3GPP TSG SA WG3 Security — S3#36 Shenzhen, China, November 23-26, 2004

Source: Siemens

Title: Clarification to VGCS/VBS ciphering mechanism

Document for: Approval

Agenda Item: VGCS

1 Coversheet

This contribution includes two attachments

Attachment 1 is for information and contains green marked changes with respect to CR002rev1 of TS 43.020 (which was approved in Malta).

Attachment 2 (CR002rev2 of TS 43.020) contains a clean version and is presented for approval and when approved will replace S3-040872 (CR002rev1).

Consequences if not approved:

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Clauses affected:	器 Annex F, Annex G				
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Other comments:	æ				

★ Contradictory text will stay in or missing information

*** Begin of change ****

Annex F (normative): Ciphering of Voice Group Call Service (VGCS) and Voice Broadcast Service (VBS)

This Annex defines the security related service and functions for VGCS and VBS in order to provide confidentiality protection to the group calls.

All data variables in this Annex are presented with the most significant substring on the left hand side and the least significant substring on the right hand side. A substring may be a bit, byte or other arbitrary length bitstring. Where a variable is broken down into a number of substrings, the leftmost (most significant) substring is numbered 0, the next most significant is numbered 1, and so on through to the least significant.

F.1 Introduction

F.1.1 Scope

In this Annex the ciphering of the voice group call service (VGCS) TS 42.068 [F1] and voice broadcast service (VBS) TS 42.069 [F4] is described. The following functions are required:

- Key derivation;
- Encryption of voice group/broadcast calls;
- The secure storage of the master group keys.

VGCS and VBS provide no authentication functions, i.e. authentication is performed implicitly via encryption/decryption since only a legitimate subscriber shall be able to encrypt and decrypt the VGCS/VBS speech call when the group call requires confidentiality protection. To include a subscriber into a voice group the required group data (including the 2 master group keys) shall be stored on the USIM, e.g. during the personalisation process or via OTA (over-the-air). To exclude a subscriber from a voice group the group data shall be deleted from the USIM. In case of a stolen or lost USIM, all USIMs of the remaining members of the voice groups that the USIM is a member of, need to be changed (e.g. via OTA or manual provisioning).

A pre-Rel-6 VGCS/VBS capable mobile shall be able to participate in an un-ciphered group call, if it is part of that group.

NOTE: The only security relevant difference between VBS and VGCS is that in the case of VBS there exists no uplink channel.

F.1.2 References

- [F1] 3GPP TS 42.068: "3rd Generation Partnership Project; Technical Specification Group Services and system Aspects; Voice Group Call Service (VGCS) Stage 1".
- [F2] 3GPP TS 43.068: "3rd Generation Partnership Project; Technical Specification Group Services and system Aspects; Voice Group Call Service (VGCS) Stage 2".
- [F3] 3GPP TS 31.102: "3rd Generation Partnership Project; Technical Specification Group Terminals; Characteristics of the USIM application".

[F4]	3GPP TS 42.069: "3rd Generation Partnership Project; Technical Specification Group Services and system Aspects; Voice Broadcast Service (VBS) - Stage 1".
[F5]	3GPP TS 43.069: "3rd Generation Partnership Project; Technical Specification Group Services and system Aspects; Voice Broadcast Service (VBS) - Stage 2".
[F6]	3GPP TS 23.003: "3rd Generation Partnership Project; Technical Specification Group Core Network; Numbering, addressing and identification".
[F7]	FIPS PUB 180-1 Secure Hash Standard.

F.1.3 Definitions and Abbreviations

F.1.3.1 Definitions

A5 Id: Identifier of the encryption algorithm which shall be used.

CELL_GLOBAL_COUNT: A counter valid for all voice group calls within a cell.

Group_Id: Unique identifier of a voice call group.

KMF: Key Modification Function. KMF derives from the short term key VSTK, the CGI and the CELL_GLOBAL_COUNT the cipher key V_Kc which is valid for that specific cell.

VSTK: Short Term Key provided by the USIM and the GCR. VSTK is derived from VSTK_RAND and V_Ki (128 bit).

VK_Id: Identifier of the Master Group Key (1 bit) of a group. There are up to 2 V_Ki per group.

VSTK_RAND: The 36-bit value that is used for derivation of a short term key VSTK.

V_Ki (**Group_Id**, **VK_Idi**): Voice Group or Broadcast Group Key (128 bit) number i::=VK_Id of group with Group_Id. This is also called Master Group Key or Group Key in this Annex.

V_Kc: Voice Group or Broadcast Ciphering Key (128 bit). V_Kc is derived from VSTK.

F.1.3.2 Abbreviations

The following list describes the abbreviations and acronyms used in this Annex.

CGI Cell Global Identifier
GCR Group Call Register
VBS Voice Broadcast Service
VGCS Voice Group Call Service

F.2 Security Requirements

The ciphering concept for VGCS, VBS fulfils following security requirements:

REQ-1: Prevent the same Voice group or Broadcast group ciphering key being used within different cells.

This requirement protects an observer of getting more information on the plaintext if different data is enciphered with the same key and COUNT (TDMA-numbers derived) in different cells.

REQ-2: The master group key shall never leave the USIM and the GCR.

Even though VGCS/VBS users should be trusted, this approach protects the 'root'-key (i.e. Master Group key) in the most secure way such that it need not be updated very frequently.

REQ-3: Prevent the reuse of COUNT with the same voice group or broadcast group ciphering key within the same cell.

The COUNT value is determined by the TDMA frame number. An overflow happens after each 3 hour and 8 minutes period. The lifetime of the used cipher key shall not be longer than the overflow period.

NOTE: This enhancement goes beyond the provided level of security of GSM-calls over a point to point channel (i.e. is not a VGCS/VBS-problem only) as long standing calls over a dedicated channel have the same characteristic of reusing the COUNT.

REQ-4: Prevent the same key stream block being used in uplink and downlink direction.

This requirement is fulfilled by Point to Point voice calls already (see clause C.1.2). By reusing the same mechanisms for uplink/downlink key stream derivation (i.e. reusing A5) the VBS/VGCS ciphering also fulfils this requirement.

F.3 Storage of the Master Group Keys and overview of flows

The master group keys (in short called group keys in this Annex) are securely stored at two locations:

- GCR: Beside other information, the GCR stores for each Group_Id a list of group keys. Each group key is uniquely identified by the Group_Id and the group key number VK_Id-(1-2);
- USIM: The USIM contains a list of 2 group keys for each Group_Id. Deletion or changing of group keys are allowed only via OTA or via USIM-personalisation.

The Short Term Key VSTK shall be deleted by the network entities after tearing down the call and by the ME on power down or UICC removal. On each new VGCS/VBS call set up, a new short term key VSTK shall be generated.

