These encoding symbols are then sent using the FEC repair packet payload-format to the receiver. The FEC repair packets are sent to a UDP destination port use an SSRC different from any of the original RTP-UDP-packets' destination port(s) as indicated by the signaling.~~SSRC, but are sent within the same RTP session. Doing so avoids non-continuous sequence numbering spaces for both the FEC repair packets and the original RTP packets.~~

The receiver recovers the original RTP-packets directly from the FEC source packets and buffers ~~it~~ them at least the min-buffer-time to allow time for the FEC repair. The receiver uses the FEC source packets to construct a (potentially incomplete) copy of the source block, using the Source FEC Payload ID and destination UDP port in each packet to determine where in the source block the packet ~~shall~~ ould be placed ~~and what value the flow ID has~~. In indication of the UPD flow (i.e. destination IP address and UDP port) the packet is part of is included in the source block with the UDP payload.

If any FEC source packets have been lost, but sufficient FEC source and FEC repair packets have been received, FEC decoding can be performed to recover the FEC source block. The original ~~RTP~~ packets UDP payload and UDP flow identity~~stream origin~~ can then be extracted from the source block and ~~buffered as normal~~be provided to the upper layer. If not enough FEC source and repair packets were received, only the original packets that were received as FEC source packets will be available. The rest of the original packets are lost.

If the MBMS UE receives a mathematically sufficient set of encoding symbols generated according to the encoder specification in Annex B for reconstruction of a source block then the decoder shall recover the entire source block. Note that the example decoder described in Annex B fulfils this requirement.

Note that the receiver must be able to buffer all the original packets and allow time for the FEC repair packets to arrive and FEC decoding to be performed before media playout begins. The min-buffer-time~~e-MIME~~ parameter specified in clause 8.3.1.9~~2.4.12~~ helps the receiver to determine a sufficient duration for initial start-up delay.

The Source and Repair FEC payload IDs are used to associate the FEC source packets and FEC repair packets, respectively, to a source block. The Source and Repair FEC payload ID formats are part of the definition of the FEC scheme. Each FEC scheme is identified by an FEC ~~encoding~~Encoding ID and, in the case of underspecified FEC schemes, FEC Instance ID, value(s) ~~for underspecified FEC encoding IDs~~. One FEC scheme for the streaming delivery method is specified in clause 8.2.2~~4~~.6. Any FEC schemes using the ~~RTP payload~~packet formats defined in this specification shall be systematic FEC codes and may use different FEC payload ID formats for FEC source packets and FEC repair packets.

The protocol architecture is illustrated in Figure 11 below.

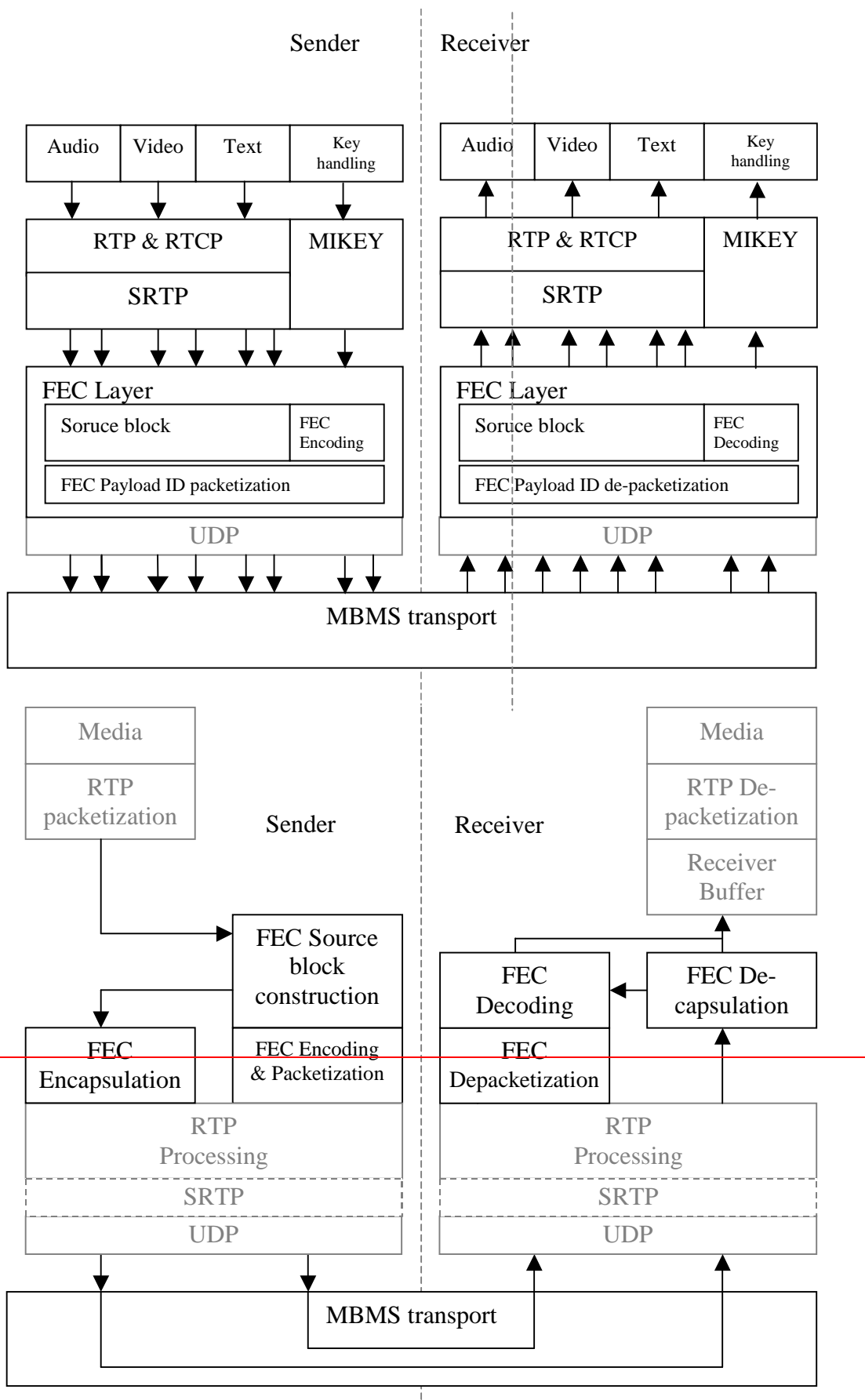


Figure 11: FEC mechanism for [the streaming RTP delivery method](#) interaction diagram

Figure 11 depicts how one or more out of several possible packet flows of different types (Audio, video, text RTP and RTCP flows, MIKEY flow) are sent to the FEC layer for protection. The source packets are modified to carry the FEC payload ID and a new flow with repair data is generated. The receiver takes the source and repair packets and buffers them to perform, if necessary, the FEC decoding. After appropriate buffering received and recovered source packets are forwarded to the higher layers. The arrows in the figure indicate distinct data flows.

8.2.2.1 Sending Terminal Operation (Informative)

It is assumed that the sender has constructed or received original ~~RTP data~~ packets for the session. These may be RTP, RTCP, MIKEY or other UDP packets. The following procedures are based on the UDP payload and the identity of the UDP flow. ~~The actual UDP header is not yet formed, but source and destination ports are already known and the different packet flows can be identified, comprising RTP header, profile specific extensions, variable sized fields such as CSRC list and extension headers, the payload headers and the payload.~~

In order to FEC protect a sequence of original data ~~such~~ packets, the sender constructs a source block as specified in section 8.2.2.6, to which the FEC algorithm is to be applied, and ~~encapsulates~~ includes the original ~~RTP~~ source packet data within FEC source packets. The following operations describe a possible way to generate compliant FEC source packet and FEC repair packet streams:

1. Each original ~~RTP~~ packet is placed in the source block. In doing so, the Source FEC Payload ID information to be included in the FEC payload ID of the FEC source packet can be determined. In the source block the identity of the packet's flow is marked using the Flow ID. See Clause 8.2.2.54 and 8.2.2.7 for details.
2. The FEC source packet is constructed according to Clause 8.2.2.43. The identity of the original flow is maintained by the source packet ~~It is identified through~~ by the use of an RTP Payload Type that indicates the use of the FEC source packet payload format encapsulating an RTP packet of the original Payload Type ~~the destination~~ a UDP port number and destination IP address, which has been advertised (for example using SDP), as carrying FEC source packets generated from an original stream of a particular protocol (e.g. RTP, RTCP, SRTP, MIKEY etc.). See Clause 8.2.2.11
3. The FEC source packet generated is sent according to ~~RTP-UDP~~ procedures defined in [UDP~~RTP~~].

When a source block is complete, the FEC encoder generates encoding symbols and places these symbols into FEC repair packets, to be conveyed in the same RTP session as the FEC source packets, but using a different SSRC to the receiver(s). These repair packets are sent using normal ~~RTP-UDP~~ procedures to a unique destination port to separate it from any of the source packet flows.

8.2.2.2 Receiving Terminal Operation (Informative)

The following describes a possible receiver algorithm, when receiving an FEC source or repair packet ~~in a given RTP session:~~

1. If a FEC source packet (~~with a Payload Type~~ UDP port that indicates this received (as indicated by the which UDP port it on which was received) ~~e-FEC source packet payload format~~) is received:
 - a. The original source FEC packet is reconstructed by removing the Source FEC Payload ID ~~and changing the Payload type back to the original Payload Type (which can be derived from the received Payload Type).~~ The resulting packet is buffered to allow time for the FEC repair.
 - b. The resulting packet is placed into the source block according to the information in the Source FEC Payload ID ~~and the source block format described in 8.2.2.35.~~ The UDP port the packet was received on is used to determine the Flow ID written into the source block.
2. If an FEC repair packet is received (as indicated by the ~~Payload Type~~ UDP port), the contained encoding symbols are placed into ~~the an~~ FEC source-encoding block according to the Repair FEC Payload ID.

3. If at least one source packet is missing (~~as detected by the gaps in the sequence number space~~), then FEC decoding may be desirable. The FEC decoder determines if the ~~source encoding~~ block constructed in steps 1 and 2 contains enough symbols from the source and repair packets for decoding and, if so, performs the decoding operation.
4. Any missing source packets that were reconstructed during the decoding operation are then buffered as normal received ~~RTP~~-packets (see step 1a above).

Note that the above procedure may result in that not all original ~~RTP~~-packets are recovered, and they must simply be marked as being lost.

Obviously, buffering and packet re-ordering are required to insert any reconstructed packets in the appropriate place in the packet sequence if that is necessary according to the used higher layer protocol (RTP, RTCP or MIKEY). To allow receivers to determine the minimal start-up buffering requirement for FEC decoding, the min-buffer-time ~~MIME~~ parameter indicates a minimum initial buffering time that is sufficient regardless of the position of the stream in which the reception starts.

8.2.2.3 ~~RTCP Statistics~~(Void)

~~RTCP Statistics, if required, shall be based on the FEC source packets and FEC repair packets and not on the original source RTP packets. The synchronization information (RTP and NTP Timestamp fields) for the source packet stream shall be identical to the original packet streams as the RTP timestamp values for the source is not modified between source and original packets.~~

~~In the context of MBMS 3GPP Rel-6, RTCP RR shall be turned off by SDP RR bandwidth modifiers [Ref].~~

8.2.2.4 ~~RTP Packet~~Payload format for FEC source ~~RTP~~ packets

The ~~RTP payload~~packet format for FEC source packets shall be used to encapsulate an original ~~RTP-UDP~~ packet. ~~This payload format is identified by the media type defined in clause 8.2.1.14.~~ As depicted in Figure 12, it consists of the ~~RTP header, the RTP payload header in the form of the~~original UDP packet, followed by the Source FEC payload ID, ~~and the Payload.~~

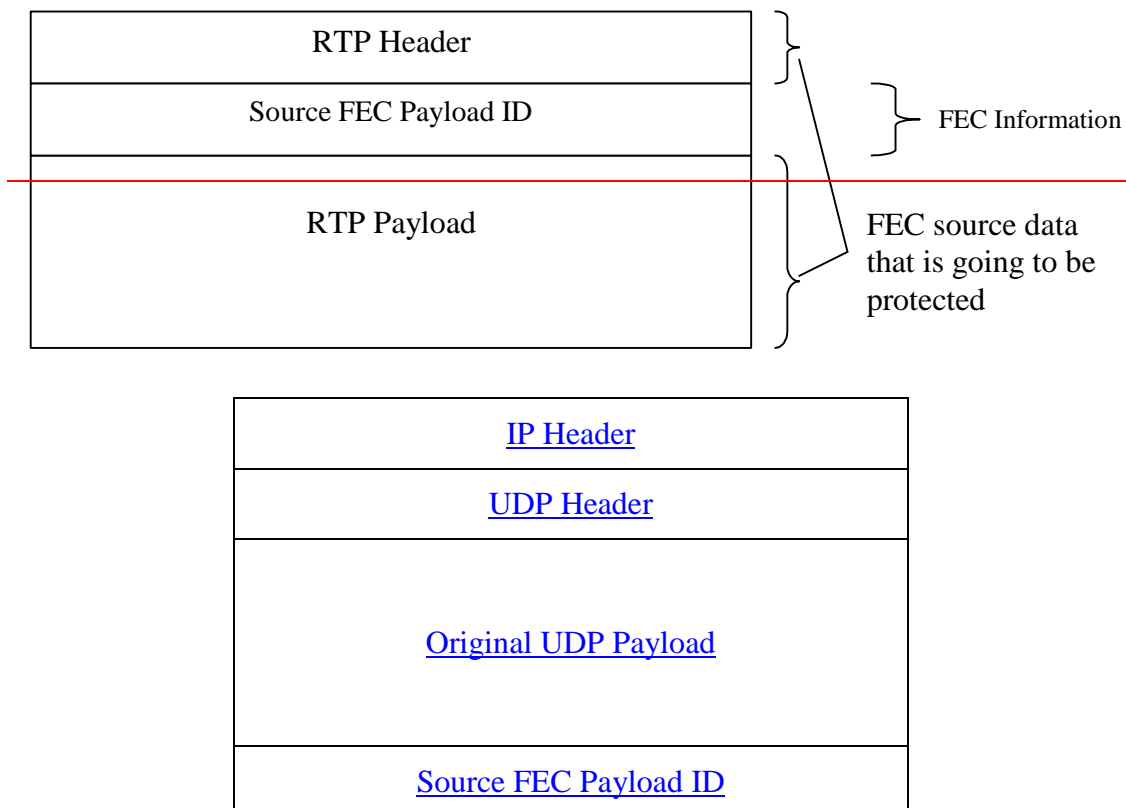


Figure 12: Structure of the FEC ~~payload packet~~ format for FEC source ~~RTP~~ packets

~~The destination IP address and UDP port shall be set to the values this flow are indicated to sent to according to the as indicated in the session control signalling. Thus This ensures that the receiver can determine which protocols and FEC P payload ID formats are used for this flow. The Payload Type field destination UDP port in the RTP UDP header shall be a value indicating the FEC source packet payload format carrying payload of the original Payload Type UDP packet.~~ The remaining fields in the IP and RTP-UDP headers shall be set according to their specifications, to the same values as those of the original source RTP-UDP packet.