The following sequence gives an overview of how the different network entities make use of the group keys (and derived information) during the establishment of a voice group/broadcast call:

- 1. during the voice group/broadcast call set up the anchor MSC sends a GCR Interrogation to the GCR containing the Group_Id;
- 2. the GCR provides on the basis of a fresh number VSTK_RAND (see Annex G) the key VSTK as described in Annex F.4. VK_Id, VSTK_RAND, VSTK, the permitted ciphering algorithm (A5_Id) and other voice group/broadcast call related information, are sent from the GCR back to the anchor MSC;
- 3. the anchor MSC sends this information to the relay MSC's via a MAP operation;
- 4. the anchor MSC and relay MSC's sends this information to the BSS using the VGCS Assignment Request or VBS Assignment Request;
- 5. the BSS sends the CELL_GLOBAL_COUNT, VSTK_RAND, Group_Id and the group key number VK_Id to the ME's via a notification procedure;
- 6. each ME generates the VSTK, on the basis of the received information from step 5, as described in clause F.4.

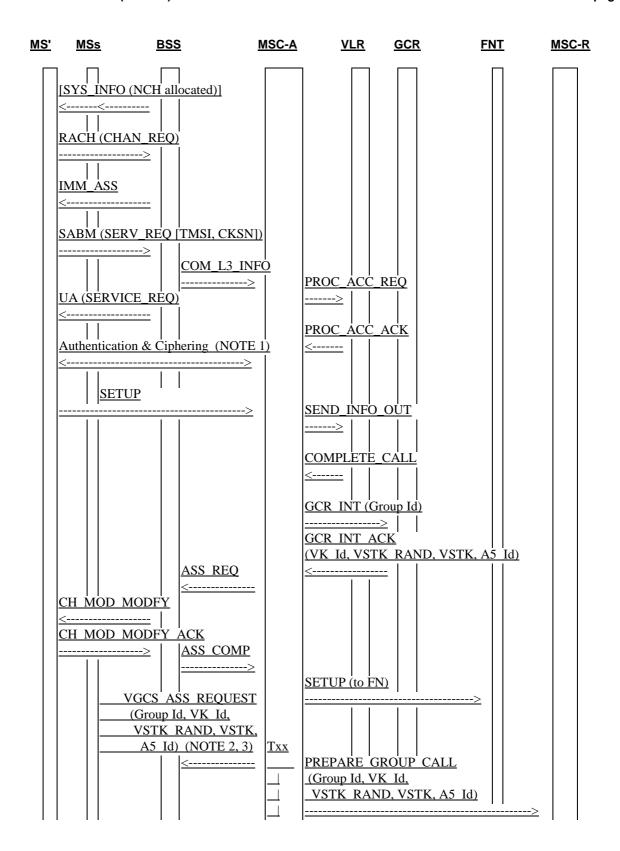
A late entrant belonging to the right Group_Id in a cell where a call is active need to pick out the notification parameters from step 5 and executes step 6.

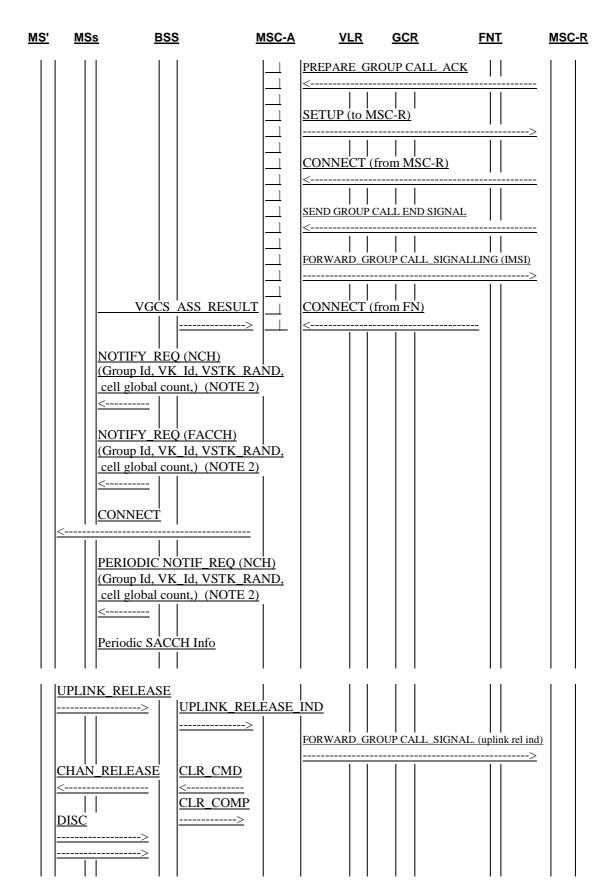
In case of inter MSC Handover of the talking subscriber the Group_Id, VSTK_RAND, VSTK and A5_Id need to be transferred via MAP Prepare Handover request message from MSC A to MSC B.

F.3.1 Distribution of ciphering data during establishment of a voice/broadcast group call

This signalling flow indicates the distribution of the VGCS parameters during the establishment of a ciphered voice group call. Figure F.3.1-1 shows the distribution of the VSTK RAND, VSTK, VK Id, A5 id and Cell Global Count between MSC, BSC and MS. The main points are:

- The Notification/NCH and Notification/FACCH are used to transfer the VSTK_RAND, VK_Id and Cell Global Count between the BSS and the MS.
- The PREPARE_GROUP_CALL is used to transfer the VSTK, VSTK_RAND, VK_Id and A5_Id between MSC-A and MSC-B.
- The VGCS/VBS Assignment Request transfers the VSTK, VSTK RAND, VK Id and A5 Id between the MSC and the BSC.





NOTE 1: If authentication and ciphering are performed, then the dedicated channel of the originator of the voice group call is ciphered with the cipher key Kc generated during the authentication procedure. If ciphering is started without authentication, the cipher key indicated with CKSN in the Service Request message is used.

- NOTE 2: The Group Id and the Group cipher key number (VK Id) are included in the Descriptive group call
- NOTE 3: The permitted ciphering algorithm (A5 Id) is included in the Encryption information.

= calling subscriber mobile station;

MSs = destination subscriber mobile stations;

FNT = fixed network user terminal;

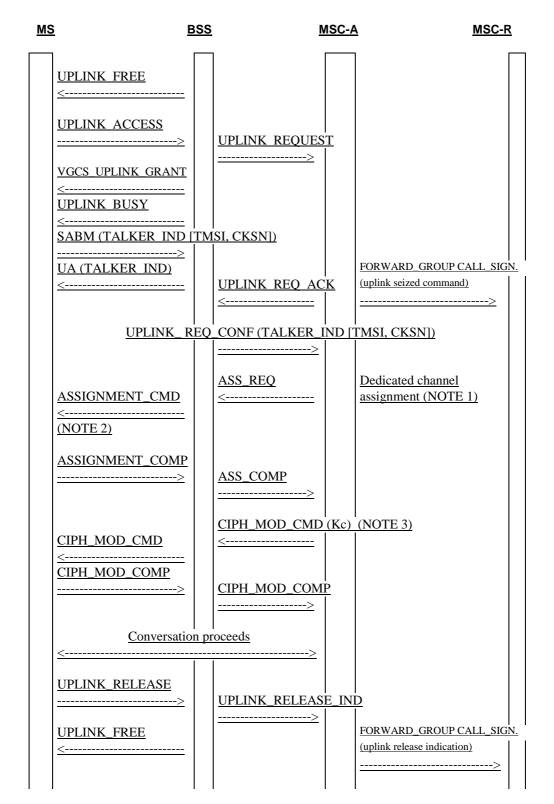
MSC-A = anchor MSC; MSC-R = relay MSC.

Figure F.3.1-1: Distribution of ciphering data during establishment of a voice group call.

F.3.2 Signalling information required for the voice group call uplink access in the anchor MSC (normal case, subsequent talker on dedicated channel)

Figure F.3.2-1 shows how the MS and the BSC determine the Cipher Key Sequence Number and Ciphering algorithm to use when the VGCS talker is on a dedicated channel. The main points are:

- The MS reads the Cipher Key Sequence Number from the USIM and passes the value to the BSC via the **TALKER INDICATION Message**
- The Cipher Key Sequence Number is passed from the BSC to the MSC via the UPLINK REQUEST CONFIRMATION message (within Layer 3 information).
- The MS and BSC are informed of the ciphering algorithm identity in the CIPHER MODE COMMAND message.



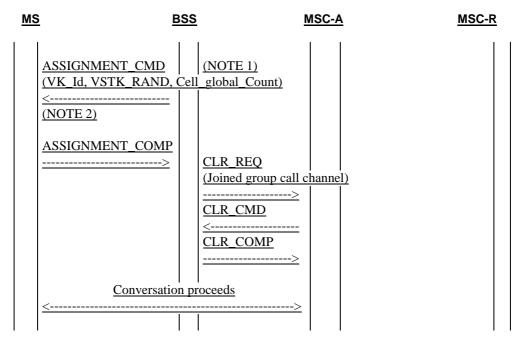
- NOTE 1: In this case the MSC decided to transfer the subsequent talker to a dedicated channel.
- NOTE 2: Upon reception of the ASSIGNMENT CMD message which transfers the MS from the group call channel to a dedicated channel, the MS starts transmission and reception on the dedicated channel in unciphered mode.
- NOTE 3: The dedicated channel of the subsequent talker is ciphered with the cipher key Kc indicated with CKSN in the Talker Indication message.

Figure F.3.2-1: Signalling information required for the voice group call uplink access in the anchor MSC (normal case, subsequent talker on dedicated channel)

F.3.3 Signalling information required to transfer the originator or subsequent talker from a dedicated channel to a group call channel

Figure F.3.3-1 shows the MS being transferred from a dedicated channel to the group channel via the ASSIGNMENT COMMAND message. The main points are:

- The group channel is ciphered with VGCS ciphering
- The VK Id, VSTK RAND and Cell Global Count are supplied in the ASSIGNMENT COMMAND message in order for the MS to calculate the voice group ciphering keys.



NOTE 1: In this case the BSC decided to transfer the originator or subsequent talker to a group call channel.

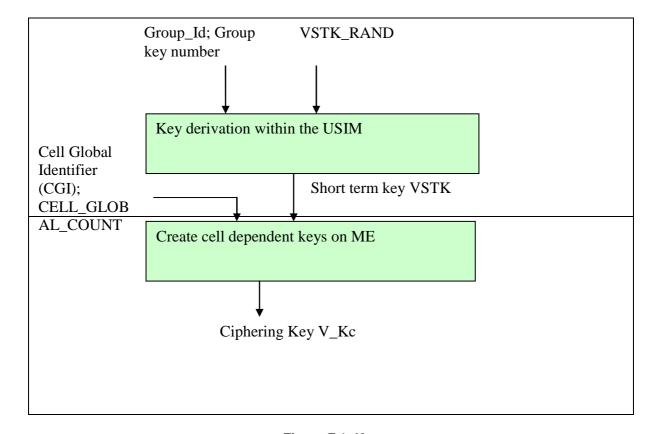
NOTE 2: Upon reception of the ASSIGNMENT CMD message, if the Group cipher key number is different from 'no ciphering', the MS derives the cipher key V_Kc and starts transmission and reception on the group call channel in ciphered mode, using V_Kc.

<u>Figure F.3.3-1: Signalling information required to transfer the originator or subsequent talker from a</u> dedicated channel to a group call channel

F.4 Key derivation

The key derivation of the encryption is performed in two steps:

- 1. derivation of a short term key VSTK on the GCR-side and USIM; VSTK_RAND generation on the GCR-side and sending it to the ME via the BSS for use on the USIM;
- 2. derivation of the actual encryption key V_Kc in the BSS and ME.



Group_Id; VK_Id VSTK_RAND

Cell Global Identifier (CGI);
CELL_GLOB AL_COUNT

Create cell dependent keys on ME

Ciphering Key V_Kc

Figure F.1: Key derivation

F.4.1 Key derivation within the USIM / GCR

This function is performed on:

- the set-up of a voice group or broadcast call by the GCR;
- entry to a voice group or broadcast call by the USIM.

On the set-up of a voice group/broadcast call the GCR generates the VSTK_RAND (See Annex G). Also an appropriate group key V_Ki (identified by VK_Id) is selected by the GCR. Using the function A8_V a short term key VSTK is derived using as input parameters:

- V_Ki (Group_Id , VK_Id);
- VSTK_RAND.

Output of A8_V is:

VSTK

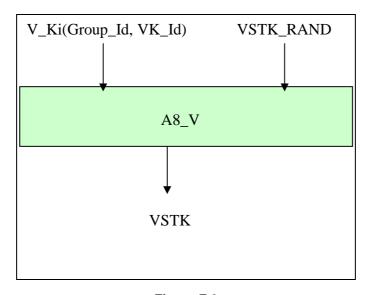


Figure F.2

The GCR sends the parameters Group_Id, VK_Id, VSTK_RAND, VSTK, A5_Id via the anchor-MSC and the relay-MSC's to the BSS. The BSS signals the Group_Id, VSTK_RAND and VK_Id to the ME.

On the ME-side, each ME sends the Group_Id of the voice group or broadcast call, the identifier of the key VK_ID and the VSTK_RAND to the USIM. The USIM performs the calculation of the short term key VSTK using the function A8_V and returns it (together with the encryption algorithm identifier A5_Id).

F.4.2 Key derivation within the ME/BSS

This function is performed by the ME on:

- entry to a voice group/broadcast call;
- cell reselection;
- changing of the value of CELL_GLOBAL_COUNT;
- Handover.

On the network side the function is performed by the BSS on

- set-up of a voice group/broadcast call in a cell;

- changing of the value of CELL_GLOBAL_COUNT.