The RTP-UDP payload shall consist of the original UDP Payload followed by the Source FEC Payload ID ~~followed by the original RTP packet's payload.~~

The Source FEC Payload ID consists of information required for the operation of the FEC algorithm. Its construction is specified in section 8.2.2.7.8.

The FEC source FEC packets over IP and UDP are indicated to be used for a flow by using anyone of the SDP protocol identifiers “UDP/MBMS-FEC/RTP/AVP”, “UDP/MBMS-FEC/RTP/SAVP” depending on the upper layer protocol RTP/AVP or RTP/SAVP respectively. If MIKEY is FEC protected and encapsulated in source packets, then it is indicated in the security description using the *fecProtection* element and the destination IP address.

8.2.2.5 RTP Payload Packet Format for Repair packets

The RTP payload packet format for FEC repair packets carries, as its payload, encoding symbols generated by the FEC encoding process. The format of a FEC repair packet is depicted in Figure 13. The RTP-UDP payload consists of the Repair FEC Payload ID, and one or more encoding symbols. ~~The RTP payload format is identified by the media type defined in Clause 8.1.2.13.~~

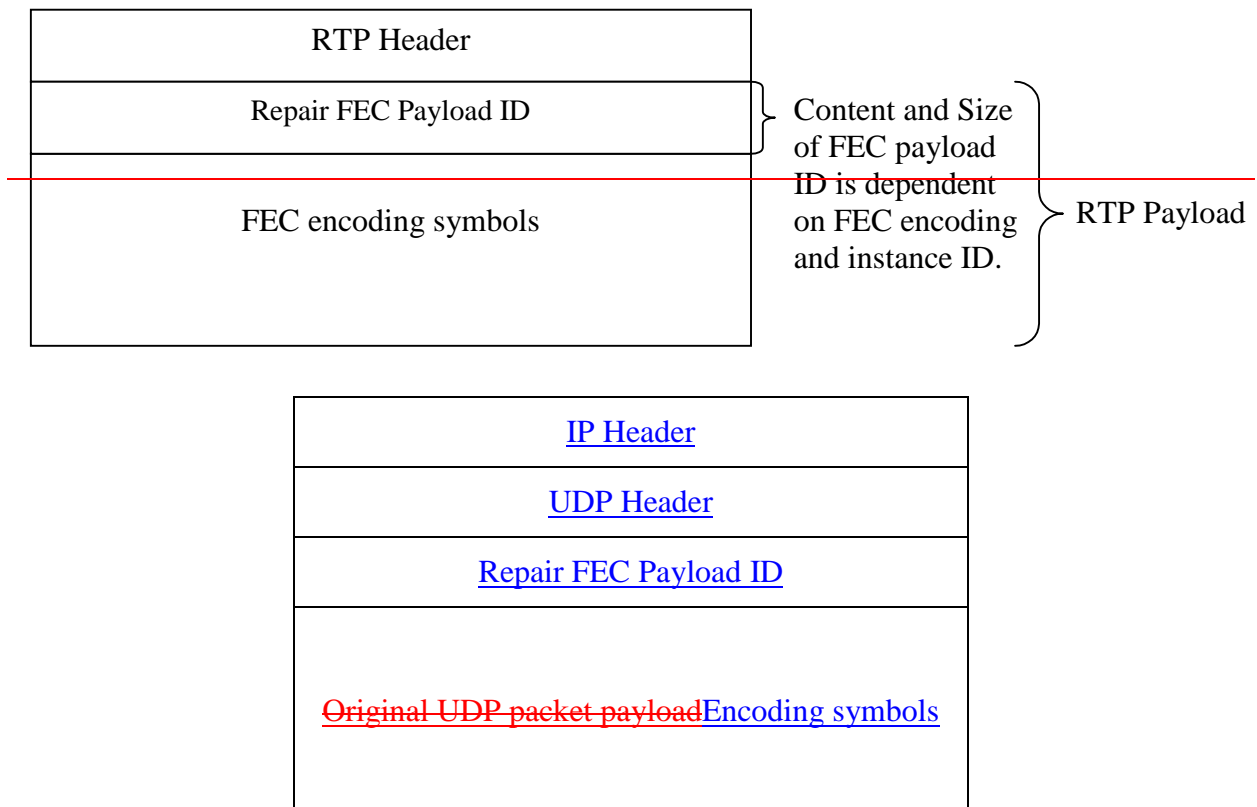


Figure 13: RTP payload structure for repair RTP packets

The repair packet sent over IP and UDP is indicated in the SDP using the protocol identifier “UDP/MBMS-REPAIR”.

The RTP header information for the FEC repair packet is set as follows:

Marker bit: The marker bit shall be set 1 for the last FEC repair packet sent for each source block, and otherwise set to 0.

Timestamp: The timestamp rate shall be 10 kHz and shall be set to a time corresponding to the packet's transmission time. The timestamp value has no use in the actual FEC protection process and is only set to a value to produce reasonable resolution for buffer handling, arrival measuring and jitter calculation.

Sequence number: Is set in accordance with RFC 3550 [6]. The sequence number is primarily used to detect losses of the FEC repair packets. All FEC repair packets for a source block shall be sent in consecutive order, and be complete before starting the transmission of the next source block.

Payload type (PT): identifies the payload format for FEC repair packets and the FEC parameters for the payload. The FEC encoding ID, FEC instance ID, and any additional FEC object transmission information is all determined through the Payload Type.

SSRC: One SSRC is used per source SSRC. The SSRC used by the FEC repair packet payload format shall be different the one used by the FEC source packets. The binding of the source SSRC to the repair SSRC shall be performed using the RTCP SDES CNAME, which shall be identical for the two SSRCs.

The Repair FEC Payload ID is of a fixed in size for a given Payload Type, as the Payload Type identifies the FEC scheme employed and its parameters (if any). A FEC scheme and its Source and Repair FEC payload ID formats are defined in 8.2.1.6.

8.2.2.6 Structure of the FEC source block

This clause defines the layout of the FEC source block.

The FEC source block shall contain at least one complete **RTP-UDP** packet payload (i.e. excluding the IP and UDP headers), and ~~two~~ three octets indicating the UDP flow from which the packet was taken and the length of the UDP RTP packet. Note: this implies that no source RTP-UDP packet be larger than the length of the FEC source block minus ~~23~~.

Let

n be the number of **RTP-UDP** packets in the source block. n is determined dynamically during the source block construction process.

R_i denote the octets of the **RTP header, payload header(s), and RTP-UDP** payload of the i th **RTP-UDP** packet to be added to the source block.

l_i be the length of R_i in octets

L_i denote two octets representing the value of l_i in network byte order (high order octet first)

f_i denote an integer “flow ID” identifying the UDP flow from which the i^{th} packet was taken

F_i denote a single octet representing the value of f_i

S_i be the smallest integer such that $s_i T \geq (l_i + 32)$

P_i denote $s_i T - (l_i + 32)$ zero octets. Note: P_i are padding octets to align the start of each **UDP-RTP** packet with the start of a symbol.

T be the source symbol size in bytes.

Then, the source block is constructed by concatenating F_i, L_i, R_i, P_i for $i = 1, 2, \dots, n$. and the source block size, $S = \sum \{s_i T, i=1, \dots, n\}$.

A UDP flow is uniquely defined by an IP source and destination address and UDP source and destination port value. The assignment of Flow ID values to UDP flows is described in Section 8.2.2.11 and 8.3.1.10.

8.2.2.7 FEC block Construction algorithm and example (informative)

When the original RTP-UDP packet is placed into the source block, the value of the UDP flow identifier, F, followed by the value of the UDP payload length, L, are first written as a single byte and two-byte value in network byte order (i.e. with high order byte first) respectively into the first available bytes in the source block, followed by the RTP-UDP packet payload itself (including the RTP header, but i.e. not including the IP/UDP headers). Following this, if the next available byte is not the first byte of a new symbol, then padding bytes up to the next symbol boundary shall be included using the value 0 in each byte. As long as any source RTP-UDP packets remain to be placed, the procedure is repeated starting each RTP-UDP flow identifier packet length indicator at the start of the next encoding symbol.

An example of forming a source block is given in figure 14 below. In this example, three RTP-UDP packets of lengths 26, 52 and 103 have been placed into a source block with symbol size T = 16 bytes. The first two packets are from UDP flow 0 and the third from UDP flow 1. Each entry in Figure 14 is a byte and the rows correspond to the source symbols and are numbered from 0 to 12. B_{i,j} denotes the (j+1)th byte of the (i+1)th RTP packet.

| | | | | | | | | | | | | | | | |
|-------------------|-------------------|-------------------|-------------------|-------------------|-------------------|-------------------|--------------------|--------------------|--------------------|-------------------|-------------------|-------------------|-------------------|-------------------|-------------------|
| 26 | 25 | | B _{0,0} | B _{0,1} | B _{0,2} | B _{0,3} | B _{0,4} | B _{0,5} | B _{0,6} | B _{0,7} | B _{0,8} | B _{0,9} | B _{0,10} | B _{0,11} | B _{0,12} |
| B _{0,13} | B _{0,14} | B _{0,15} | B _{0,16} | B _{0,17} | B _{0,18} | B _{0,19} | B _{0,20} | B _{0,21} | B _{0,22} | B _{0,23} | B _{0,24} | B _{0,25} | 0 | 0 | 0 |
| 52 | 51 | | B _{1,0} | B _{1,1} | B _{1,2} | B _{1,3} | B _{1,4} | B _{1,5} | B _{1,6} | B _{1,7} | B _{1,8} | B _{1,9} | B _{1,10} | B _{1,11} | B _{1,12} |
| B _{1,13} | B _{1,14} | B _{1,15} | B _{1,16} | B _{1,17} | B _{1,18} | B _{1,19} | B _{1,20} | B _{1,21} | B _{1,22} | B _{1,23} | B _{1,24} | B _{1,25} | B _{1,26} | B _{1,27} | B _{1,28} |
| B _{1,29} | B _{1,30} | B _{1,31} | B _{1,32} | B _{1,33} | B _{1,34} | B _{1,35} | B _{1,36} | B _{1,37} | B _{1,38} | B _{1,39} | B _{1,40} | B _{1,41} | B _{1,42} | B _{1,43} | B _{1,44} |
| B _{1,45} | B _{1,46} | B _{1,47} | B _{1,48} | B _{1,49} | B _{1,50} | B _{1,51} | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 103 | 102 | | B _{2,0} | B _{2,1} | B _{2,2} | B _{2,3} | B _{2,4} | B _{2,5} | B _{2,6} | B _{2,7} | B _{2,8} | B _{2,9} | B _{2,10} | B _{2,11} | B _{2,12} |
| B _{2,13} | B _{2,14} | B _{2,15} | B _{2,16} | B _{2,17} | B _{2,18} | B _{2,19} | B _{2,20} | B _{2,21} | B _{2,22} | B _{2,23} | B _{2,24} | B _{2,25} | B _{2,26} | B _{2,27} | B _{2,28} |
| B _{2,29} | B _{2,30} | B _{2,31} | B _{2,32} | B _{2,33} | B _{2,34} | B _{2,35} | B _{2,36} | B _{2,37} | B _{2,38} | B _{2,39} | B _{2,40} | B _{2,41} | B _{2,42} | B _{2,43} | B _{2,44} |
| B _{2,45} | B _{2,46} | B _{2,47} | B _{2,48} | B _{2,49} | B _{2,50} | B _{2,51} | B _{2,52} | B _{2,53} | B _{2,54} | B _{2,55} | B _{2,56} | B _{2,57} | B _{2,58} | B _{2,59} | B _{2,60} |
| B _{2,61} | B _{2,62} | B _{2,63} | B _{2,64} | B _{2,65} | B _{2,66} | B _{2,67} | B _{2,68} | B _{2,69} | B _{2,70} | B _{2,71} | B _{2,72} | B _{2,73} | B _{2,74} | B _{2,75} | B _{2,76} |
| B _{2,77} | B _{2,78} | B _{2,79} | B _{2,80} | B _{2,81} | B _{2,82} | B _{2,83} | B _{2,84} | B _{2,85} | B _{2,86} | B _{2,87} | B _{2,88} | B _{2,89} | B _{2,90} | B _{2,91} | B _{2,92} |
| B _{2,93} | B _{2,94} | B _{2,95} | B _{2,96} | B _{2,97} | B _{2,98} | B _{2,99} | B _{2,100} | B _{2,101} | B _{2,102} | 0 | 0 | 0 | 0 | 0 | 0 |

Figure 14: Source block consisting of 3 source RTP packets of lengths 26, 52 and 103 bytes.