For each cell the BSS and ME calculate an encryption key V_Kc using the key modification function KMF. Input parameter of the KMF are:

- VSTK: the short term key for this voice call group and this call;
- CGI: the cell global identifier which identifies a cell world-wide uniquely;
- CELL_GLOBAL_COUNT: this parameter shall be incremented by the BSS when the TDMA-frame-number wraps around.

NOTE: The MS and network SHALL be aligned regarding the value of the CELL_GLOBAL_COUNT. In case of transmissions on the FACCH, this requires that the network transmits a part of the whole of the TDMA frame number together with the CELL_GLOBAL_COUNT.

The output of the key modification function is the actually cipher key V_Kc.

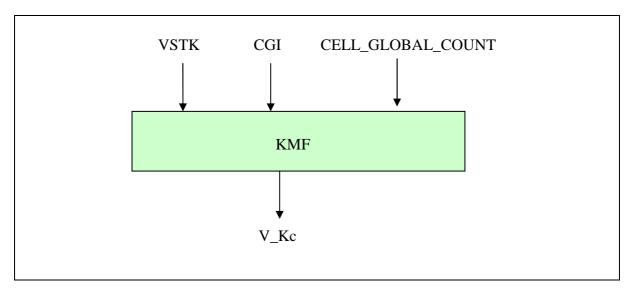


Figure F.3

To provide the required information to the ME the parameters CELL_GLOBAL_COUNT and CGI are included in various messages from the BSS to the ME (i.e. CELL_GLOBAL_COUNT on the NCH, FACCH and PCH, and the CGI on the BCCH and the FACCH).

F.4.3 Encryption algorithm selection

The encryption algorithm identifier A5_Id is stored in the GCR and the USIM. For each group key $V_Ki(Group_Id, VK_Idi)$ there is a unique A5_Id.

A5_Id is transmitted from the GCR to the BSS. The ME fetches the A5_Id together with the VSTK from the USIM.

NOTE 1: It is possible that different algorithm identifiers are bound to different V_Ki of the same group.

NOTE 2: The algorithm identifier A5_Id stored in the GCR and on the USIM shall match with the encryption capabilities of the ME's used by the group and the BSS where the voice group calls are allowed to take place.

F.4.4 Algorithm requirements

F.4.3.1 A8 V

The key derivation function A8_V has the following input and output parameter:

Input Parameter:

VSTK_RAND: 36 bit value (see annex G);

V_Ki (Group_Id, VK_Idi): 128 bit secret key;

Output:

VSTK: 128 bit short term key

A8_V is an operator specific algorithm. The calculation time for A8_V shall not exceed 500 ms.

A8_V is implemented in the GCR and on the USIM.

F.4.3.1 KMF

The key derivation function KMF has the following input and output parameter:

Input Parameter:

VSTK: 128 bit short term key;

CGI: the cell global identifier: 56 bit (TS 23.003 [F6]);

CELL_GLOBAL_COUNT: 2 bit.

Output:

V_Kc 128 bit encryption key.

The KMF is implemented in the BSS and in the ME.

The specification of KMF can be found in clause F.6

F.5 Encryption of voice group calls

For the encryption of a voice group call the same encryption algorithms are used as for a normal GSM speech call. Which algorithm out of the algorithm suite A5/x is used is determined by the identifier $A5_Id$, which is stored on the USIM (together with the group key $V_Ki(Group_Id, VK_Id\frac{1}{2})$). The algorithm A5/X is used in the same way as in the GSM (see clause -C.1) using the key V_Kc as encryption/decryption key Kc as input to A5/x.

If the key length KL of the encryption algorithm A5/X is shorter than the length of V_Kc (128 bit) then only bits [0] to [KL-1] the KL least significant KL bits of V_Kc are used.

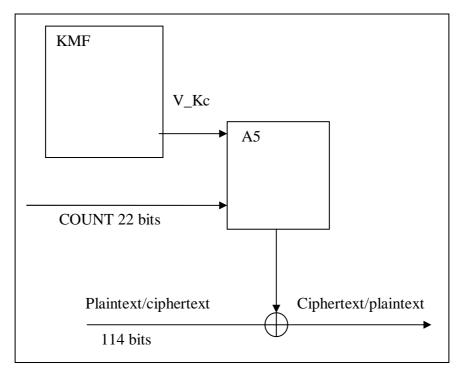


Figure F.4

F.6 Specification of the Key Modification Function (KMF)

SHA-1 (FIPS PUB 180-1 [F7]) is used for generating V_Kc:

V_Kc= SHA-1(VSTK | CGI | CELL_GLOBAL_COUNT | VSTK)

From the 160 bit output of SHA-1, the $\frac{128}{\text{bits}}$ numbered as [0] to [127] least significant bits are taken as 128 bit V_Kc.

Annex G (informative): Generation of VSTK_RAND

All data variables in this Annex are presented with the most significant substring on the left hand side and the least significant substring on the right hand side. A substring may be a bit, byte or other arbitrary length bitstring. Where a variable is broken down into a number of substrings, the leftmost (most significant) substring is numbered 0, the next most significant is numbered 1, and so on through to the least significant.

Since the length of VSTK_RAND (36 bits) is small, care should be taken that a VSTK_RAND isn't generated twice (so-called collision) during the lifetime of V_Ki. On the other hand, the predictability of VSTK_RAND shall be avoided. The following scheme could be used in order to generate 4096 VSTK_RAND for each V_Ki with a probability $< 10^{-6}$ that a collision occurs.

NOTE: A collision probability of <10⁻⁴ could still give a sufficient security margin and may allow, depending on the VSTK_RAND structure that is chosen, that more VSTK can be generated from one V_Ki.

The GCR maintains a COUNTER (12 bits) for each voice group. After each generation of a VSTK_RAND for a specific voice group, COUNTER for that voice group is incremented by one.

The left most 12 bits (COUNTER) of VSTK_RAND are set to COUNTER. The remaining <u>right most</u> 24 bits (RANDOM) are generated randomly, i.e. unpredictably for each new VSTK_RAND.

Therefore VSTK_RAND = COUNTER | RANDOM.

NOTE: For security reasons, any adopted scheme shall contain at least 24 true random bits. The length of RANDOM shall be at least 24 bits.

If COUNTER wraps around, a new V_Ki is required for that group.

Table G.1 gives the maximum number of voice group calls that are possible with a with a full random generated VSTK_RAND:

Table G.1: Maximum number of voice group calls that are possible with a with a full random generated VSTK RAND

Length of VSTK_RAND	Max collision prob for fixed V_Ki	Number of calls
36	10 ⁻⁶	TBD 371
36	10 ⁻⁴	3707 TBD

Table G.2 gives the maximum number of voice group calls that are possible with a VSTK_RAND, as structured in this annex.