8.2.2.8 MBMS FEC scheme definition

This clause defines a FEC encoding ~~scheme~~ scheme for MBMS forward error correction as defined in Annex B for the streaming delivery method. This scheme and is identified by the FEC encoding ID [TBA]. It utilized the method for forming FEC source block as defined in Clause 8.2.1.3. It defines two different FEC ~~payload ID formats~~, one for FEC source ~~RTP~~ packets and another for FEC repair packets ~~encoding symbols that are used with the corresponding RTP payload formats.~~

NOTE: This clause will require some rewording after the final decision on the FEC scheme(s)

8.2.2.9 Source FEC Payload ID

The Source FEC payload ID is composed as follows:

| | |
|---------------------------|--------------------------|
| Source Block Number (SBN) | Encoding Symbol ID (ESI) |
|---------------------------|--------------------------|

Source Block Number (SBN), (8 ~~or~~ 16 bits): An integer identifier for the source block that the source data within the packet relates to. The I-D of the source block the media packet belongs to.

Encoding Symbol ID (ESI), (8 ~~or~~ 16 bits): The starting symbol index of the source packet in the source block.

Figure 15: Source FEC Payload ID

The interpretation of the Source Block Number and Encoding Symbol Identifier is defined in Annex B.

NOTE: The Source Block number could be 1 or 2 bytes and is independent of the FEC scheme—its length is a tradeoff decision between wrap-around resilience and overhead. It would also be possible to allow both lengths and distinguish by SDP. This needs to be decided.

NOTE: as far as we understand the 2 FEC proposals, the length of SBN and ESI is as follows:

| | Raptor codes | 2D-Reed-Solomon |
|-----|--------------|-----------------|
| SBN | 1 or 2 | 1 or 2 |
| ESI | 2 | 1 |

8.2.2.10 Repair FEC payload ID

The structure of the Repair FEC Payload ID is as follows:

| |
|------------------------------------|
| Source Block Number SBN |
| Encoding Symbol ID ESI |

| | |
|--------------------------------------|--------------------------|
| Source block length SBL | |
| Encoding block length EBL | |
| Symbol Length T | |
| Source Block Number (SBN) | Encoding Symbol ID (ESI) |
| Source Block Length (SBL) | |

Source Block Number (SBN), (8 or 16 bits): ~~An integer identifier for the source block that the repair symbols within the packet relate to. The ID of the source block the media packet belongs to.~~

Encoding Symbol ID (ESI), (8 or 16 bits): ~~integer identifier for the encoding symbols within the packet. The starting symbol index of the source packet in the source block.~~

Source Block Length (SBL), (8, 16, or 32 bits): The number of source symbols in the source block. The interpretation of the Source Block Number, Encoding Symbol Identifier and Source Block Length is defined in Annex B.

~~**Encoding Block Length (EBL), (not present, 8 bits, 16 bits or 32 bits):** The total number of symbols (source symbols of the source block plus encoding symbols generated from the source block) that are sent for the source block.~~

~~**Symbol Length (T), (not present, 8 or 16 bits):** The length of a symbol in bytes.~~

Figure 16: Repair FEC Payload ID

~~NOTE 1: The rest of this clause requires rewording after the FEC schemes supported have been selected.~~

~~NOTE 2: As far as we understand the 2 FEC proposals, the length of the various FEC payload ID fields is as follows:~~

| | Raptor codes | 2D Reed-Solomon |
|--|-------------------------|----------------------------|
| SBN | 4 or 2 | 4 or 2 |
| ESI | 2 | 4 |
| SBL | 2 | 4 |
| EBL | 0 | 4 |
| T | 2 (*) | 4 |
| NOTE: The (*) denotes fields that could probably be of zero length here, and conveyed in SDP. | | |

8.2.2.10a FEC Object Transmission Information

The FEC Object Transmission information shall consist of:

- the FEC Encoding ID
- the maximum source block length, in symbols
- the symbol size, in bytes

The symbol size and maximum source block length shall be encoded into a 4 octet field defined as follows:

| | |
|--------------------------|-----------------------------|
| Symbol Size (<i>T</i>) | Maximum Source Block Length |
|--------------------------|-----------------------------|

Symbol Size (*T*) (16 bits): The size of an encoding symbol, in bytes.

Maximum Source Block Length (16 bits): The maximum length of a source block, in symbols.

The interpretation of *T* is defined in Annex B.

The Source Block Length signalled within the Repair FEC Payload ID of any packet of a stream shall not exceed the Maximum Source Block Length signalled within the FEC Object Transmission Information for the stream.

The FEC Object Transmission Information shall be communicated as described in Section 8.2.2.14.

8.2.2.11 Hypothetical FEC Decoder

This clause specifies the hypothetical FEC decoder and its use to check packet stream and MBMS receiver conformance.

The hypothetical FEC decoder uses the packet stream, the transmission time of each packet, the initial buffering delay, and the SDP for the stream as inputs. The packet stream from the beginning of the FEC source block until the end of the stream shall comply with the hypothetical reference decoder as specified below when the initial buffer delay equals to the value of the min-buffer-time ~~MIME~~-parameter.

The maximum buffer size for MBMS streaming is 1 Mbytes. The default hypothetical FEC decoding buffer size is equal to 1Mbytes.

For the packet stream, the buffer occupancy level of the hypothetical FEC decoding buffer shall not exceed the value of the buf-size ~~MIME~~-parameter, when it is present in the SDP, or the default FEC decoding buffer size, when the buf-size ~~MIME~~-parameter is not present in the SDP. The output of the hypothetical FEC decoder shall comply with the RTP payload and decoding specifications of the media format.

The hypothetical FEC decoder operates as follows:

- 1) The hypothetical FEC decoding buffer is initially empty.
- 2) Each FEC source packet and FEC repair packet, ~~including its RTP header~~, starting from the first packet in transmission order, is inserted to a FEC source block at its transmission time. The FEC source block generation is done as specified in clause 8.2.2.51-6. The FEC source block resides in the hypothetical FEC decoding buffer.
- 3) When both the last FEC source packet and the last FEC repair packet of an FEC source block are transmitted, any elements of the FEC source block that are not original ~~RTP-UDP~~ packets (e.g. FEC repair packets and potential padding bytes) are removed from the hypothetical FEC decoding buffer.
- 4) Original ~~RTP-UDP~~ packets are not removed from the hypothetical FEC decoding buffer before the signaled initial buffering delay has expired. Then, the first original ~~RTP-UDP~~ packet in sequence number order is output and removed from the hypothetical FEC decoding buffer immediately. Each succeeding original ~~RTP-UDP~~ packet is output and removed when the following conditions are true:
 - i. The following time (in seconds) since the removal of the previous packet has elapsed:

$$8 * (\text{size of the previous original } \del{\text{RTP-UDP}} \text{ packet with including UDP/IP header overhead in bytes}) / (1000 * (\text{value of "b=AS" SDP attribute for the stream}))$$
 - ii. All the packets in the same FEC source block as the original ~~RTP-UDP~~ packet have been transmitted.

An MBMS client shall be capable of receiving a packet stream that complies with the hypothetical FEC decoder. Furthermore, in the case of RTP packets, when an MBMS client complies with the requirements for the media decoding of the packet stream, it shall be able to de-packetize and decode the packet stream and output decoded data at the correct rate specified by the RTP timestamps of the received packet stream.

8.2.2.12 FEC encoding procedures

~~NOTE: Here the exact procedures of the FEC scheme is to be defined. A reference to the detailed annex will be put when the FEC scheme is decided..~~

FEC encoding shall be performed using the MBMS forward error correction code defined in Annex B.

8.2.2.13 Signalling

The signalling for streaming FEC consists of several components:

- If several user services are bundled together they are indicated as a sequence of services in the User Service Bundle Description. See section 5.2.2.4

- A separate SDP describing the FEC repair stream and all the flow IDs. Referenced from the User Service Bundle Description. See section 5.2.2.6 and 8.2.2.13
- SDP protocol identifiers and attributes to indicate the usage of the source packet format, how the FEC payload ID is configured and other FEC parameters such as minimal buffering delay, for the RTP/RTCP streams. See section 8.2.2.12
- Security description extensions to indicate usage of FEC source packet format, and the FEC parameters. See section 5.2.2.3 and 8.2.2.12.

~~The two different RTP payload formats requires each a media type to identify them and define their respective set of parameters. The media types are:~~

~~-media type audio, video, or text/rtp-mbms-fec-repair~~

~~-media type audio, video, or text/rtp-mbms-fec-source~~

~~Their associated set of parameters are defined in annex C.~~

~~The use of the FEC Source Packet format defined in 8.2.2.4 is indicated in SDP, in the case of RTP and RTCP streams and in the service description in the case of MIKEY.~~

~~The FEC Repair packet stream is also signalled in SDP along with the definition of the UDP Flow IDs that will be used within the source block to distinguish different UDP flows. The user service description contains either a single service or several bundled services. All of the streaming delivery methods and security descriptions that are present in within the *bundleDescription* element must be considered when configuring the FEC operations. This includes RTP, RTCP and MIKEY flows. A receiver intending to perform FEC decoding to cover for packet losses shall receive all the flows that are indicated to be sent as FEC source packets, even if the flows are in a service currently not played out. A receiver intending to use FEC shall also receive the FEC repair stream as described by the FEC Repair Stream Description. The delivery method's session description, and the security description both carry the FEC source packet configuration information: FEC encoding ID, FEC instance ID, and FEC OTI information. The FEC repair packet stream is configured using the similar methods as for the source packets, with the addition of the Flow ID information and buffer delay parameter.~~

8.2.2.13a SDP for FEC source packet streams

~~To indicate the presence of the FEC layer between IP/UDP and, RTP or SRTP a SDP protocol identifier³ is used. Instead of the normal RTP/AVP and RTP/SAVP protocol identifiers, 'UDP/MBMS-FEC/RTP/AVP' and 'UDP/MBMS-FEC/RTP/SAVP' are defined respectively. Both these protocol identifiers shall use the FMT space rules that are used for RTP/AVP and RTP/SAVP respectively, i.e. payload types used in the RTP session is listed. The protocol identifiers are defined in Appendix C1.~~

~~The FEC parameters, FEC encoding ID, FEC instance ID and FEC-OTI-Extension information are is signalled using the mechanism defined in 8.3.1.9. The "a=FEC" SDP attribute shall be used to indicate the single definition that is used for each media block component.~~

~~For MIKEY messages the service protection description is used to indicate when FEC source packet shall be used, see section 5.2.2.3. The FEC parameter used is also defined in the service protection description. As all MIKEY packets from all user services arrive on the same port, the receiver must use the destination address to separate FEC protected packets from not FEC protected packets. This requires that all MIKEY packets sent to a specific destination address are either is FEC protected or not. There can't be any mixing of Note that it is not possible to mix protected and non-protected ing packets within a single stream as there is no header field mechanism to determine if whether they are protected or not.~~

8.2.2.14 Mapping the Media types to SDP SDP for FEC repair packet streams

~~The repair packet stream is indicated in SDP using a media block with the protocol identifier "UDP/MBMS-REPAIR". The media type shall be "application". The FEC parameters, FEC encoding ID, FEC instance ID, FEC-OTI-Extension information and repair parameters (min-buffer-time) is are signalled using the mechanisms defined in 8.3.1.9. The information carried in the MIME media type specification has a specific mapping to fields in the Session Description Protocol (SDP) [6], which is commonly used to describe RTP sessions. When SDP is used to specify sessions employing the FEC payload format, the mapping is as follows:~~

- o The Media type (application, audio, video, or text) goes in SDP "m=" as the media name.
- o The Media subtype (payload format name) goes in SDP "a=rtpmap" as the encoding name. The RTP clock rate in "a=rtpmap" SHALL be 10000 for audio/rtp-mbms-fec-repair, video/rtp-mbms-fec-repair, or text/rtp-mbms-fec-repair, and according to the rate parameter for audio/rtp-mbms-fec-source, video/rtp-mbms-fec-source, or text/rtp-mbms-fec-source.
- ~~— The protocol shall be set to 'UDP/MBMS-FEC/RTP/AVP' or 'UDP/MBMS-FEC/RTP/SAVP' to indicate the use of the FEC Source packet format with RTP or SRTP.~~
- o The FEID and FIID parameters are indicated using the "a=FEC-declaration" (see clause 7.3.2.8) and the FOTI parameter by the "a=FEC-OTI-extension" (see clause 8.3.1.9). To bind the correct FEC declaration and corresponding FEC OTI-extension to a payload type, an "a=fmtp" line parameter value pair "FEC-ref="<fec-ref>" shall be used. Where the <fec-ref> value is equal to the fec-ref number used in the SDP attributes FEC-declaration and FEC-OTI-extension.
- o Any remaining parameters go in the SDP "a=fmtp" attribute by copying them directly from the MIME media type string as a semicolon-separated list of parameter=value pairs.

~~These payload formats is only intended to be used in declarative SDP use cases and no offer/answer negotiation procedures is defined. A single FEC declaration shall be indicated using the "a=FEC" attribute.~~

The mapping of the FEC source block flow ID (see 8.2.2.5) to the destination IP address and UDP port are done using the SDP attribute "a=mbms-flowid" defined in section 8.3.1.10.

8.2.1.15 Signalling Example of SDP for FEC

This section contains a complete signalling example for a session using FEC with a Service description, a SDP for the streaming delivery method, a SDP for the FEC repair stream, and a security description.