Table G.2: Maximum number of voice group calls that are possible with a VSTK_RAND

Total challenge length	Length of counter	Length of random part	Max collision prob for fixed V_Ki	Max collision prob for one fixed counter	Number of calls for one fixed counter	Total number of calls for fixed V_Ki
36	1 <u>2</u> 4	24	10 ⁻⁶		1	4096
36	1 <u>2</u> 4	24	10 ⁻⁴	26.4410 × 10	1	4096

Explanation of the columns of table G.2:

Max collision probability for fixed V_K i: what we have determined, for security reasons, should be the maximum probability that the same value of VSTK RAND (and hence the same value of VSTK) is used twice before the value of V_K i is changed. 10^{-6} is a strong security setting; 10^{-4} is not quite so strong, but probably adequate.

Max collision probability for one fixed counter: suppose that VSTK RAND is made up of N c counter bits and N r random bits. We assume that the counter part will take all possible 2^{N_c} values before V Ki is updated. Having selected

 $\frac{\text{our required "Max collision prob for fixed V_Ki", this is the corresponding maximum permitted probability that the }{\text{same value of the N_r random bits (and hence the same value of VSTK) is used twice for a fixed value of the N_c }{\text{counter bits.}}$

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Clauses affected:

Other comments:

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Other specs affected:

*** Begin of change ****

Annex F (normative): Ciphering of Voice Group Call Service (VGCS) and Voice Broadcast Service (VBS)

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A pre-Rel-6 VGCS/VBS capable mobile shall be able to participate in an un-ciphered group call, if it is part of that group.

NOTE: The only security relevant difference between VBS and VGCS is that in the case of VBS there exists no uplink channel.

F.1.2 References

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- [F3] 3GPP TS 31.102: "3rd Generation Partnership Project; Technical Specification Group Terminals; Characteristics of the USIM application".

[F4]	3GPP TS 42.069: "3rd Generation Partnership Project; Technical Specification Group Services and system Aspects; Voice Broadcast Service (VBS) - Stage 1".
[F5]	3GPP TS 43.069: "3rd Generation Partnership Project; Technical Specification Group Services and system Aspects; Voice Broadcast Service (VBS) - Stage 2".
[F6]	3GPP TS 23.003: "3rd Generation Partnership Project; Technical Specification Group Core Network; Numbering, addressing and identification".
[F7]	FIPS PUB 180-1 Secure Hash Standard.

F.1.3 Definitions and Abbreviations

F.1.3.1 Definitions

A5 Id: Identifier of the encryption algorithm which shall be used.

CELL_GLOBAL_COUNT: A counter valid for all voice group calls within a cell.

Group_Id: Unique identifier of a voice call group.

KMF: Key Modification Function. KMF derives from the short term key VSTK, the CGI and the CELL_GLOBAL_COUNT the cipher key V_Kc which is valid for that specific cell.

VSTK: Short Term Key provided by the USIM and the GCR. VSTK is derived from VSTK_RAND and V_Ki (128 bit).

VK_Id: Identifier of the Master Group Key (1 bit) of a group. There are up to 2 V_Ki per group.

VSTK_RAND: The 36-bit value that is used for derivation of a short term key VSTK.

V_Ki (**Group_Id**, **VK_Idi**): Voice Group or Broadcast Group Key (128 bit) number i::=VK_Id of group with Group_Id. This is also called Master Group Key or Group Key in this Annex.

V_Kc: Voice Group or Broadcast Ciphering Key (128 bit). V_Kc is derived from VSTK.

F.1.3.2 Abbreviations

The following list describes the abbreviations and acronyms used in this Annex.

CGI Cell Global Identifier
GCR Group Call Register
VBS Voice Broadcast Service
VGCS Voice Group Call Service

F.2 Security Requirements

The ciphering concept for VGCS, VBS fulfils following security requirements:

REQ-1: Prevent the same Voice group or Broadcast group ciphering key being used within different cells.

This requirement protects an observer of getting more information on the plaintext if different data is enciphered with the same key and COUNT (TDMA-numbers derived) in different cells.

REQ-2: The master group key shall never leave the USIM and the GCR.

Even though VGCS/VBS users should be trusted, this approach protects the 'root'-key (i.e. Master Group key) in the most secure way such that it need not be updated very frequently.

REQ-3: Prevent the reuse of COUNT with the same voice group or broadcast group ciphering key within the same cell.

The COUNT value is determined by the TDMA frame number. An overflow happens after each 3 hour and 8 minutes period. The lifetime of the used cipher key shall not be longer than the overflow period.

NOTE: This enhancement goes beyond the provided level of security of GSM-calls over a point to point channel (i.e. is not a VGCS/VBS-problem only) as long standing calls over a dedicated channel have the same characteristic of reusing the COUNT.

REQ-4: Prevent the same key stream block being used in uplink and downlink direction.

This requirement is fulfilled by Point to Point voice calls already (see clause C.1.2). By reusing the same mechanisms for uplink/downlink key stream derivation (i.e. reusing A5) the VBS/VGCS ciphering also fulfils this requirement.

F.3 Storage of the Master Group Keys and overview of flows

The master group keys (in short called group keys in this Annex) are securely stored at two locations:

- GCR: Beside other information, the GCR stores for each Group_Id a list of group keys. Each group key is uniquely identified by the Group_Id and the group key number VK_Id-(1-2);
- USIM: The USIM contains a list of 2 group keys for each Group_Id. Deletion or changing of group keys are allowed only via OTA or via USIM-personalisation.

The Short Term Key VSTK shall be deleted by the network entities after tearing down the call and by the ME on power down or UICC removal. On each new VGCS/VBS call set up, a new short term key VSTK shall be generated.

The following sequence gives an overview of how the different network entities make use of the group keys (and derived information) during the establishment of a voice group/broadcast call:

- 1. during the voice group/broadcast call set up the anchor MSC sends a GCR Interrogation to the GCR containing the Group_Id;
- 2. the GCR provides on the basis of a fresh number VSTK_RAND (see Annex G) the key VSTK as described in Annex F.4. VK_Id, VSTK_RAND, VSTK, the permitted ciphering algorithm (A5_Id) and other voice group/broadcast call related information, are sent from the GCR back to the anchor MSC;
- 3. the anchor MSC sends this information to the relay MSC's via a MAP operation;
- 4. the anchor MSC and relay MSC's sends this information to the BSS using the VGCS Assignment Request or VBS Assignment Request;
- 5. the BSS sends the CELL_GLOBAL_COUNT, VSTK_RAND, Group_Id and the group key number VK_Id to the ME's via a notification procedure;
- 6. each ME generates the VSTK, on the basis of the received information from step 5, as described in clause F.4.

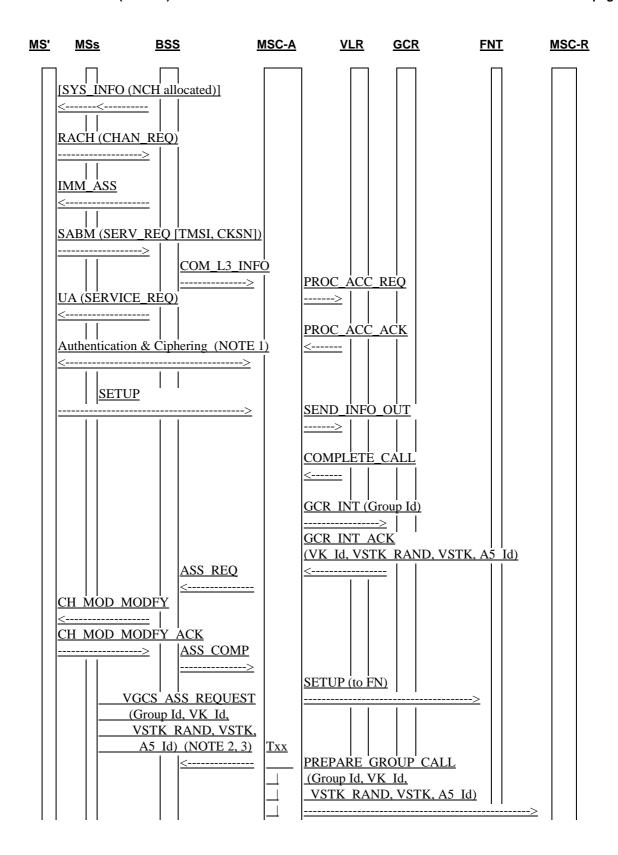
A late entrant belonging to the right Group_Id in a cell where a call is active need to pick out the notification parameters from step 5 and executes step 6.

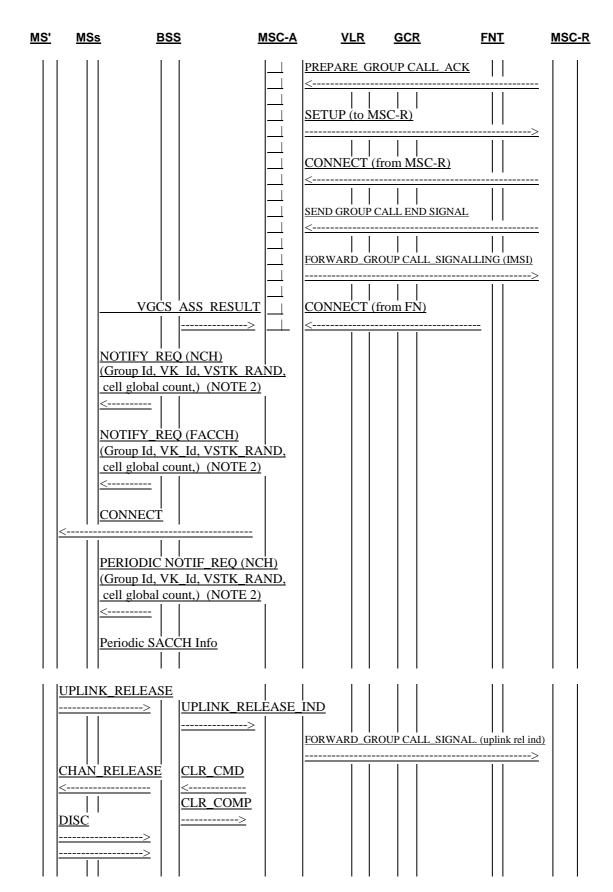
In case of inter MSC Handover of the talking subscriber the Group_Id, VSTK_RAND, VSTK and A5_Id need to be transferred via MAP Prepare Handover request message from MSC A to MSC B.

F.3.1 Distribution of ciphering data during establishment of a voice/broadcast group call

This signalling flow indicates the distribution of the VGCS parameters during the establishment of a ciphered voice group call. Figure F.3.1-1 shows the distribution of the VSTK RAND, VSTK, VK Id, A5 id and Cell Global Count between MSC, BSC and MS. The main points are:

- The Notification/NCH and Notification/FACCH are used to transfer the VSTK_RAND, VK_Id and Cell Global Count between the BSS and the MS.
- The PREPARE_GROUP_CALL is used to transfer the VSTK, VSTK_RAND, VK_Id and A5_Id between MSC-A and MSC-B.
- The VGCS/VBS Assignment Request transfers the VSTK, VSTK RAND, VK Id and A5 Id between the MSC and the BSC.





NOTE 1: If authentication and ciphering are performed, then the dedicated channel of the originator of the voice group call is ciphered with the cipher key Kc generated during the authentication procedure. If ciphering is started without authentication, the cipher key indicated with CKSN in the Service Request message is used.

- NOTE 2: The Group Id and the Group cipher key number (VK Id) are included in the Descriptive group call
- NOTE 3: The permitted ciphering algorithm (A5 Id) is included in the Encryption information.

= calling subscriber mobile station;

MSs = destination subscriber mobile stations;

FNT = fixed network user terminal;

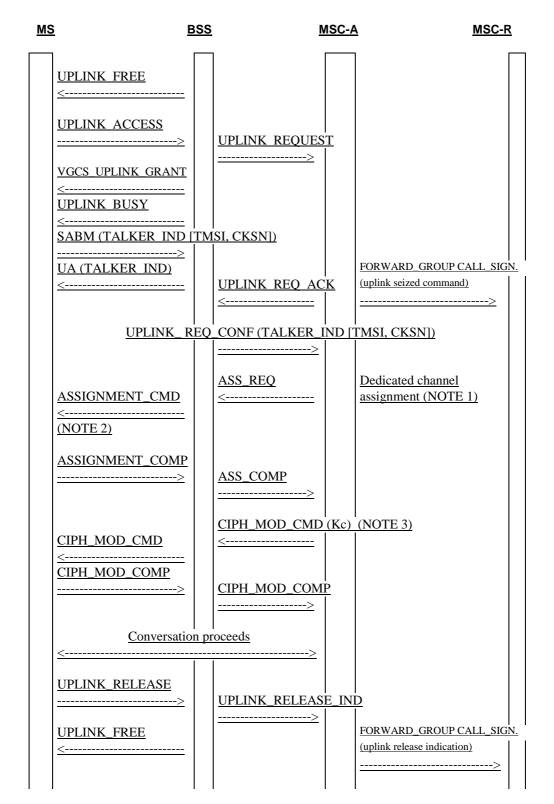
MSC-A = anchor MSC; MSC-R = relay MSC.

Figure F.3.1-1: Distribution of ciphering data during establishment of a voice group call.

F.3.2 Signalling information required for the voice group call uplink access in the anchor MSC (normal case, subsequent talker on dedicated channel)

Figure F.3.2-1 shows how the MS and the BSC determine the Cipher Key Sequence Number and Ciphering algorithm to use when the VGCS talker is on a dedicated channel. The main points are:

- The MS reads the Cipher Key Sequence Number from the USIM and passes the value to the BSC via the **TALKER INDICATION Message**
- The Cipher Key Sequence Number is passed from the BSC to the MSC via the UPLINK REQUEST CONFIRMATION message (within Layer 3 information).
- The MS and BSC are informed of the ciphering algorithm identity in the CIPHER MODE COMMAND message.