The top element is the security description that

```
<?xml version="1.0" encoding="UTF-8"?>
<bundleDescription
  xmlns="www.example.com/3gppUserServiceDescription"
  xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"
  fecDescriptionURI="http://www.example.com/3gpp/mbms/session1-fec.sdp">
  <userServiceDescription
    userServiceId="urn:3gpp:0010120123hotdog">
    <deliveryMethod
      sessionDescriptionURI="http://www.example.com/3gpp/mbms/session1.sdp"
      protectionDescriptionURI="http://www.example.com/3gpp/mbms/sec-descript"/>
    </userServiceDescription>
  </bundleDescription>
```

The security description has the URI: http://www.example.com/3gpp/mbms/sec-descript

```
<?xml version="1.0" encoding="UTF-8"?>
<securityDescription
  xmlns="www.example.com/3gppSecurityDescription"
  xmlns:xs="http://www.w3.org/2001/XMLSchema-instance"
  confidentialityProtection="true"
  integrityProtection="true"
  uiccKeyManagement="true">
  <keyManagement
    waitTime="5"
    maxBackOff="10">
    <serverURI =http://register.operator.umts/ />
    <serverURI ="http:// register2.operator.umts/" />
  </keyManagement>
  <keyId identity="<someMSKIdA>" mediaFlow=FF1E:03AD::7F2E:172A:1E24:4002/>
  <keyId identity="<someMSKIdB>" mediaFlow=FF1E:03AD::7F2E:172A:1E24:4004/>
  <fecProtection
    fecEncodingId="130"
    fecInstanceId="0"
    fecOtiExtension="1SCxWEMNe397m24SwgyRhg==" />
</securityDescription>
```


An example of how ~~an~~the SDP <http://www.example.com/3gpp/mbms/session1.sdp> could look for a session containing two media streams that are FEC protected. In this example we have assumed an audiovisual stream, using 56 kbps for video and 12 kbps for audio. ~~In addition another 300 bits/second of RTCP packets from the source is used for the each of the sessions. We further assume that we send redundant packets for the video part at 6 kbps and redundant packets for the audio part at 3 kbps.~~ Hence, the total media session bandwidth is $56 + \underline{126} + \underline{0.3} + \underline{0.3} + 12 + 3 = \underline{77.68.6}$ kbps. ~~In addition another 1200 bits/second of RTCP packets from the source is used for the both sessions.~~

~~The FEC encoding symbols payloads does also declare that the minimal required buffering time for either of the streams are 2.6 seconds. As the value is defined as a limitation on the sender side, further buffering to handle jitter in the transmission is also likely to be required.~~

```
v=0
o=ghost 2890844526 2890842807 IN IP4 2001:210:1:2:240:96FF:FE25:8EC9192.168.10.10
s=3GPP MBMS Streaming SDP Example
i=Example of MBMS streaming SDP file
u=http://www.infoserver.example.com/ae600
e=ghost@mailserver.example.com
c=IN IP4 64 FF1E:03AD::7F2E:172A:1E24
224.1.2.3
t=3034423619 3042462419
b=AS:6277
b=TIAS: 60500
a=maxprate: 25
a=source-filter: incl IN IP6 * 2001:210:1:2:240:96FF:FE25:8EC9
a=FEC-declaration:0 encoding-id=130; instance-id=0
a=FEC-OTI-extension:0 1SCxWEMNe397m24SwgyRhg==

m=video 4002 UDP/MBMS-FEC/RTP/AVP 97-96-100
b=TIAS:5500062
b=RR:0
b=RS:6300
a=rtpmap:96 H263-2000/90000
a=fmtp:96 profile=3;level=10
a=framesize:96 176-144
a=FEC:0
a=maxprate:15
a=rtpmap: 97 rtp-mbms-fec-source/90000
a=fmtp:97 opt=96; FEID=129; FHD=12435; FOTI="1SCxWEMNe397m24SwgyRhg=="
a=rtpmap: 100 rtp-mbms-fec-repair/10000
a=fmtp:100 FEID=129; FHD=12435; FOTI="1SCxWEMNe397m24SwgyRhg=="; min-  
buffer-time=2600
m=audio 4004 UDP/MBMS-FEC/RTP/AVP 99-98 101
b=TIAS:1150045
b=RR:0
b=RS:6300
a=rtpmap:98 AMR/8000
a=fmtp:98 octet-align=1
a=FEC:0
a=maxprate:10
a=rtpmap: 99 rtp-mbms-fec-source/8000
a=fmtp: 99 opt=98; FEID=129; FHD=12435; FOTI="1SCxWEMNe397m24SwgyRhg=="
a=rtpmap: 101 rtp-mbms-fec-repair/10000
a=fmtp: 101 FEID=129; FHD=12435; FOTI="1SCxWEMNe397m24SwgyRhg=="; min-  
buffer-time=2600
```

[The FEC stream used to protect the above RTP sessions and a MIKEY key stream has the below SDP \(http://www.example.com/3gpp/mbms/session1-fec.sdp\):](http://www.example.com/3gpp/mbms/session1-fec.sdp)

```

v=0
o=ghost 2890844526 2890842807 IN IP6 2001:210:1:2:240:96FF:FE25:8EC9
s=3GPP MBMS Streaming FEC SDP Example
i=Example of MBMS streaming SDP file
u=http://www.infoserver.example.com/ae600
e=ghost@mailserver.example.com
c=IN IP6 FF1E:03AD::7F2E:172A:1E24
t=3034423619 3042462419
b=AS:15
a=FEC-declaration:0 encoding-id=131; instance-id=0
a=FEC-OTI-extension:0 1SCxWEMNe397m24SwgyRhg==
a=mbms-repair: 0 min-buffer-time=2600
a=source-filter: incl IN IP6 * 2001:210:1:2:240:96FF:FE25:8EC9
m=application 4006 UDP/MBMS-REPAIR *
b=AS:15
a=FEC:0
a=mbms-apid: 1=FF1E:03AD::7F2E:172A:1E24/4002, 2=FF1E:03AD::7F2E:172A:1E24/4003,
3=FF1E:03AD::7F2E:172A:1E24/4004, 4=FF1E:03AD::7F2E:172A:1E24/4005,
5=FF1E:03AD::7F2E:172A:1E24/2269

```

***** Next Change *****

8.3.1.8 FEC Parameters

The FEC encoding ID and instance ID is provided using the "a=FEC-declaration" attribute defined in Clause 7.3.2.8. Any OTI information for that FEC encoding ID and instance ID is provided with below defined FEC OTI attribute.

The FEC OTI attribute must be immediately preceded by the "a=FEC-declaration" attribute (and so can be session-level and media-level). The fec-ref maps the oti-extension to the FEC-declaration OTI it extends. The purpose of the oti-extension is to define FEC code specific OTI required for RTP receiver FEC payload configuration, exact contents are FEC code specific and need to be specified by each FEC code using this attribute.

The syntax for the attributes in ABNF [23] is:

```

sdp-fec-oti-extension-line = "a=FEC-OTI-extension:" fec-ref SP oti-extension CRLF
fec-ref = 1*3DIGIT (the SDP-internal identifier for the associated FEC-declaration).
oti-extension = base64
base64      = *base64-unit [base64-pad]
base64-unit = 4base64-char
base64-pad  = 2base64-char "=" / 3base64-char "="
base64-char = ALPHA / DIGIT / "+" / "/"

```

To provide the FEC repair packets with additional, non FEC specific parameters, a session and media level SDP attribute is defined.

```

sdp-fec-parameter-line = "a=mbms-repair: 0*1SP fec-ref SP parameter-list CRLF
parameter-list = parameter-spec *(1*SP parameter-spec)
parameter-spec = name "=" value;
name = 1*(ALPHA / DIGIT)
value = 1*(safe) ; safe defined in RFC 2327

```

Currently one FEC non code-specific parameter is defined:

min-buffer-time: This FEC buffering parameter specifies the minimum receiver buffer time (delay) needed to ensure that FEC repair has time to happen regardless of the FEC source block of the stream from which the reception starts. The value is in milliseconds and represents the wallclock time between the reception of the first FEC source or repair packet of a FEC source block, whichever is earlier in transmission order, and the wallclock time when media decoding can safely start.

The parameters name and value is defined in ABNF as follows:

Min-buffer-time-parameter-name = "min-buffer-time"

Min-buffer-time-parameter-value = 1*8DIGIT ;Wallclock time in milliseconds.

The FEC declaration and FEC OTI information utilized in a specific ~~RTP payload type~~ source or repair packet is indicated using the FEC-ref number in the a=~~fmtp~~fec lines as described in clause 8.2.4.2.4.12 and 8.2.2.13.

8.3.1.9 FEC Flow ID attribute

To indicate the mapping between destination IP address and UDP port number and FEC source block flow IDs, the "a=mbms-flowid" SDP attribute is defined. Each flowID that may be used in source block within the bundled sessions shall be included. It is a media level attribute that shall be present in any SDP media block using the "UDP/MBMS-REPAIR" protocol identifier.

The syntax for the attributes in ABNF [23] is:

Sdp-mbms-flowid-attr = "a=mbms-flowid:" *WSP flow-id-spec *(", " *WSP flow-id-spec) CRLF

flow-id-spec = flowID "=" address-spec "/" port-spec

address-spec = multicast-address ; As defined by RFC 3266

port-spec = 1*5DIGIT

8.3.2 SDP Example for Streaming Session

Here is a full example of SDP description describing a FLUTE session:

```
v=0
o=ghost 2890844526 2890842807 IN IP4 192.168.10.10
s=3GPP MBMS Streaming SDP Example
i=Example of MBMS streaming SDP file
u=http://www.infoserver.example.com/ae600
e=ghost@mailserver.example.com
c=IN IP6 FF1E:03AD::7F2E:172A:1E24
t=3034423619 3042462419
b=AS:77
a=mbms-mode:broadcast 1234
a=source-filter: incl IN IP6 * 2001:210:1:2:240:96FF:FE25:8EC9
a=FEC-declaration:0 encoding-id=130; instance-id=0
a=FEC-OTI-extension:0 1SCxWEMNe397m24SwgyRhg==
a=FEC-declaration:1 encoding-id=131; instance-id=2
a=FEC-OTI-extension:1 1SCxWEMNe397m24SwgyRhg==
m=video 4002 UDP/MBMS-FEC/RTP/AVP 97-96-100
b=TIAS:62000
b=RR:0
b=RS:600
a=maxprate:17
a=rtpmap:96 H264/90000
a=fmtp:96 profile-level-id=42A01E; packetization-mode=1; sprop-parameter-sets=Z0IACpZTBYmI,aMljiA==
a=rtpmap:97 rtp-mbms-fec-source/90000
```

```
a=fmtp:97 opt=96; FEC-ref=0
a=rtpmap: 100 rtp-mbms fec-source/10000
a=fmtp:100 FEC-ref=1; min-buffer-time=2600
m=audio 4004 UDP/MBMS-FEC/RTP/AVP 99-98-101
b=TIAS:15120
b=RR:0
b=RS:600
a=maxprate:10
a=rtpmap:98 AMR/8000
a=fmtp:98 octet-align=1
a=rtpmap: 99 rtp-mbms fec-source/8000
a=fmtp: 99 opt=98; FEC-ref=0 "
a=rtpmap: 101 rtp-mbms fec-source/10000
a=fmtp:101 FEC-ref=1; min-buffer-time=2600a=FEC:0
```

Annex <A> (normative): FLUTE Support Requirements

This clause provides a table representation of the requirement levels for different features in FLUTE. Table 1 includes requirements for an MBMS client and an MBMS server for FLUTE support as well as the requirements for a FLUTE client and a FLUTE server according to the FLUTE protocol [9]. The terms used in Table 1 are described underneath.

Table 1: Overview of the FLUTE support requirements in MBMS servers and clients

| | FLUTE Client support requirement as per [9]. | MBMS FLUTE Client support requirement as per present document | FLUTE Server use requirement as per [9]. | MBMS FLUTE Server use requirement as per present document |
|--|---|--|--|---|
| FLUTE Blocking Algorithm | Required | Required | Strongly recommended | Required |
| Symbol Encoding Algorithm | Compact No-Code algorithm required. Other FEC building blocks are undefined optional plug-ins. | Compact No-Code algorithm required. [TBD] MBMS Forward Error Correction required | Compact No-Code algorithm is the default option. Other FEC building blocks are undefined optional plug-ins. | Compact No-Code algorithm is the default option. MBMS Forward Error Correction. [TBD] |
| Congestion Control Building Block (CCBB) / Algorithm | Congestion Control building blocks undefined. | Single channel support required | Single channel without additional CCBB given for the controlled network scenario. | Single channel support required |
| Content Encoding for FDT Instances | Optional | Not applicable | Optional | Not applicable |
| A flag active (header) | Required | Required | Optional | Optional |
| B flag active (header) | Required | Required | Optional | Optional |
| T flag active and SCT field (header) | Optional | Optional | Optional | Optional |
| R flag active and ERT field (header) | Optional | Optional | Optional | Optional |
| Content-Location attribute (FDT) | Required | Required | Required | Required |
| TOI (FDT) | Required | Required | Required | Required |
| FDT Expires attribute (FDT) | Required | Required | Required | Required |
| Complete attribute (FDT) | Required | Required | Optional | Optional |
| FEC-OTI-Maximum-Source-Block-Length | Required | Required | Required | Required |
| FEC-OTI-Encoding-Symbol-Length | Required | Required | Required | Required |
| FEC-OTI-Max-Number-of-Encoding-Symbols. | Required | Required | Required | Required |
| FEC-OTI-FEC-Instance-ID | Required | [TBD] | Required | [TBD] |
| FEC-OTI-Scheme-Specific-Info | n/a | Required | n/a | Required if MBMS FEC used |

The following are descriptions of the above terms:

- Blocking algorithm: The blocking algorithms is used for the fragmentation of files. It calculates the source blocks from the source files.
- Symbol Encoding algorithm: The symbol encoding algorithm is used for the fragmentation of files. It calculates encoding symbols from source blocks for Compact No-Code FEC. It may also be used for other FEC schemes.
- Congestion Control Building Block: A building block used to limit congestion by using congestion feedback, rate regulation and receiver controls [17].

- Content Encoding for FDT Instances: FDT Instance may be content encoded for more efficient transport, e.g. using ZLIB.
- A flag: The Close Session flag for indicating the end of a session to the receiver in the ALC/LCT header.
- B flag: The Close Object flag is for indicating the end of an object to the receiver in the ALC/LCT header.
- T flag: The T flag is used to indicate the use of the optional “Sender Current Time (SCT)” field (when T=1) in the ALC/LCT header.
- R flag: The R flag is used to indicate the use of the optional “Expected Residual Time (ERT) field in the ALC/LCT header.
- Content Location attribute: This attribute provides a URI for the location where a certain piece of content (or file) being transmitted in a FLUTE session is located.
- Transport Object Identifier (TOI): The TOI uniquely identifies the object within the session from which the data in the packet was generated.
- FDT Expires attribute: Indicates to the receiver the time until which the information in the FDT is valid.
- Complete attribute: This may be used to signal that the given FDT Instance is the last FDT Instance to be expected on this file delivery session.
- FEC-OTI-Maximum-Source-Block-Length: This parameter indicates the maximum number of source symbols per source block.
- FEC-OTI-Encoding-Symbol-Length: This parameter indicates the length of the Encoding Symbol in bytes.
- FEC-OTI-Max-Number-of-Encoding-Symbols: This parameter indicates the maximum number of Encoding Symbols that can be generated for a source block.
- FEC-OTI-FEC-Instance-ID: This field is used to indicate the FEC Instance ID, if a FEC scheme is used.
- [FEC-OTI-Scheme-Specific-Info: Carries Object Transmission Information which is specific to the FEC scheme in use.](#)

***** Next Change *****

Annex B (normative): FEC encoder and decoder specification

~~NOTE:—This annex will provide the detailed encoder and decoder FEC specification when FEC code selection is complete. Only Raptor codes and/or 2D Reed-Solomon codes are considered.~~

[This Annex specifies the systematic Raptor forward error correction code and its application to MBMS \[7\]. Raptor is a fountain code, i.e., as many encoding symbols as needed can be generated by the encoder on-the-fly from the source symbols of a block. The decoder is able to recover the source block from any set of encoding symbols only slightly more in number than the number of source symbols.](#)

[The code described in this document is a Systematic code, that is, the original source symbols are sent unmodified from sender to receiver, as well as a number of repair symbols.](#)

B.1 Definitions, Symbols and abbreviations

B.1.1 Definitions

For the purposes of this Annex, the following terms and definitions apply.

Source block: a block of K source symbols which are considered together for Raptor encoding purposes.

Source symbol: the smallest unit of data used during the encoding process. All source symbols within a source block have the same size.

Encoding symbol: a symbol that is included in a data packet. The encoding symbols consist of the source symbols and the repair symbols. Repair symbols generated from a source block have the same size as the source symbols of that source block.

Systematic code: a code in which the source symbols are included as part of the encoding symbols sent for a source block.

Repair symbol: the encoding symbols sent for a source block that are not the source symbols. The repair symbols are generated based on the source symbols.

Intermediate symbols: symbols generated from the source symbols using an inverse encoding process. The repair symbols are then generated directly from the intermediate symbols. The encoding symbols do not include the intermediate symbols, i.e., intermediate symbols are not included in data packets.

Symbol: a unit of data. The size, in bytes, of a symbol is known as the symbol size.

Encoding symbol group: a group of encoding symbols that are sent together, i.e., within the same packet whose relationship to the source symbols can be derived from a single Encoding Symbol ID.

Encoding Symbol ID: information that defines the relationship between the symbols of an encoding symbol group and the source symbols.

Encoding packet: data packets that contain encoding symbols

Sub-block: a source block is sometime broken into sub-blocks, each of which is sufficiently small to be decoded in working memory. For a source block consisting of K source symbols, each sub-block consists of K sub-symbols, each symbol of the source block being composed of one sub-symbol from each sub-block.

Sub-symbol: part of a symbol. Each source symbol is composed of as many sub-symbols as there are sub-blocks in the source block.

Source packet: data packets that contain source symbols.

Repair packet: data packets that contain repair symbols.

B.1.2. Symbols

$i, j, x, h, a, b, d, v, m$ represent positive integers

$\text{ceil}(x)$ denotes the smallest positive integer which is greater than or equal to x

$\text{choose}(i, j)$ denotes the number of ways j objects can be chosen from among i objects without repetition

$\text{floor}(x)$ denotes the largest positive integer which is less than or equal to x

$i \% j$ denotes i modulo j

$X \wedge Y$ denotes, for equal-length bit strings X and Y , the bitwise exclusive-or of X and Y

A denote a symbol alignment parameter. Symbol and sub-symbol sizes are restricted to be multiples of A .

A^T denotes the transposed matrix of matrix A

| | |
|---------------------------------|--|
| A^{-1} | denotes the inverse matrix of matrix A |
| K | denotes the number of symbols in a single source block |
| K_{MAX} | denotes the maximum number of source symbols that can be in a single source block. Set to 8192. |
| L | denotes the number of pre-coding symbols for a single source block |
| S | denotes the number of LDPC symbols for a single source block |
| H | denotes the number of Half symbols for a single source block |
| C | denotes an array of intermediate symbols, $C[0], C[1], C[2], \dots, C[L-1]$ |
| C' | denotes an array of source symbols, $C'[0], C'[1], C'[2], \dots, C'[K-1]$ |
| X | a non-negative integer value |
| V_0, V_1 | two arrays of 4-byte integers, $V_0[0], V_0[1], \dots, V_0[255]$ and $V_1[0], V_1[1], \dots, V_1[255]$ |
| $\text{Rand}[X, i, m]$ | a pseudo-random number generator |
| $\text{Deg}[v]$ | a degree generator |
| $\text{LTEnc}[K, C, (d, a, b)]$ | a LT encoding symbol generator |
| $\text{Trip}[K, X]$ | a triple generator function |
| G | the number of symbols within an encoding symbol group |
| N | the number of sub-blocks within a source block |
| T | the symbol size in bytes. If the source block is partitioned into sub-blocks, then $T = T' \cdot N$. |
| T' | the sub-symbol size, in bytes. If the source block is not partitioned into sub-blocks then T' is not relevant. |
| F | the file size, for file download, in bytes |
| I | the sub-block size in bytes |
| P | for file download, the payload size of each packet, in bytes, that is used in the recommended derivation of the file download transport parameters. For streaming, the payload size of each repair packet, in bytes, that is used in the recommended derivation of the streaming transport parameters. |
| Q | $Q = 65521$, i.e., Q is the largest prime smaller than 2^{16} |
| Z | the number of source blocks, for file download |
| $J(K)$ | the systematic index associated with K |
| G | denotes any generator matrix |
| I_S | denotes the $S \times S$ identity matrix |
| $0_{S \times H}$ | denotes the $S \times H$ zero matrix |

B.1.3 Abbreviations

For the purposes of the present document, the following abbreviations apply:

| | |
|-------------|--|
| <u>ESI</u> | <u>Encoding Symbol ID</u> |
| <u>LDPC</u> | <u>Low Density Parity Check</u> |
| <u>LT</u> | <u>Luby Transform</u> |
| <u>SBN</u> | <u>Source Block Number</u> |
| <u>SBL</u> | <u>Source Block Length (in units of symbols)</u> |

B.2. Overview

The Raptor forward error correction code can be applied to both the MBMS file delivery and MBMS streaming applications described in the main body of this document. Raptor code aspects which are specific to each of these applications are discussed in Sections B.3 and B.4 of this document.

The principle component of the systematic Raptor code is the basic encoder described in Section B.5. First, it is described how to derive values for a set of intermediate symbols from the original source symbols such that knowledge of the intermediate symbols is sufficient to reconstruct the source symbols. Secondly, the encoder produces repair symbols which are each the exclusive OR of a number of the intermediate symbols. The encoding symbols are the combination of the source and repair symbols. The repair symbols are produced in such a way that the intermediate symbols and therefore also the source symbols can be recovered from any sufficiently large set of encoding symbols.

This document defines the systematic Raptor code encoder. A number of possible decoding algorithms are possible. An efficient decoding algorithm is provided in Section B.8.

The construction of the intermediate and repair symbols is based in part on a pseudo-random number generator described in Section B.5. This generator is based on a fixed set of 512 random numbers which must be available to both sender and receiver. These are provided in Section B.7.

Finally, the construction of the intermediate symbols from the source symbols is governed by a 'systematic index', values of which are provided in Section B.6 for source block sizes from 4 source symbols to $K_{MAX} = 8192$ source symbols.

B.3. File download

B.3.1. Source block construction

B.3.1.1. General

In order to apply the Raptor encoder to a source file, the file may be broken into $Z > 1$ blocks, known as *source blocks*. The Raptor encoder is applied independently to each source block. Each source block is identified by a unique integer Source Block Number (SBN), where the first source block has SBN zero, the second has SBN one, etc. Each source block is divided into a number, K , of *source symbols* of size T bytes each. Each source symbol is identified by a unique integer Encoding Symbol Identifier (ESI), where the first source symbol of a source block has ESI zero, the second has ESI one, etc.

Each source block with K source symbols is divided into $N > 1$ sub-blocks, which are small enough to be decoded in the working memory. Each sub-block is divided into K sub-symbols of size T' .

Note that the value of K is not necessarily the same for each source block of a file and the value of T' may not necessarily be the same for each sub-block of a source block. However, the symbol size T is the same for all source blocks of a file and the number of symbols, K is the same for every sub-block of a source block. Exact partitioning of the file into source blocks and sub-blocks is described in B.3.1.2 below.

Figure B.3.1.1.-1 shows an example source block placed into a two dimensional array, where each entry is a T' -byte sub-symbol, each row is a sub-block and each column is a source symbol. In this example, the value of T' is the same for every sub-block. The number shown in each sub-symbol entry indicates their original order within the source block. For example, the sub-symbol numbered K contains bytes $T' \cdot K$ through $T' \cdot (K+1) - 1$ of the source block. Then, source symbol i is the concatenation of the i th sub-symbol from each of the sub-blocks, which corresponds to the sub-symbols of the source block numbered $i, K+i, 2 \cdot K+i, \dots, (N-1) \cdot K+i$.

| | | | | | | | |
|-------------------------------|---------------------------------|---------------------------------|------------|------------|------------|------------|---------------------------------|
| <u>0</u> | <u>1</u> | <u>2</u> | <u>...</u> | <u>...</u> | <u>...</u> | <u>...</u> | <u>$K-1$</u> |
| <u>K</u> | <u>$K+1$</u> | <u>$K+2$</u> | <u>...</u> | <u>...</u> | <u>...</u> | <u>...</u> | <u>$2 \cdot K-1$</u> |
| <u>$2 \cdot K$</u> | <u>$2 \cdot K+1$</u> | <u>$2 \cdot K+2$</u> | <u>...</u> | <u>...</u> | <u>...</u> | <u>...</u> | <u>$3 \cdot K-1$</u> |

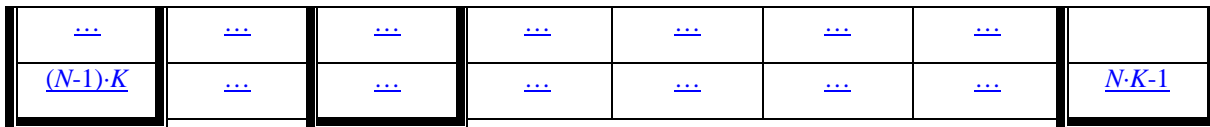


Figure B.3.1.1-1 – Source symbols from sub-symbols– the 3 highlighted columns show source symbols 0, 2 and $K-1$

B.3.1.2 Source block and sub-block partitioning

The construction of source blocks and sub-blocks is determined based on five input parameters, F , A , T , Z and N and a function Partition[]. The five input parameters are defined as follows:

- F the size of the file, in bytes
- A a symbol alignment parameter, in bytes
- T the symbol size, in bytes, which must be a multiple of A
- Z the number of source blocks
- N the number of sub-blocks in each source block

These parameters shall be set so that $\text{ceil}(\text{ceil}(F/T)/Z) \leq K_{MAX}$. Recommendations for derivation of these parameters are provided in Section B.3.4.

The function Partition[] takes a pair of integers (I, J) as input and derives four integers (I_L, I_S, J_L, J_S) as output. Specifically, the value of Partition[I, J] is a sequence of four integers (I_L, I_S, J_L, J_S) , where $I_L = \text{ceil}(I/J)$, $I_S = \text{floor}(I/J)$, $J_L = I - I_S \cdot J$ and $J_S = J - J_L$. Partition[] derives parameters for partitioning a block of size I into J approximately equal sized blocks. Specifically, J_L blocks of length I_L and J_S blocks of length I_S .

The source file shall be partitioned into source blocks and sub-blocks as follows:

Let,

$$K_L = \text{ceil}(F/T)$$

$$(K_L, K_S, Z_L, Z_S) = \text{Partition}[K_L, Z]$$

$$(T_L, T_S, N_L, N_S) = \text{Partition}[T/A, N]$$

Then, the file shall be partitioned into $Z = Z_L + Z_S$ contiguous source blocks, the first Z_L source blocks each having length $K_L \cdot T$ bytes and the remaining Z_S source blocks each having $K_S \cdot T$ bytes.

If $K_L \cdot T > F$ then for encoding purposes, the last symbol shall be padded at the end with $K_L \cdot T - F$ zero bytes.

Next, each source block shall be divided into $N = N_L + N_S$ contiguous sub-blocks, the first N_L sub-blocks each consisting of K contiguous sub-symbols of size of $T_L \cdot A$ and the remaining N_S sub-blocks each consisting of K contiguous sub-symbols of size of $T_S \cdot A$. The symbol alignment parameter A ensures that sub-symbols are always a multiple of A bytes.

Finally, the m th symbol of a source block consists of the concatenation of the m th sub-symbol from each of the N sub-blocks.