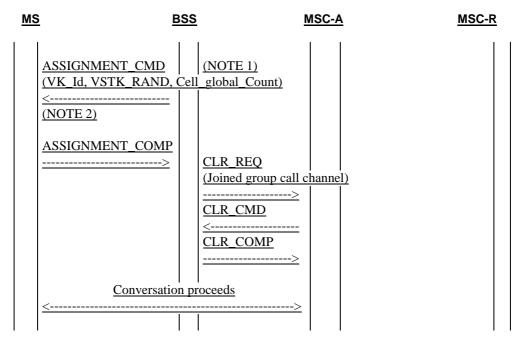
- NOTE 1: In this case the MSC decided to transfer the subsequent talker to a dedicated channel.
- NOTE 2: Upon reception of the ASSIGNMENT CMD message which transfers the MS from the group call channel to a dedicated channel, the MS starts transmission and reception on the dedicated channel in unciphered mode.
- NOTE 3: The dedicated channel of the subsequent talker is ciphered with the cipher key Kc indicated with CKSN in the Talker Indication message.

Figure F.3.2-1: Signalling information required for the voice group call uplink access in the anchor MSC (normal case, subsequent talker on dedicated channel)

F.3.3 Signalling information required to transfer the originator or subsequent talker from a dedicated channel to a group call channel

Figure F.3.3-1 shows the MS being transferred from a dedicated channel to the group channel via the ASSIGNMENT COMMAND message. The main points are:

- The group channel is ciphered with VGCS ciphering
- The VK Id, VSTK RAND and Cell Global Count are supplied in the ASSIGNMENT COMMAND message in order for the MS to calculate the voice group ciphering keys.



NOTE 1: In this case the BSC decided to transfer the originator or subsequent talker to a group call channel.

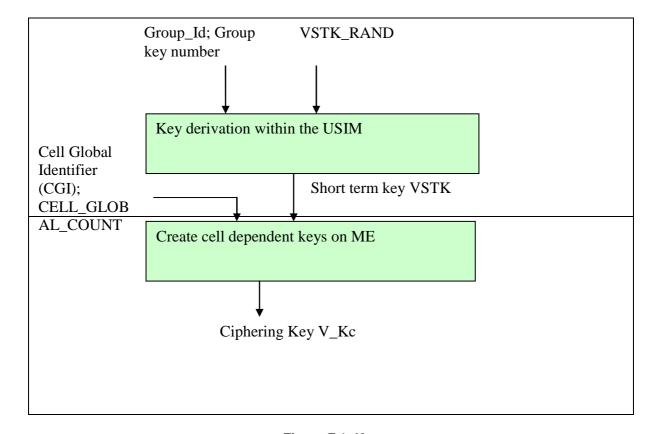
NOTE 2: Upon reception of the ASSIGNMENT CMD message, if the Group cipher key number is different from 'no ciphering', the MS derives the cipher key V_Kc and starts transmission and reception on the group call channel in ciphered mode, using V_Kc.

<u>Figure F.3.3-1: Signalling information required to transfer the originator or subsequent talker from a</u> dedicated channel to a group call channel

F.4 Key derivation

The key derivation of the encryption is performed in two steps:

- 1. derivation of a short term key VSTK on the GCR-side and USIM; VSTK_RAND generation on the GCR-side and sending it to the ME via the BSS for use on the USIM;
- 2. derivation of the actual encryption key V_Kc in the BSS and ME.



Cell Global Identifier (CGI);
CELL_GLOB AL_COUNT

Ciphering Key V_Kc

Group_Id; VK_Id VSTK_RAND

Key derivation within the USIM

Short term key VSTK

Create cell dependent keys on ME

Figure F.1: Key derivation

F.4.1 Key derivation within the USIM / GCR

This function is performed on:

- the set-up of a voice group or broadcast call by the GCR;
- entry to a voice group or broadcast call by the USIM.

On the set-up of a voice group/broadcast call the GCR generates the VSTK_RAND (See Annex G). Also an appropriate group key V_Ki (identified by VK_Id) is selected by the GCR. Using the function A8_V a short term key VSTK is derived using as input parameters:

- V_Ki (Group_Id , VK_Id);
- VSTK_RAND.

Output of A8_V is:

VSTK

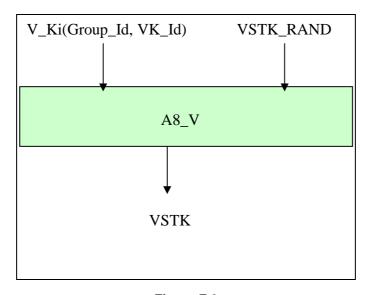


Figure F.2

The GCR sends the parameters Group_Id, VK_Id, VSTK_RAND, VSTK, A5_Id via the anchor-MSC and the relay-MSC's to the BSS. The BSS signals the Group_Id, VSTK_RAND and VK_Id to the ME.

On the ME-side, each ME sends the Group_Id of the voice group or broadcast call, the identifier of the key VK_ID and the VSTK_RAND to the USIM. The USIM performs the calculation of the short term key VSTK using the function A8_V and returns it (together with the encryption algorithm identifier A5_Id).

F.4.2 Key derivation within the ME/BSS

This function is performed by the ME on:

- entry to a voice group/broadcast call;
- cell reselection;
- changing of the value of CELL_GLOBAL_COUNT;
- Handover.

On the network side the function is performed by the BSS on

- set-up of a voice group/broadcast call in a cell;

- changing of the value of CELL_GLOBAL_COUNT.

For each cell the BSS and ME calculate an encryption key V_Kc using the key modification function KMF. Input parameter of the KMF are:

- VSTK: the short term key for this voice call group and this call;
- CGI: the cell global identifier which identifies a cell world-wide uniquely;
- CELL_GLOBAL_COUNT: this parameter shall be incremented by the BSS when the TDMA-frame-number wraps around.

NOTE: The MS and network SHALL be aligned regarding the value of the CELL_GLOBAL_COUNT. In case of transmissions on the FACCH, this requires that the network transmits a part of the whole of the TDMA frame number together with the CELL_GLOBAL_COUNT.

The output of the key modification function is the actually cipher key V_Kc.

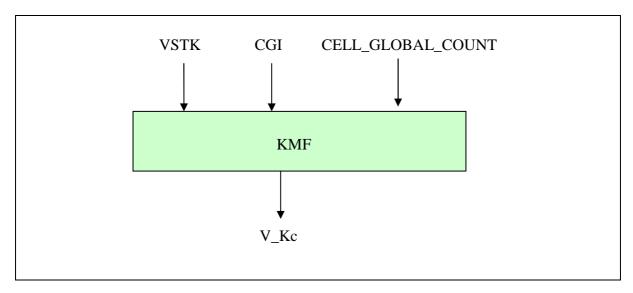


Figure F.3

To provide the required information to the ME the parameters CELL_GLOBAL_COUNT and CGI are included in various messages from the BSS to the ME (i.e. CELL_GLOBAL_COUNT on the NCH, FACCH and PCH, and the CGI on the BCCH and the FACCH).