B.3.2. Encoding packet construction

B.3.2.1. General

Each encoding packet contains the following information:

- Source Block Number (SBN)
- Encoding Symbol ID (ESI)
- encoding symbol(s)

Each source block is encoded independently of the others. Source blocks are numbered consecutively from zero.

Encoding Symbol ID values from 0 to $K-1$ identify the source symbols. Encoding Symbol IDs from K onwards identify repair symbols.

B.3.2.2 Encoding packet construction

Each encoding packet either consists entirely of source symbols (source packet) or entirely of repair symbols (repair packet). A packet may contain any number of symbols from the same source block. In the case that the last symbol in the packet includes padding bytes added for FEC encoding purposes then these bytes need not be included in the packet. Otherwise, only whole symbols shall be included.

The Encoding Symbol ID, X , carried in each source packet is the Encoding Symbol ID of the first source symbol carried in that packet. The subsequent source symbols in the packet have Encoding Symbol IDs, $X+1$ to $X+G-1$, in sequential order, where G is the number of symbols in the packet.

Similarly, the Encoding Symbol ID, X , placed into a repair packet is the Encoding Symbol ID of the first repair symbol in the repair packet and the subsequent repair symbols in the packet have Encoding Symbol IDs $X+1$ to $X+G-1$ in sequential order, where G is the number of symbols in the packet.

Note that it is not necessary for the receiver to know the total number of repair packets. The G repair symbol triples $(d[0], a[0], b[0]), \dots, (d[G-1], a[G-1], b[G-1])$ for the repair symbols placed into a repair packet with ESI X are computed using the Triple generator defined in B.5.3.4 as follows:

For each $i = 0, \dots, G-1$

$$(d[i], a[i], b[i]) = \text{Trip}[K, X+i]$$

The G repair symbols to be placed in repair packet with ESI X are calculated based on the repair symbol triples as described in Section B.5.3 using the intermediate symbols C and the LT encoder $\text{LTenc}[K, C, (d[i], a[i], b[i])]$.

B.3.3. Transport

This section describes the information exchange between the Raptor encoder/decoder and any transport protocol making use of Raptor forward error correction for file delivery.

The Raptor encoder and decoder for file delivery require the following information from the transport protocol:

- The file size, F , in bytes
- The symbol alignment parameter, A
- The symbol size, T , in bytes, which must be a multiple of A
- The number of source blocks, Z
- The number of sub-blocks in each source block, N

The Raptor encoder for file delivery additionally requires:

- the file to be encoded, F bytes

The Raptor encoder supplies the transport protocol with encoding packet information consisting, for each packet, of:

- Source Block Number (SBN)
- Encoding Symbol ID (ESI)
- encoding symbol(s)

The transport protocol shall communicate this information transparently to the Raptor decoder.

Suitable transport protocols based on FLUTE/ALC and HTTP are defined in this specification.

B.3.4. Recommended Example Parameters (informative)

B.3.4.1 Parameter derivation algorithm

This section provides recommendations for the derivation of the four transport parameters, A , T , Z and N . This recommendation is based on the following input parameters:

- F the file size, in bytes
- W a target on the sub-block size, in bytes
- P the maximum packet payload size, in bytes, which is assumed to be a multiple of A
- A the symbol alignment factor, in bytes
- K_{MAX} the maximum number of source symbols per source block.
- K_{MIN} a minimum target on the number of symbols per source block
- G_{MAX} a maximum target number of symbols per packet

Based on the above inputs, the transport parameters T , Z and N are calculated as follows:

Let,

$$G = \min\{\text{ceil}(P \cdot K_{MIN}/F), P/A, G_{MAX}\} \quad \text{- the approximate number of symbols per packet}$$

$$T = \text{floor}(P/(A \cdot G)) \cdot A$$

$$K_t = \text{ceil}(F/T) \quad \text{- the total number of symbols in the file}$$

$$Z = \text{ceil}(K_t/K_{MAX})$$

$$N = \min\{\text{ceil}(\text{ceil}(K_t/Z) \cdot T/W), T/A\}$$

The values of G and N derived above should be considered as lower bounds. It may be advantageous to increase these values, for example to the nearest power of two. In particular, the above algorithm does not guarantee that the symbol size, T , divides the maximum packet size, P , and so it may not be possible to use the packets of size exactly P . If, instead, G is chosen to be a value which divides P/A , then the symbol size, T , will be a divisor of P and packets of size P can be used.

Recommended settings for the input parameters, W , A , K_{MIN} and G_{MAX} are as follows:

$$W = 256 \text{ KB} \quad A = 4 \quad K_{MIN} = 1024 \quad G_{MAX} = 10$$

B.3.4.2 Examples

The above algorithm leads to transport parameters as shown in Table B.3.4.2-1 below, assuming the recommended values for W , A , K_{MIN} and G_{MAX} and $P = 512$:

| <u>File size F</u> | <u>G</u> | <u>Symbol size T</u> | <u>$G \cdot T$</u> | <u>K_t</u> | <u>Source blocks Z</u> | <u>Sub-blocks N</u> | <u>K_L</u> | <u>K_S</u> | <u>T_L/A</u> | <u>T_S/A</u> |
|---------------------------------|-----------------------|-----------------------------------|-------------------------------|-------------------------|-------------------------------------|----------------------------------|-------------------------|-------------------------|---------------------------|---------------------------|
| 100 KB | 6 | 84 | 504 | 1,220 | 1 | 1 | 1,220 | 1,220 | N/A | N/A |
| 100 KB | 8 | 64 | 512 | 1,600 | 1 | 1 | 1,600 | 1,600 | N/A | N/A |
| 300 KB | 2 | 256 | 512 | 1,200 | 1 | 2 | 1,200 | 1,200 | 128 | 128 |
| 1,000 KB | 1 | 512 | 512 | 2,000 | 1 | 5 | 2,000 | 2,000 | 104 | 100 |

| | | | | | | | | | | |
|------------------|----------|------------|------------|---------------|----------|-----------|--------------|--------------|-----------|-----------|
| <u>3,000 KB</u> | <u>1</u> | <u>512</u> | <u>512</u> | <u>6,000</u> | <u>1</u> | <u>12</u> | <u>6,000</u> | <u>6,000</u> | <u>44</u> | <u>40</u> |
| <u>10,000 KB</u> | <u>1</u> | <u>512</u> | <u>512</u> | <u>20,000</u> | <u>3</u> | <u>14</u> | <u>6,666</u> | <u>6,667</u> | <u>40</u> | <u>36</u> |

Table B.3.4.2-1

B.4. Streaming

B.4.1. Source block construction

A source block is constructed by the transport protocol, for example as defined in this document, making use of the Systematic Raptor Forward Error Correction code. The symbol size, T , to be used for source block construction and the repair symbol construction are provided by the transport protocol. The parameter T shall be set so that the number of source symbols in any source block is at most K_{MAX} .

Recommended parameters are presented in section B.4.4.

B.4.2. Encoding packet construction

As described in B.4.3., each repair packet contains the following information:

- Source Block Number (SBN)
- Encoding Symbol ID (ESI)
- Source Block Length (SBL)
- repair symbol(s)

The number of repair symbols contained within a repair packet is computed from the packet length. The ESI values placed into the repair packets and the repair symbol triples used to generate the repair symbols are computed as described in Section B.3.2.2..

B.4.3. Transport

This section describes the information exchange between the Raptor encoder/decoder and any transport protocol making use of Raptor forward error correction for streaming.

The Raptor encoder for streaming requires the following information from the transport protocol for each source block:

- The symbol size, T , in bytes
- The number of symbols in the source block, K
- The Source Block Number (SBN)
- The source symbols to be encoded, $K \cdot T$ bytes

The Raptor encoder supplies the transport protocol with encoding packet information consisting, for each repair packet, of:

- Source Block Number (SBN)
- Encoding Symbol ID (ESI)
- Source Block Length (SBL)
- repair symbol(s)

The transport protocol shall communicate this information transparently to the Raptor decoder.

A suitable transport protocol is defined in this specification.

B.4.4. Recommended Example Parameters (informative)

B.4.4.1 Parameter derivation algorithm

This section provides recommendations for the derivation of the transport parameter T . This recommendation is based on the following input parameters:

- B the maximum source block size, in bytes
- P the maximum repair packet payload size, in bytes, which is a multiple of A
- A the symbol alignment factor, in bytes
- K_{MAX} the maximum number of source symbols per source block.
- K_{MIN} a minimum target on the number of symbols per source block
- G_{MAX} a maximum target number of symbols per repair packet

A requirement on these inputs is that $\text{ceil}(B/P) \leq K_{MAX}$. Based on the above inputs, the transport parameter T is calculated as follows:

Let,

$$\underline{G = \min\{\text{ceil}(P \cdot K_{MIN}/B), P/A, G_{MAX}\}} \quad \text{- the approximate number of symbols per packet}$$

$$\underline{T = \text{floor}(P/(A \cdot G)) \cdot A}$$

The value of T derived above should be considered as a guide to the actual value of T used. It may be advantageous to ensure that T divides into P , or it may be advantageous to set the value of T smaller to minimize wastage when full size repair symbols are used to recover partial source symbols at the end of lost source packets (as long as the maximum number of source symbols in a source block does not exceed K_{MAX}). Furthermore, the choice of T may depend on the source packet size distribution, e.g., if all source packets are the same size then it is advantageous to choose T so that the actual payload size of a repair packet P' , where P' is a multiple of T , is equal to (or as few bytes as possible larger than) the number of bytes each source packet occupies in the source block.

Recommended settings for the input parameters, A , K_{MIN} and G_{MAX} are as follows:

$$\underline{A = 4 \quad K_{MIN} = 1024 \quad G_{MAX} = 10}$$

B.4.4.2 Examples

The above algorithm leads to transport parameters as shown in Table B.4.4.2-1 below, assuming the recommended values for A , K_{MIN} and G_{MAX} and $P = 512$:

| <u>Max source block size B</u> | <u>G</u> | <u>Symbol size T</u> | <u>$G \cdot T$</u> |
|---|-----------------------|-----------------------------------|-------------------------------|
| <u>40 KB</u> | <u>10</u> | <u>48</u> | <u>480</u> |
| <u>160 KB</u> | <u>4</u> | <u>128</u> | <u>512</u> |
| <u>640 KB</u> | <u>1</u> | <u>512</u> | <u>512</u> |

Table B.4.4.2-1

B.5. Systematic Raptor encoder

B.5.1. Encoding overview

The systematic Raptor encoder is used to generate *repair symbols* from a source block that consists of K *source symbols*.

Symbols are the fundamental data units of the encoding and decoding process. For each source block (sub-block) all symbols (sub-symbols) are the same size. The atomic operation performed on symbols (sub-symbols) for both encoding and decoding is the exclusive-or operation.

Let $C'[0], \dots, C'[K-1]$ denote the K source symbols.

Let $C[0], \dots, C[L-1]$ denote L intermediate symbols.

The first step of encoding is to generate a number, $L > K$, of intermediate symbols from the K source symbols. In this step, K source triples $(d[0], a[0], b[0]), \dots, (d[K-1], a[K-1], b[K-1])$ are generated using the Trip[] generator as described in Section B.5.4.4. The K source triples are associated with the K source symbols and are then used to determine the L intermediate symbols $C[0], \dots, C[L-1]$ from the source symbols using an inverse encoding process. This process can be realized by a Raptor decoding process.

Certain “pre-coding relationships” must hold within the L intermediate symbols. Section B.5.2 describes these relationships and how the intermediate symbols are generated from the source symbols.

Once the intermediate symbols have been generated, repair symbols are produced and one or more repair symbols are placed as a group into a single data packet. Each repair symbol group is associated with an Encoding Symbol ID (ESI) and a number, G , of encoding symbols. The ESI is used to generate a triple of three integers, (d, a, b) for each repair symbol, again using the Trip[] generator as described in Section B.5.4.4. This is done as described in Sections B.3 and B.4 using the generators described in Section B.5.4. Then, each (d,a,b) -triple is used to generate the corresponding repair symbol from the intermediate symbols using the $\text{LTEnc}[K, C[0], \dots, C[L-1], (d,a,b)]$ generator described in Section B.5.4.3.

B.5.2. First encoding step: Intermediate Symbol Generation

B.5.2.1 General

The first encoding step is a pre-coding step to generate the L intermediate symbols $C[0], \dots, C[L-1]$ from the source symbols $C'[0], \dots, C'[K-1]$. The intermediate symbols are uniquely defined by two sets of constraints:

1. The intermediate symbols are related to the source symbols by a set of *source symbol triples*. The generation of the source symbol triples is defined in Section B.5.2.2 using the the Trip[] generator as described in Section B.5.4.4.
2. A set of pre-coding relationships hold within the intermediate symbols themselves. These are defined in Section B.5.2.3.

The generation of the L intermediate symbols is then defined in Section 5.2.4.

B.5.2.2 Source symbol triples

Each of the K source symbols is associated with a triple $(d[i], a[i], b[i])$ for $0 \leq i < K$. The source symbol triples are determined using the Triple generator defined in Section B.5.4.4 as:

For each $i, 0 \leq i < K$

$$(d[i], a[i], b[i]) = \text{Trip}[K, i]$$

B.5.2.3 Pre-coding relationships

The pre-coding relationships amongst the L intermediate symbols are defined by expressing the last $L-K$ intermediate symbols in terms of the first K intermediate symbols.

The last $L-K$ intermediate symbols $C[K], \dots, C[L-1]$ consist of S LDPC symbols and H Half symbols. The values of S and H are determined from K as described below. Then $L = K + S + H$.

Let

X be the smallest positive integer such that $X \cdot (X-1) = 2 \cdot K$.

S be the smallest prime integer such that $S \geq \text{ceil}(0.01 \cdot K) + X$

H be the smallest integer such that $\text{choose}(H, \text{ceil}(H/2)) > K + S$

$H' = \text{ceil}(H/2) \cdot L = K + S + H$

$C[0], \dots, C[K-1]$ denote the first K intermediate symbols

$C[K], \dots, C[K+S-1]$ denote the S LDPC symbols, initialised to zero

$C[K+S], \dots, C[L-1]$ denote the H Half symbols, initialised to zero

The S LDPC symbols are defined to be the values of $C[K], \dots, C[K+S-1]$ at the end of the following process:

For $i = 0, \dots, K-1$ do

$a = 1 + (\text{floor}(i/S) \% (S-1))$

$b = i \% S$

$C[K+b] = C[K+b] \wedge C[i]$

$b = (b+a) \% S$

$C[K+b] = C[K+b] \wedge C[i]$

$b = (b+a) \% S$

$C[K+b] = C[K+b] \wedge C[i]$

The H Half symbols are defined as follows:

Let

$g[i] = i \wedge (\text{floor}(i/2))$ for all positive integers i

Note: $g[i]$ is the Gray sequence, in which each element differs from the previous one in a single bit position

$g[j,k]$ denote the j^{th} element, $j=0, 1, 2, \dots$, of the subsequence of $g[i]$ whose elements have exactly k non-zero bits in their binary representation

Then, the Half symbols are defined as the values of $C[K+S], \dots, C[L-1]$ after the following process:

For $h = 0, \dots, H-1$ do

For $j = 0, \dots, K+S-1$ do

If bit h of $g[j, H']$ is equal to 1 then $C[h+K+S] = C[h+K+S] \wedge C[j]$.

B.5.2.4 Intermediate symbols

B.5.2.4.1 Definition

Given the K source symbols $C'[0], C'[1], \dots, C'[K-1]$ the L intermediate symbols $C[0], C[1], \dots, C[L-1]$ are the uniquely defined symbol values that satisfy the following conditions:

1. The K source symbols $C'[0], C'[1], \dots, C'[K-1]$ satisfy the K constraints

$$C'[i] \equiv \text{LTEnc}[K, (C[0], \dots, C[L-1]), (d[i], a[i], b[i])], \text{ for all } i, 0 \leq i < K.$$

2. The L intermediate symbols $C[0], C[1], \dots, C[L-1]$ satisfy the pre-coding relationships defined in B.5.2.3.

B.5.2.4.2 Example method for calculation of intermediate symbols (informative)

This subsection describes a possible method for calculation of the L intermediate symbols $C[0], C[1], \dots, C[L-1]$ satisfying the constraints in B.5.2.4.1

The generator matrix \mathbf{G} for a code which generates N output symbols from K input symbols is an $N \times K$ matrix over $GF(2)$, where each row corresponds to one of the output symbols and each column to one of the input symbols and where the i^{th} output symbol is equal to the sum of those input symbols whose column contains a non-zero entry in row i .

Then, the L intermediate symbols can be calculated as follows:

Let

\mathbf{C} denote the column vector of the L intermediate symbols, $C[0], C[1], \dots, C[L-1]$.

\mathbf{D} denote the column vector consisting of $S+H$ zero symbols followed by the K source symbols $C'[0], C'[1], \dots, C'[K-1]$

Then the above constraints define an $L \times L$ matrix over $GF(2)$, \mathbf{A} , such that:

$$\mathbf{A} \cdot \mathbf{C} = \mathbf{D}$$

The matrix \mathbf{A} can be constructed as follows:

Let:

\mathbf{G}_{LDPC} be the $S \times K$ generator matrix of the LDPC symbols. So,

$$\mathbf{G}_{\text{LDPC}} \cdot (C[0], \dots, C[K-1])^T = (C[K], \dots, C[K+S-1])^T$$

\mathbf{G}_{Half} be the $H \times (K+S)$ generator matrix of the Half symbols. So,

$$\mathbf{G}_{\text{Half}} \cdot (C[0], \dots, C[S+K-1])^T = (C[K+S], \dots, C[K+S+H-1])^T$$

\mathbf{I}_S be the $S \times S$ identity matrix

\mathbf{I}_H be the $H \times H$ identity matrix

$\mathbf{0}_{S \times H}$ be the $S \times H$ zero matrix

\mathbf{G}_{LT} be the $K \times L$ generator matrix of the encoding symbols generated by the LT Encoder.

So,

$$\mathbf{G}_{\text{LT}} \cdot (C[0], \dots, C[L-1])^T = (C'[0], C'[1], \dots, C'[K-1])^T$$

i.e. $\mathbf{G}_{\text{LT},i} = 1$ if and only if $C[i]$ is included in the symbols which are XORed to produce $\text{LTEnc}[K, (C[0], \dots, C[L-1]), (d[i], a[i], b[i])]$.

Then:

The first S rows of \mathbf{A} are equal to $\mathbf{G}_{\text{LDPC}} \mid \mathbf{I}_S \mid \mathbf{0}_{S \times H}$.

The next H rows of \mathbf{A} are equal to $\mathbf{G}_{\text{Half}} \parallel \mathbf{I}_H$.

The remaining K rows of \mathbf{A} are equal to \mathbf{G}_{LT} .

The matrix \mathbf{A} is depicted in the figure below:

| | | | |
|-----|----------------------------|----------------|---------------------------|
| | K | S | H |
| S | \mathbf{G}_{LDPC} | \mathbf{I}_S | $\mathbf{Z}_{S \times H}$ |
| H | \mathbf{G}_{Half} | | \mathbf{I}_H |
| K | \mathbf{G}_{LT} | | |

Figure B.5.2.5.2-1: The matrix \mathbf{A}

The intermediate symbols can then be calculated as:

$$\mathbf{C} = \mathbf{A}^{-1} \cdot \mathbf{D}$$

The source triples are generated such that for any K matrix \mathbf{A} has full rank and is therefore invertible. This calculation can be realized by applying a Raptor decoding process to the K source symbols $C'[0], C'[1], \dots, C'[K-1]$ to produce the L intermediate symbols $C[0], C[1], \dots, C[L-1]$.

To efficiently generate the intermediate symbols from the source symbols, it is recommended that an efficient decoder implementation such as that described in Section B.8 be used. The source symbol triples are designed to facilitate efficient decoding of the source symbols using that algorithm.

B.5.3. Second encoding step: LT encoding

In the second encoding step, the repair symbol with ESI X is generated by applying the generator $\text{LTEnc}[K, (C[0], C[1], \dots, C[L-1]), (d, a, b)]$ defined in Section B.5.4 to the L intermediate symbols $C[0], C[1], \dots, C[L-1]$ using the triple $(d, a, b) = \text{Trip}[K, X]$ generated according to Sections B.3.2.2 and B.4.2.

B.5.4. Generators

B.5.4.1 Random Generator

The random number generator $\text{Rand}[X, i, m]$ is defined as follows, where X is a non-negative integer, i is a non-negative integer and m is a positive integer and the value produced is an integer between 0 and $m-1$. Let V_0 and V_1 be arrays of 256 entries each, where each entry is a 4-byte unsigned integer. These arrays are provided in Section B.7.

Then,

$$\text{Rand}[X, i, m] = (V_0[(X+i) \% 256] \wedge V_1[(\text{floor}(X/256)+i) \% 256]) \% m$$

B.5.4.2 Degree Generator

The degree generator $\text{Deg}[v]$ is defined as follows, where v is an integer that is at least 0 and less than $2^{20} = 1048576$.

In Table 7.3.2-1, find the index j such that $f[j-1] \leq v < f[j]$

$$\text{Deg}[v] = d[j]$$

| <u>Index j</u> | <u>$f[j]$</u> | <u>$d[j]$</u> |
|-----------------------------|--------------------------|--------------------------|
| <u>0</u> | <u>0</u> | <u>--</u> |
| <u>1</u> | <u>10241</u> | <u>1</u> |
| <u>2</u> | <u>491582</u> | <u>2</u> |
| <u>3</u> | <u>712794</u> | <u>3</u> |
| <u>4</u> | <u>831695</u> | <u>4</u> |
| <u>5</u> | <u>948446</u> | <u>10</u> |
| <u>6</u> | <u>1032189</u> | <u>11</u> |
| <u>7</u> | <u>1048576</u> | <u>40</u> |

Table 7.3.2-1 – Defines the degree distribution for encoding symbols

B.5.4.3 LT Encoding Symbol Generator

The encoding symbol generator $LTEnc[K, (C[0], C[1], \dots, C[L-1]), (d, a, b)]$ takes the following inputs:

K is the number of source symbols (or sub-symbols) for the source block (sub-block). Let L be derived from K as described in Section B.5.2, and let L' be the smallest prime integer greater than or equal to L .

$(C[0], C[1], \dots, C[L-1])$ is the array of L intermediate symbols (sub-symbols) generated as described in Section B.5.2

(d, a, b) is a source triple determined using the Triple generator defined in Section B.5.3.4, whereby

d is an integer denoting an encoding symbol degree

a is an integer between 1 and $L'-1$ inclusive

b is an integer between 0 and $L'-1$ inclusive

The encoding symbol generator produces a single encoding symbol as output, according to the following algorithm:

While $(b \geq L)$ do $b = (b + a) \% L'$

$LTEnc[K, (C[0], C[1], \dots, C[L-1]), (d, a, b)] = C[b]$.

For $j = 1, \dots, \min(d-1, L-1)$ do

$b = (b + a) \% L'$

While $(b \geq L)$ do $b = (b + a) \% L'$

$LTEnc[K, (C[0], C[1], \dots, C[L-1]), (d, a, b)] = LTEnc[K, (C[0], C[1], \dots, C[L-1]), (d, a, b)] \wedge C[b]$

B.5.4.4 Triple generator

The triple generator $Trip[K, X]$ takes the following inputs:

K The number of source symbols

X An encoding symbol ID

Let

L be determined from K as described in Section B.5.2

L' be the smallest prime that is greater than or equal to L

$Q = 65521$, the largest prime smaller than 2^{16} .

$J(K)$ be the systematic index associated with K , as defined in Section B.7

The output of the triple generator is a triples, (d, a, b) determined as follows:

B.8.2 Decoding a source block

B.8.2.1 General

It is assumed that the decoder knows the structure of the source block it is to decode, including the symbol size, T , and the number K of symbols in the source block.

From the algorithms described in Sections B.5, the Raptor decoder can calculate the total number $L = K+S+H$ of pre-coding symbols and determine how they were generated from the source block to be decoded. In this description it is assumed that the received encoding symbols for the source block to be decoded are passed to the decoder. Furthermore, for each such encoding symbol it is assumed that the number and set of intermediate symbols whose exclusive-or is equal to the encoding symbol is passed to the decoder. In the case of source symbols, the source symbol triples described in Section B.5.2.2 indicate the number and set of intermediate symbols which sum to give each source symbol.

Let $N \geq K$ be the number of received encoding symbols for a source block and let $M = S+H+N$. The following M by L bit matrix \mathbf{A} can be derived from the information passed to the decoder for the source block to be decoded. Let \mathbf{C} be the column vector of the L intermediate symbols, and let \mathbf{D} be the column vector of M symbols with values known to the receiver, where the first $S+H$ of the M symbols are zero-valued symbols that correspond to LDPC and Half symbols (these are check symbols for the LDPC and Half symbols, and not the LDPC and Half symbols themselves), and the remaining N of the M symbols are the received encoding symbols for the source block. Then, \mathbf{A} is the bit matrix that satisfies $\mathbf{A} \cdot \mathbf{C} = \mathbf{D}$, where here \cdot denotes matrix multiplication over GF[2]. In particular, $A[i,j] = 1$ if the intermediate symbol corresponding to index j is exclusive-ORed into the LDPC, Half or encoding symbol corresponding to index i in the encoding, or if index i corresponds to a LDPC or Half symbol and index j corresponds to the same LDPC or Half symbol. For all other i and j , $A[i,j] = 0$.

Decoding a source block is equivalent to decoding \mathbf{C} from known \mathbf{A} and \mathbf{D} . It is clear that \mathbf{C} can be decoded if and only if the rank of \mathbf{A} over GF[2] is L . Once \mathbf{C} has been decoded, missing source symbols can be obtained by using the source symbol triples to determine the number and set of intermediate symbols which must be exclusive-ORed to obtain each missing source symbol.

The first step in decoding \mathbf{C} is to form a decoding schedule. In this step \mathbf{A} is converted, using Gaussian elimination (using row operations and row and column reorderings) and after discarding $M - L$ rows, into the L by L identity matrix. The decoding schedule consists of the sequence of row operations and row and column re-orderings during the Gaussian elimination process, and only depends on \mathbf{A} and not on \mathbf{D} . The decoding of \mathbf{C} from \mathbf{D} can take place concurrently with the forming of the decoding schedule, or the decoding can take place afterwards based on the decoding schedule.

The correspondence between the decoding schedule and the decoding of \mathbf{C} is as follows. Let $c[0] = 0, c[1] = 1, \dots, c[L-1] = L-1$ and $d[0] = 0, d[1] = 1, \dots, d[M-1] = M-1$ initially.

- Each time row i of \mathbf{A} is exclusive-ORed into row i' in the decoding schedule then in the decoding process symbol $D[d[i]]$ is exclusive-ORed into symbol $D[d[i']]$.
- Each time row i is exchanged with row i' in the decoding schedule then in the decoding process the value of $d[i]$ is exchanged with the value of $d[i']$.
- Each time column j is exchanged with column j' in the decoding schedule then in the decoding process the value of $c[j]$ is exchanged with the value of $c[j']$.

From this correspondence it is clear that the total number of exclusive-ORs of symbols in the decoding of the source block is the number of row operations (not exchanges) in the Gaussian elimination. Since \mathbf{A} is the L by L identity matrix after the Gaussian elimination and after discarding the last $M - L$ rows, it is clear at the end of successful decoding that the L symbols $D[d[0]], D[d[1]], \dots, D[d[L-1]]$ are the values of the L symbols $C[c[0]], C[c[1]], \dots, C[c[L-1]]$.