F.4.3 Encryption algorithm selection

The encryption algorithm identifier A5_Id is stored in the GCR and the USIM. For each group key $V_Ki(Group_Id, VK_Idi)$ there is a unique A5_Id.

A5_Id is transmitted from the GCR to the BSS. The ME fetches the A5_Id together with the VSTK from the USIM.

NOTE 1: It is possible that different algorithm identifiers are bound to different V_Ki of the same group.

NOTE 2: The algorithm identifier A5_Id stored in the GCR and on the USIM shall match with the encryption capabilities of the ME's used by the group and the BSS where the voice group calls are allowed to take place.

F.4.4 Algorithm requirements

F.4.3.1 A8 V

The key derivation function A8_V has the following input and output parameter:

Input Parameter:

VSTK_RAND: 36 bit value (see annex G);

V_Ki (Group_Id, <u>VK_Id</u>): 128 bit secret key;

Output:

VSTK: 128 bit short term key

A8_V is an operator specific algorithm. The calculation time for A8_V shall not exceed 500 ms.

A8_V is implemented in the GCR and on the USIM.

F.4.3.1 KMF

The key derivation function KMF has the following input and output parameter:

Input Parameter:

VSTK: 128 bit short term key;

CGI: the cell global identifier: 56 bit (TS 23.003 [F6]);

CELL_GLOBAL_COUNT: 2 bit.

Output:

V_Kc 128 bit encryption key.

The KMF is implemented in the BSS and in the ME.

The specification of KMF can be found in clause F.6

F.5 Encryption of voice group calls

For the encryption of a voice group call the same encryption algorithms are used as for a normal GSM speech call. Which algorithm out of the algorithm suite A5/x is used is determined by the identifier $A5_Id$, which is stored on the USIM (together with the group key $V_Ki(Group_Id, VK_Id\frac{1}{2})$). The algorithm A5/X is used in the same way as in the GSM (see clause -C.1) using the key V_Kc as encryption/decryption key Kc as input to A5/x.

If the key length KL of the encryption algorithm A5/X is shorter than the length of V_Kc (128 bit) then only bits [0] to [KL-1] the KL least significant KL bits of V_Kc are used.

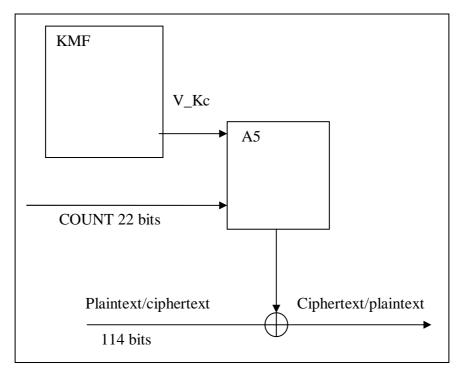


Figure F.4

F.6 Specification of the Key Modification Function (KMF)

SHA-1 (FIPS PUB 180-1 [F7]) is used for generating V_Kc:

V_Kc= SHA-1(VSTK | CGI | CELL_GLOBAL_COUNT | VSTK)

From the 160 bit output of SHA-1, the $\frac{128}{\text{bits}}$ numbered as [0] to [127] least significant bits are taken as 128 bit V_Kc.

Annex G (informative): Generation of VSTK_RAND

All data variables in this Annex are presented with the most significant substring on the left hand side and the least significant substring on the right hand side. A substring may be a bit, byte or other arbitrary length bitstring. Where a variable is broken down into a number of substrings, the leftmost (most significant) substring is numbered 0, the next most significant is numbered 1, and so on through to the least significant.

Since the length of VSTK_RAND (36 bits) is small, care should be taken that a VSTK_RAND isn't generated twice (so-called collision) during the lifetime of V_Ki. On the other hand, the predicta+bility of VSTK_RAND shall be avoided. The following scheme could be used in order to generate 4096 VSTK_RAND for each V_Ki with a probability $< 10^{-6}$ that a collision occurs.

NOTE: A collision probability of $<10^{-4}$ could still give a sufficient security margin and may allow, depending on the VSTK_RAND structure that is chosen, that more VSTK can be generated from one V_Ki.

The GCR maintains a COUNTER (12 bits) for each voice group. After each generation of a VSTK_RAND for a specific voice group, COUNTER for that voice group is incremented by one.

The left most 12 bits (COUNTER) of VSTK_RAND are set to COUNTER. The remaining <u>right most</u> 24 bits (RANDOM) are generated randomly, i.e. unpredictably for each new VSTK_RAND.

Therefore VSTK_RAND = COUNTER | RANDOM.

NOTE: For security reasons, any adopted scheme shall contain at least 24 true random bits. The length of RANDOM shall be at least 24 bits.

If COUNTER wraps around, a new V_Ki is required for that group.

Table G.1 gives the maximum number of voice group calls that are possible with a with a full random generated VSTK_RAND:

Table G.1: Maximum number of voice group calls that are possible with a with a full random generated VSTK RAND

Length of VSTK_RAND	Max collision prob for fixed V_Ki	Number of calls
36	10 ⁻⁶	TBD 371
36	10 ⁻⁴	<u>3707</u> TBD

Table G.2 gives the maximum number of voice group calls that are possible with a VSTK_RAND, as structured in this annex.

Table G.2: Maximum number of voice group calls that are possible with a VSTK_RAND

Total challenge length	Length of counter	Length of random part	Max collision prob for fixed V_Ki	Max collision prob for one fixed counter	Number of calls for one fixed counter	Total number of calls for fixed V_Ki
36	1 <u>2</u> 4	24	10 ⁻⁶		1	4096
36	1 <u>2</u> 4	24	10 ⁻⁴	26.4410 × 10	1	4096

Explanation of the columns of table G.2:

Max collision probability for fixed V_K i: what we have determined, for security reasons, should be the maximum probability that the same value of VSTK RAND (and hence the same value of VSTK) is used twice before the value of V_K i is changed. 10^{-6} is a strong security setting; 10^{-4} is not quite so strong, but probably adequate.

Max collision probability for one fixed counter: suppose that VSTK RAND is made up of N c counter bits and N r random bits. We assume that the counter part will take all possible 2^{N_c} values before V Ki is updated. Having selected

 $\frac{\text{our required "Max collision prob for fixed V_Ki", this is the corresponding maximum permitted probability that the }{\text{same value of the N_r random bits (and hence the same value of VSTK) is used twice for a fixed value of the N_c }{\text{counter bits.}}$

*** End of change ****