The order in which Gaussian elimination is performed to form the decoding schedule has no bearing on whether or not the decoding is successful. However, the speed of the decoding depends heavily on the order in which Gaussian elimination is performed. (Furthermore, maintaining a sparse representation of \mathbf{A} is crucial, although this is not described here). The remainder of this section describes an order in which Gaussian elimination could be performed that is relatively efficient.

B.8.2.2 First Phase

The first phase of the Gaussian elimination the matrix \mathbf{A} is conceptually partitioned into submatrices. The submatrix sizes are parameterized by non-negative integers i and u which are initialized to 0. The submatrices of \mathbf{A} are:

- (1) The submatrix \mathbf{I} defined by the intersection of the first i rows and first i columns. This is the identity matrix at the end of each step in the phase.
- (2) The submatrix defined by the intersection of the first i rows and all but the first i columns and last u columns. All entries of this submatrix are zero.
- (3) The submatrix defined by the intersection of the first i columns and all but the first i rows. All entries of this submatrix are zero.
- (4) The submatrix \mathbf{U} defined by the intersection of all the rows and the last u columns.
- (5) The submatrix \mathbf{V} formed by the intersection of all but the first i columns and the last u columns and all but the first i rows.

Figure A.2.2-1 illustrates the submatrices of \mathbf{A} . At the beginning of the first phase $\mathbf{V} = \mathbf{A}$. In each step, a row of \mathbf{A} is chosen.

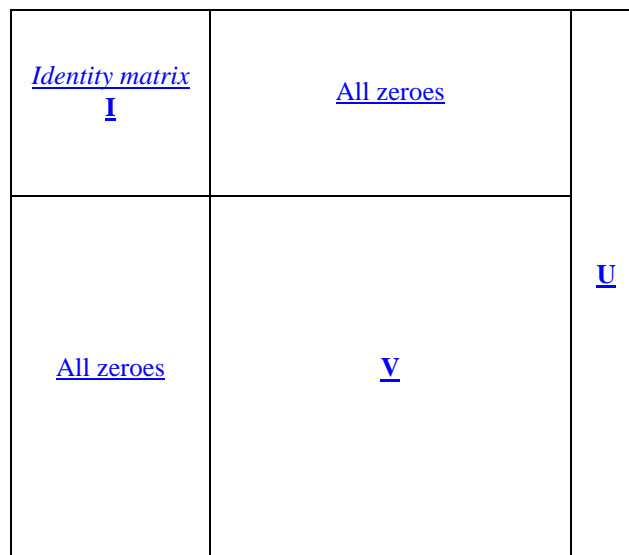


Figure A.2.2-1 – Submatrices of \mathbf{A} in the first phase

The following graph defined by the structure of \mathbf{V} is used in determining which row of \mathbf{A} is chosen. The columns that intersect \mathbf{V} are the nodes in the graph, and the rows that have exactly 2 ones in \mathbf{V} are the edges of the graph that connect the two columns (nodes) in the positions of the two ones. A component in this graph is a maximal set of nodes (columns) and edges (rows) such that there is a path between each pair of nodes/edges in the graph. The size of a component is the number of nodes (columns) in the component.

There are at most L steps in the first phase. The phase ends successfully when $i + u = L$, i.e., when \mathbf{V} and the all zeroes submatrix above \mathbf{V} have disappeared and \mathbf{A} consists of \mathbf{I} , the all zeroes submatrix below \mathbf{I} , and \mathbf{U} . The phase ends unsuccessfully in decoding failure if at some step before \mathbf{V} disappears there is no non-zero row in \mathbf{V} to choose in that step. In each step, a row of \mathbf{A} is chosen as follows:

- If all entries of \mathbf{V} are zero then no row is chosen and decoding fails.
- Let r be the minimum integer such that at least one row of \mathbf{A} has exactly r ones in \mathbf{V} .
- If $r \neq 2$ then choose a row with exactly r ones in \mathbf{V} with minimum original degree among all such rows.
- If $r = 2$ then choose any row with exactly 2 ones in \mathbf{V} that is part of a maximum size component in the graph defined by \mathbf{V} .

After the row is chosen in this step the first row of \mathbf{A} that intersects \mathbf{V} is exchanged with the chosen row so that the chosen row is the first row that intersects \mathbf{V} . The columns of \mathbf{A} among those that intersect \mathbf{V} are reordered so that one of the r ones in the chosen row appears in the first column of \mathbf{V} and so that the remaining $r-1$ ones appear in the last columns of \mathbf{V} . Then, the chosen row is exclusive-ORed into all the other rows of \mathbf{A} below the chosen row that have a one in the first column of \mathbf{V} . Finally, i is incremented by 1 and u is incremented by $r-1$, which completes the step.

B.8.2.3 Second Phase

The submatrix \mathbf{U} is further partitioned into the first i rows, $\mathbf{U}_{\text{upper}}$, and the remaining $M-i$ rows, $\mathbf{U}_{\text{lower}}$. Gaussian elimination is performed in the second phase on $\mathbf{U}_{\text{lower}}$ to either determine that its rank is less than u (decoding failure) or to convert it into a matrix where the first u rows is the identity matrix (success of the second phase). Call this u by u identity matrix \mathbf{I}_u . The $M-L$ rows of \mathbf{A} that intersect $\mathbf{U}_{\text{lower}} - \mathbf{I}_u$ are discarded. After this phase \mathbf{A} has L rows and L columns.

B.8.2.4 Third Phase

After the second phase the only portion of \mathbf{A} which needs to be zeroed out to finish converting \mathbf{A} into the L by L identity matrix is $\mathbf{U}_{\text{upper}}$. The number of rows i of the submatrix $\mathbf{U}_{\text{upper}}$ is generally much larger than the number of columns u of $\mathbf{U}_{\text{upper}}$. To zero out $\mathbf{U}_{\text{upper}}$ efficiently, the following precomputation matrix \mathbf{U}' is computed based on \mathbf{I}_u in the third phase and then \mathbf{U}' is used in the fourth phase to zero out $\mathbf{U}_{\text{upper}}$. The u rows of \mathbf{I}_u are partitioned into $\text{ceil}(u/8)$ groups of 8 rows each. Then, for each group of 8 rows all non-zero combinations of the 8 rows are computed, resulting in $2^8 - 1 = 255$ rows (this can be done with $2^8 - 8 - 1 = 247$ exclusive-ors of rows per group, since the combinations of Hamming weight one that appear in \mathbf{I}_u do not need to be recomputed). Thus, the resulting precomputation matrix \mathbf{U}' has $\text{ceil}(u/8) \cdot 255$ rows and u columns. Note that \mathbf{U}' is not formally a part of matrix \mathbf{A} , but will be used in the fourth phase to zero out $\mathbf{U}_{\text{upper}}$.

B.8.2.5 Fourth Phase

For each of the first i rows of \mathbf{A} , for each group of 8 columns in the $\mathbf{U}_{\text{upper}}$ submatrix of this row, if the set of 8 column entries in $\mathbf{U}_{\text{upper}}$ are not all zero then the row of the precomputation matrix \mathbf{U}' that matches the pattern in the 8 columns is exclusive-ORed into the row, thus zeroing out those 8 columns in the row at the cost of exclusive-oring one row of \mathbf{U}' into the row.

After this phase \mathbf{A} is the L by L identity matrix and a complete decoding schedule has been successfully formed. Then, as explained in Section C.2.1, the corresponding decoding consisting of exclusive-ORing known encoding symbols can be executed to recover the intermediate symbols based on the decoding schedule.

The triples associated with all source symbols are computed according to B.5.2.2. The triples for received source symbols are used in the decoding. The triples for missing source symbols are used to determine which intermediate symbols need to be exclusive-ORed to recover the missing source symbols.

***** Next change *****

C.1 Registration of SDP media type "~~audio, video, or text/rtp-mbms-fec-repair~~" Protocol Identifiers for Source packet

This specification defines two new SDP protocol identifiers for source packets. As the registration rules requires these to be registered by an RFC, there will be an RFC referencing the definitions here.

Protocol identifier "UDP/MBMS-FEC/RTP/AVP" identifies a protocol combination of UDP [7], FEC source packets (see section 8.2.2.3), RTP [6] using the AVP profile [78]. This protocol identifier shall use the FMT space rules that are used for RTP/AVP.

Protocol identifier "UDP/MBMS-FEC/RTP/SAVP" identifies a protocol combination of UDP [7], FEC source packets (see section 8.2.2.3), and RTP [6] using the SAVP profile [77]. This protocol identifier shall use the FMT space rules that are used for RTP/AVP.

This media type represents the RTP payload format defined in Clause 8.2.1.4.

Type name: application

Subtype name: rtp-mbms-fec-repair

Required parameters:

~~FEID: The FEC Encoding ID used in this instantiation. Expressed either as an integer or a string without white space.~~

~~min-buffer-time: This FEC buffering attribute specifies the minimum RTP receiver buffer time (delay) needed to ensure that FEC repair has time to happen regardless of the FEC source block of the stream from which the reception starts. The value is in milliseconds and represents the walleck time between the reception of the first FEC source or repair packet of a FEC source block, whichever is earlier in transmission order, and the walleck time when media decoding can safely start.~~

Optional parameters:

~~FID: The FEC Instance ID used in this instantiation. Expressed either as an integer or a string without white space. Parameter shall be present for all FEC encoding IDs that are not fully specified.~~

~~FOTI: The additional FEC object transmission information if any that the FEC instantiation uses as defined by the FEID and FID. The content is expected to be binary data but is not necessarily so, however it shall always be BASE64 encoded in the parameter value.~~

~~buf-size: This FEC buffering attribute specifies the actual memory requirement for the size of the hypothetical FEC decoding buffer that is sufficient for correct reception of the stream. The parameter can be used for optimizing the memory allocation in receivers or specifying buffer size share for the hypothetical FEC decoding buffers of different media explicitly. The parameter value is expressed in terms of number of bytes.~~

Encoding considerations:

~~The binary parameter FOTI shall be encoded using BASE 64. The RTP payload format is a binary one, however as it is restricted to usage over RTP, no special considerations are needed.~~

Restrictions on usage:

~~This format is only defined for transfer over RTP RFC3550.~~

Security considerations:

~~This format carries protection data and instructions used in FEC decoding. Thus alteration or insertion of packets can cause wide spread corruption of recovered data. Thus authentication and integrity protection of the format is appropriate. See also clause 7 of RFC 3452.~~

Interoperability considerations:

~~The greatest interoperability issue with this format is that it virtually supports any systematic FEC encoding scheme. Thus the issue is to ensure that both sender and receiver are capable of using the same FEC codes.~~

Published specification:

~~3GPP Technical specification TS 26.346~~

Applications which use this media type:

~~MBMS terminals capable of receiving the MBMS streaming delivery method.~~

Additional information:

Magic number(s): N/A

File extension(s): N/A

~~Macintosh File Type Code(s): N/A~~

~~Person & email address to contact for further information:~~

~~Magnus Westerlund (magnus.westerlund@ericsson.com)~~

~~Intended usage:~~

~~COMMON~~

~~Author:~~

~~3GPP SA4~~

~~Change controller:~~

~~3GPP TSG SA~~

C.2 Registration of [SDP Protocol identifier for repair packets](#) ~~media type "audio, video, or text/rtp-mbms-fec-source"~~

~~This media type represents the RTP payload format defined in Clause 8.1.2.4.~~

~~Type name: audio, video, or text~~

~~Subtype name: rtp-mbms-fec-source~~

~~Required parameters:~~

~~opt: — The original payload type (OPT) of the media that is tagged by this instantiation of the format. This is an unsigned integer which range depends on the RTP profile in use, but commonly 0-127.~~

~~rate: — An integer value, equal to the RTP timestamp rate value for the payload type specified by "opt".~~

~~FEID: — The FEC Encoding ID used in this instantiation. Expressed either as an integer or a string without white space.~~

~~Optional parameters:~~

~~FHID: — The FEC Instance ID used in this instantiation. Expressed either as an integer or a string without white space. Parameter shall be present for all FEC encoding IDs that are not fully specified.~~

~~FOTI: — The additional FEC object transmission information if any that the FEC instantiation uses as defined by the FEID and FHID. The content is expected to be binary data but is not necessarily so, however it shall always be BASE64 encoded in the parameter value.~~

~~Encoding considerations:~~

~~This format is a binary one, however as it is restricted to usage over RTP, no special considerations are needed.~~

~~Restrictions on usage:~~

~~This type is only defined for transfer over RFC3550.~~

~~Security considerations:~~

~~This format carries source data and instructions used in FEC decoding. Thus alteration or insertion of packets can cause wide spread corruption of recovered data. Thus authentication and integrity protection of the format is appropriate. See also clause 7 of RFC 3452.~~

~~Interoperability considerations:~~

~~The greatest interoperability issue with this format is that it supports any FEC encoding that follows the rules of RFC 3452. Thus the issue is to ensure that both sender and receiver are capable of using the same FEC codes.~~

~~Published specification:~~

~~3GPP Technical specification TS 26.346~~

~~Applications which use this media type:~~

~~MBMS terminals capable of receiving streaming media.~~

~~Additional information:~~

~~Magic number(s): N/A~~

~~File extension(s): N/A~~

~~Macintosh File Type Code(s): N/A~~

~~Person & email address to contact for further information:~~

~~Magnus Westerlund (magnus.westerlund@ericsson.com)~~

~~Intended usage:~~

~~COMMON~~

~~Author:~~

~~3GPP SA4~~

~~Change controller:~~

~~3GPP TSG SA~~

[This specification defines one new SDP protocol identifier for FEC repair packets. As the registration rules requires these to be registered by an RFC, there will be an RFC referencing the definitions within this specification.](#)

[Protocol identifier "UDP/MBMS-REPAIR" identifies a protocol combination of UDP \[7\], FEC repair packets \(see section 8.2.2.4\). The FMT string is not used and shall be set to "*".](#)

***** End of changes *****